[RE027] China-based APT Mustang Panda might still have continued their attack activities against organizations in Vietnam

blog.vincss.net/re027-china-based-apt-mustang-panda-might-still-have-continued-their-attack-activities-against-organizations-in-vietnam/

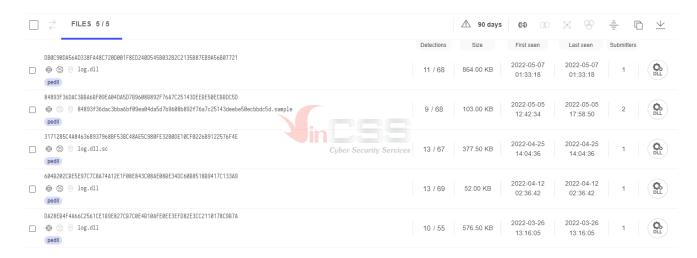
20/05/2022



1. Executive Summary

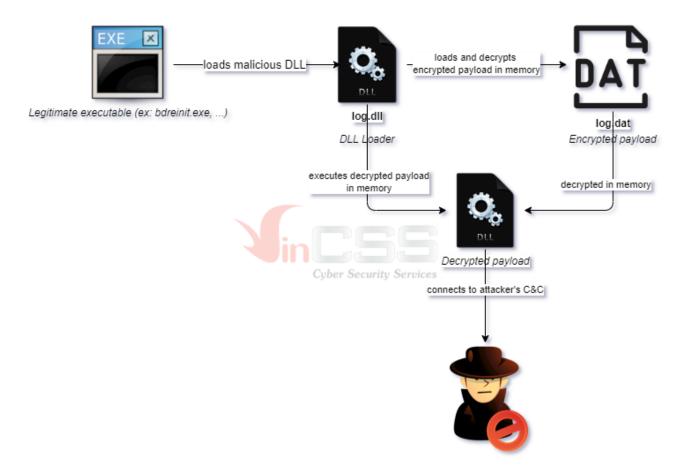
At VinCSS, through continuous cyber security monitoring, hunting malware samples and evaluating them to determine the potential risks, especially malware samples targeting Vietnam. Recently, during hunting on <u>VirusTotal's</u> platform and performing scan for specific byte patterns related to the **Mustang Panda (PlugX)**, we discovered a series of malware samples, suspected to be relevant to APT Mustang Panda, that was uploaded from Vietnam.

All of these samples share the same name as "log.dll" and have a rather low detection rate.



Based on the above information, we infer that there is a possibility that malware has been infected in certain orgs in Vietnam, so we decided to analyze these malware samples. During analysis, based on the detected indicators, we continue to investigate and set the scenario of the attack campaign.

A general overview of the execution flow demonstrated as follow:



Our blog includes:

- Technical analysis of the log.dll file.
- Technical analysis of shellcode decrypted from log.dat.
- Analyze PlugX DII as well as decrypt PlugX configuration information.

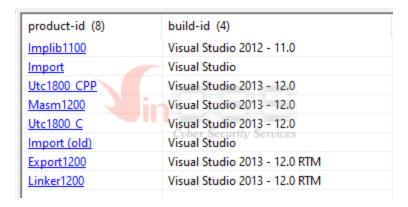
2. Analyze the log.dll

In the list of hunted samples above, we choose the one with hash: <u>3171285c4a846368937968bf53bc48ae5c980fe32b0de10cf0226b9122576f4e</u>

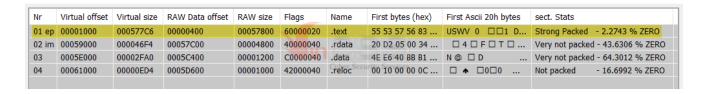
This sample was submitted to VirusTotal from Vietnam on 2022-04-25 14:04:36 UTC



The information from the Rich Header suggests that it is likely compiled with **Visual Studio 2012/2013**:



By checking the sections information, we can see that it is packed or the code is obfuscated:



Sample has the original name IjAt.dll, and it exports two functions LogFree and LogInit:

Offset	Name		Value	e Meanin	g				
5BC90	Characteris	tics	0						
5BC94	TimeDateSta	.mp	622DA	6ED Sunday	, 13.03.2022 (08:10:21 UTC			
5BC98	MajorVersio	n	0						
5BC9A	MinorVersio	n	0		_				
5BC9C	Name		5D0C0	ljAt.d	u				
5BCA0	Base		1						
5BCA4	NumberOfFun	ctions	2						
5BCA8	NumberOfNam	es	2						
5BCAC	AddressOfFu	nctions	5D0B8	5D0B8					
5BCB0	AddressOfNa	mes	5 D 0C0						
5BCB4	AddressOfNa	meOrdinals	5D0C8	ber Security	Services				
Exported	Functions [2 entries]						
Offset	Ordinal	Function	n RVA	Name RVA	Name	Forwarder			
5BCB8	1	1000		5D0D5	LogFree				
5BCBC	2	4E5E0		5D0DD	LogInit				

Load sample into IDA, analyze the code of the two functions above:

LogFree function:

Looking at this function, it can be seen that its code has been completely obfuscated by Obfuscator-LLVM, using the Control Flow Flattening technique:



After further analysis, I found that this function has no special task.

LogInit function:

This function will call the **LogInit_0** function:

```
text:1004E5E0 ; Exported entry
                                                                              2void _stdcall LogInit()
                                2. LogInit
text:1004E5E0
                                                                            4 LogInit_0();
text:1004E5E0
text:1004E5E0 ; Attributes: thunk
text:1004E5E0
text:1004E5E0 ; void _stdcall LogInit()
text:1004E5E0
                            public LogInit
text:1004E5E0 LogInit
                             proc near
text:1004E5E0
text:1004E5E0
                                            🔙 ; TAGS: ['Enum', 'FileWIN'
text:1004E5E0 LogInit
text:1004E5E0
```

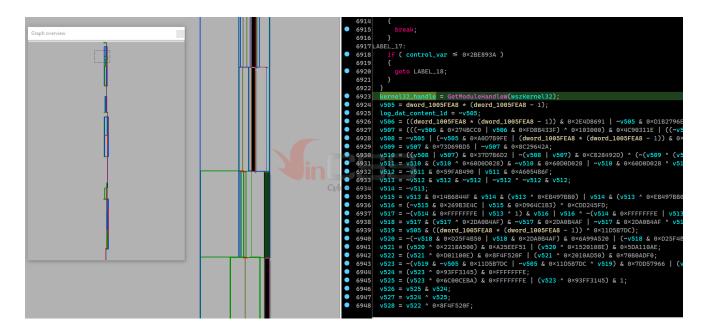
Similar to the above, the code at the **LogInit_0** function has also been completely obfuscated, it takes a long time for IDA to decompile the code of this function:



The primary task of the **LogInit_0** function is to call the function **f_read_content_of_log_dat_file_to_buf** for reading the content of **log.dat** file and execute the decrypted shellcode:

```
proc near
jmp
                23 calls, 1 strings
                calls:
                                     call dword ptr[eax]
                                     call ds:CloseHandle ; call CloseHandle call ds:CreateFileA ; call CreateFileA to open file
                                     call ds:ReadFile ; call ReadFile to read file content
                                     call _strncmp ; call _strncmp to compare string
                                    call dword ptr[eax] ; exec decrypted payload/shellcode
                                     call ds:CloseHandle ; call CloseHandle call ds:DeleteFileA ; call DeleteFileA
                                     call ds:CloseHandle ; call CloseHandle
                                     call ds:DeleteFileA ; call DeleteFileA
                                     call f_read_content_of_log_dat_file_to_buf ; call f_read_content_of_log_dat_file
                                     call ds:GetModuleHandleA ; call GetModuleHandleA to retrieve kernel32.dll handl
                                     call ds:GetProcAddress ; retrieve api address
                                     call eax ; call API func
                                     call ds:ExpandEnvironmentStringsA ; call ExpandEnvironmentStringsA
                                     call ds:CreateFileA ; call CreateFileA for retrieving handle to create tmp file
                                     call _strlen ; call _strlen
                                     call ds:WriteFile ; call WriteFile to write content to file
                                     call ds:ExpandEnvironmentStringsA; call ExpandEnvironmentStringsA
                                     call ds:CreateFileA ; call CreateFileA
                                     call _strlen ; call _strlen
                                     call ds:WriteFile ; call WriteFile
                                     call __security_check_cookie(x)
with Hex View-1. Pr
                strings:
```

f_read_content_of_log_dat_file_to_buf's code is also completely obfuscated:



The major task of this function as the following:

- Call the GetModuleHandleW function to retrieve the handle of kernel32.dll.
- Call the GetProcAddress function to get the addresses of the APIs: VirtualAlloc,
 GetModuleFileNameA, CreateFileA, ReadFile.

Use the above APIs to retrieve the path to the log.dat file and read the contents of this
file into the allocated memory.

```
control_var = 0x/A/CA244;

mov ecx, [eba+decrunted_shallcode]

mov edx, 11

mov edx, 11

mov edx, 11

mov eax, 7A'

call ds:GetModuleHandleW; call GetModuleHandleW to retrieve handle of kernel32.dll

cmov eax, 7A'

call ds:GetProcAddress; retrieve VirtualAlloc addr

cmov eax, edx

call ds:GetProcAddress; retrieve GetModuleFileNameA

cmp eax, 0Et call ds:GetProcAddress; retrieve CreateFileA addr

call ds:GetProcAddress; retrieve ReadFile addr

call ds:GetProcAddress; retrieve ReadFile addr

call [esp+1FCh+GetModuleFileNameA]; call GetModuleFileNameA to retrieve full path of module that load malware dll

call f_strstr; Returns a pointer to the first occurrence of a search string in a string.

call eax; call CreateFileA for open file but not retrieve file handle

call eax; call CreateFileA for repaing log, dat content to allocated buffer

call eax; call CreateFileA to retrieve handle to log, dat file

call ds:GetFileSize; call GetFileSize to retrieve size of log.dat

call eax; call VirtualAlloc to allocate buffer with buf's size equal size of log.dat

call eax; call StratA; call IstratA to build full path to log.dat
```

Decode the contents of **log.dat** into shellcode so that this shellcode is then executed by the call from the **LogInit_0**function.

3. Shellcode analysis

Based on the information analyzed above, we know that the **log.dll** file will read the content from the **log.dat** file and decrypt it into shellcode for further execution. Relying on this indicator, we continue to hunt **log.dat** file on VirusTotal which restrict the scope of submission source from Vietnam.

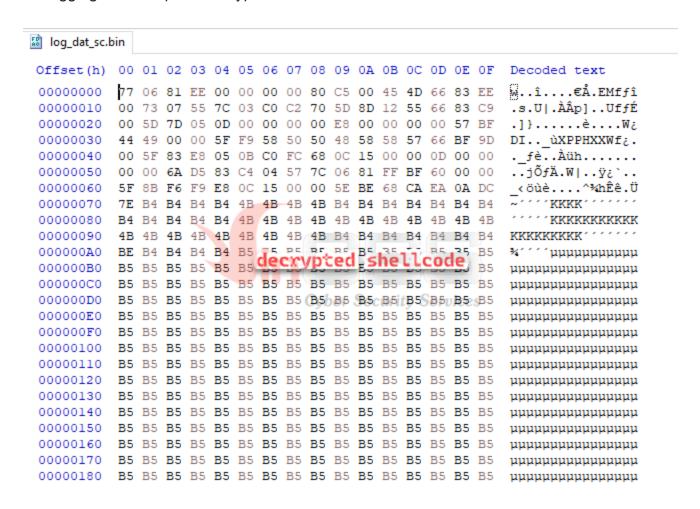
The results are following:



With the above results, at the time of analysis, we selected the **log.dat** file (2de77804e2bd9b843a826f194389c2605cfc17fd2fafde1b8eb2f819fc6c0c84) was submitted to VirusTotal on **2022-04-20 12:33:19 UTC** (5 days before the above **log.dll** file).



Debugging and dump the decrypted shellcode look like this:

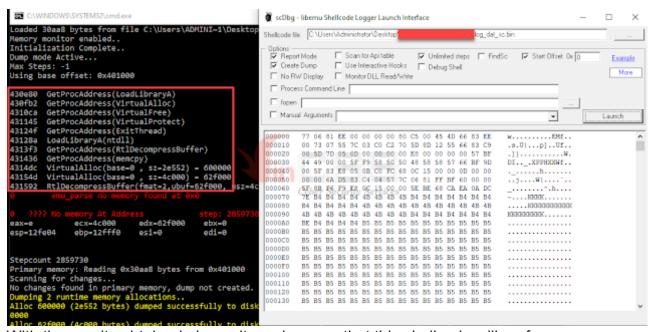


I use two tools, <u>FLOSS</u> and <u>scdbg</u> to get an overview of this shellcode. The results can be seen in the screenshots below:

```
FLOSS static Unicode strings

FLOSS decoded 2 strings
(EAA
&EAA

FLOSS extracted 8 stackstrings
VirtualProtect
VirtualAlloc
ExitThread
memcpy
ntdll
LoadLibraryA
VirtualFree
RtlDecompressBuffer
```



With the results obtained above, it can be seen that this shellcode will perform memory allocation and then call the **RtIDecompressBuffer** function to decompress the data with the compression format is **COMPRESSION_FORMAT_LZNT1**.

By using IDA to analyze this shellcode, its main task is to decompress a DII into memory and call the exported function of this DII to execute. The function that does this task is named **f_load_dII_from_memory**:

```
proc near ; CODE XREF: sub_403575+18†p
                                                                                                        2 int _usercall sub_431AE4@<eax>(int a1@<eax>)
text:00431AE4
                                          push
                                                             h ; shellcode size
text:00431AE9
                                          push
                                                     eax ; ptr_call_addr
                                                                                                                                 [esp-10h] [ebp-10h]
                                                                                                                            [esp-Ch] [ebp-Ch]
[esp-8h] [ebp-8h]
text:00431AEA
text:00431AEE
                                          stc
text:00431AEF
                                                                                                                            [esp-4h] [ebp-4h]
                                          stc
text:00431AF0
                                          test
text:00431AF2
                                          call
                                                                                                            return f_load_dll_from_memory(a1, 0×30AA8, v2, v3, v4, v5);
                                                                                                 10}
text:00431AF7
                                          retn
text:00431AF7 sub_431AE4
                                                                                 21 calls, 0 strings
text:00431AF7
                                                                                   alls:
                                                                                                             call [ebp+GetProcAddress]
call [ebp+GetProcAddress]
call [ebp+GetProcAddress]
call [ebp+GetProcAddress]
                                                                                                              call [ebp+0
                                                                                                              call [ebp+LoadLibraryA]
                                                                                                             call [ebp+GetProcAddress]
call [ebp+GetProcAddress]
call [ebp+VirtualAlloc]
                                                                                                             call [ebp+VirtualAlloc]
call [ebp+VirtualAlloc]
call [ebp+NtlbecompressBu
call [ebp+VirtualAlloc]
call [ebp+LoadLibraryA]
call [ebp+GetProcAddress]
call [ebp+GetProcAddress]
                                                                                                              call [ebp+VirtualProtect]
                                                                                                             call ecx ; call to DllEntryPoint
call [ebp+exported_func] ; call to PlugX exported function
call [ebp+virtualFree]
```

The code in this function will first get the base address of **kernel32.dll** based on the precalculated hash value is **0x6A4ABC5B**. This hash value has also been mentioned by us <u>in this analysis</u>.

```
kernel32\_base\_addr = 0;
GetProcAddress = 0;
pLdr = NtCurrentPeb()→Ldr;
 or ( ldr_entry = pLdr→InMemoryOrderModuleList.Flink; ldr_entry; ldr_entry = ADJ(ldr_entry)→InMemoryOrderLinks.Flink )
  wszDllName = ADJ(ldr_entry)→BaseDllName.Buffer;
  dll_name_length = ADJ(ldr_entry)→BaseDllName.Length;
  calced_hash = 0;
    calced_hash = __ROR4__(calced_hash, 13);
    if ( *wszDllName < '
      calced_hash += *wszDllName;
                                                                // calced_hash + letter
      calced_hash = calced_hash + *wszDllName - 0×20;
                                                                // calced_hash + upper_letter
    wszDllName = (wszDllName + 1);
    --dll_name_length;
   vhile ( dll_name_length );
  if ( calced_hash = 0 \times 6A4ABC5B )
                                                           python .\brute_force_Dll_name.py
Found dll kernel32.dll of 0x6a4abc5b
Found dll ntdll.dll of 0x3cfa685d
    kernel32_base_addr = ADJ(ldr_entry)→DllBase;
   (!kernel32_base_addr)
```

Next it will retrieve the address of GetProcAddress:

```
or ( i = 0; i < export_dir_va→NumberOfNames; ++i )
 szAPIName = kernel32_base_addr + pFuncsNamesAddr[i];
 if ( *szAPIName =
  && szAPIName[1] =
  && szAPIName[2] =
  && szAPIName[3] =
  && szAPIName[4] =
  && szAPIName[5] =
   && szAPIName[6] =
  && szAPIName[7] =
  && szAPIName[8] =
   && szAPIName[9] =
   GetProcAddress = (kernel32_base_addr
                  + *(kernel32_base_addr
                    + 4 * *(kernel32_base_addr + 2 * i + export_dir_va→AddressOfNameOrdinals)
                    + export_dir_va → AddressOfFunctions));
  (!GetProcAddress)
```

By using the <u>stackstring</u> technique, the shellcode constructs the names of the APIs and gets the addresses of the following API functions:

```
strcpy(szLoadLibraryA, "LoadLibraryA");
120
121
122
123
         LoadLibraryA = GetProcAddress(kernel32_base_addr, szLoadLibraryA);
         if ( !LoadLibraryA )
        return 3;
// "VirtualAlloc" → (size: 13)
strcpy(szVirtualAlloc, "VirtualAlloc");
   124
         VirtualAlloc = GetProcAddress(kernel32_base_addr, szVirtualAlloc);
         if ( !VirtualAlloc )
   130
         // "VirtualFree" → (size: 12)
strcpy(szVirtualFree, "VirtualFree");
         VirtualFree = GetProcAddress(kernel32_base_addr, szVirtualFree);
         if (!VirtualFree)
         // "VirtualProtect" → (size: 15)
strcpy(szVirtualProtect, "VirtualProtect");
138139
         VirtualProtect = GetProcAddress(kernel32_base_addr, szVirtualProtect);
         if ( !VirtualProtect )
           return 6;
         // "ExitThread" → (size: 11)
strcpy(szExitThread, "ExitThread");
• 142
                                                                                    LoadLibraryA
                                                                                    VirtualAlloc
         if ( !GetProcAddress(kernel32_base_addr, szExitThread) )
                                                                                    VirtualFree
           return 6;
                                                                                    VirtualProtect
         strcpy(szntdll, "ntdll");
ntdll_handle = LoadLibraryA(szntdll);
                                                                                    ExitThread
                                                                                    RtlDecompressBuffer
         if ( !ntdll_handle )
                                                                                    memcpy
         strcpy(szRtlDecompressBuffer, "RtlDecompressBuffer");
```

Next, the shellcode performs a memory allocation (**compressed_buf**) of size **0x2E552**, then reads data from offset **0x1592** (on disk) and executes an xor loop with a key is **0x72** to fill data into the **compressed_buf**. In fact, the size of **compressed_buf** is **0x2E542**, but its first 16 bytes are used to store information about **signature**, **uncompressed_size**, compressed_size, so **0x10** is added.

Shellcode continues to allocate memory (uncompressed_buf) of size 0x4C000 and calls the RtIDecompressBuffer function to decompress the data at the compressed_buf into uncompressed_buf with the compression format is COMPRESSION_FORMAT_LZNT1.

```
ptr_enc_compressed_dll_addr = 0×1592 (offset on disk)
signature = 0×C7EA9B1C
signature = *ptr_enc_compressed_dll_addr;
xor_key = signature - 0×7979A9AA;
// dd 0B598E96Eh

// dd 0C7EA9B1Ch → signature

// dd 0004C000h → uncompressed_size

// dd 2E542h → compressed_size;
for (j = 0; j < 0 \times 10; ++j)
  config_info_buf[j] = xor_key ^ ptr_enc_compressed_dll_addr[j];// xor_key = 0×72
if ( signature # computed_signature )
     urn 0×A;
dwSize = computed_compressed_size + 0×10;
                                                                  // dwSize = 0 \times 2E552
compressed_buf = VirtualAlloc(0, computed_compressed_size + 0×10, MEM
if ( !compressed_buf )
  return 0×B;
xor_key = signature - 0×7979A9AA;
For (k = 0; k < dwSize; ++k)
 *(&compressed_buf->decoded_buffer + k) = xor_key ^ ptr_enc_compressed_dll_addr[k];
uncompressed_buf = VirtualAlloc(0, uncompressed_buf_size, ME
if ( !uncompressed_buf )
  return 0×C;
final_uncompressed_size = 0;
// decompress dll payload to m
if ( RtlDecompressBuffer(
       uncompressed_buf
       uncompressed_buf_size,
       &compressed_buf→compressed_buf,
       compressed_buf → compressed_size,
                                                                   // 0×2E542
       &final_uncompressed_size) )
  return 0×D;
   ( uncompressed_buf_size ≠ final_uncompressed_size )
```

Based on the above analysis results, it is easy to get the extracted DII file (however, the file header information was destroyed):

decompression	essed_c	dll_40	2000.	dum	р												
Offset (h) 00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	Decoded text
00000000	6C	41	76	62	42	48	6A	44	4C	75	4D	42	54	6B	57	57	1AvbBHjDLuMBTkWW
00000010	45	78	5A	45	4F	6F	54	65	79	70	75	44	63	4B	4E	45	ExZEOoTeypuDcKNE
00000020	74	6C	73	50	61	48	48	78	69	5A	7A	4A	6E	4E	6E	74	tlsPaHHxiZzJnNnt
00000030	69	49	46	4C	42	43	4F	59	50	58	54	00	E0	00	00	00	iIFLBCOYPXT.à
00000040	78	43	52	55	6A	44	62	52	4E	4C	58	4A	76	73	47	79	xCRUjDbRNLXJvsGy
00000050	75	4F	77	76	55	59	55	76	76	46	58	5A	77	7A	42	55	uOwvUYUvvFXZwzBU
00000060	70	6F	4B	48	4D	75	50	46	45	45	67	45	73	67	71	61	poKHMuPFEEgEsgqa
00000070			75	4C	6E	6C	53	52	74	69	51	72	7A	63		49	ViuLnlSRtiQrzcLI
00000080		7A	61	55	6E	5A	6A	78	79	45	51	62	6D	76	42	69	izaUnZjxyEQbmvBi
00000090			67	72	75	55	64	46	4E	6C	78	78	50	6F		64	SOgruUdFNlxxPoPd
000000A		72	75	68	61	69	67	6F	61	58	52	71	4E	59		6C	uruhaigoaXRqNYcl
000000B0		4E	58	72	4C	44	42	69	48	49	65	67	56	43		48	uNXrLDBiHIegVCuH
00000000		73	77	48	68	53	6B	45	72	4B	77	68	55	6C	52	78	wswHhSkErKwhUlRx
000000D0			6B	46	42	64	59	79	4C	6E	79	72	50	52		54	LDkFBdYyLnyrPRqT
000000E0		6C		00	4C	01	03	00	30	83	1E	53	00	00	00	00	S1LOf.S
000000F0		00	00	00	E0	0.0	02	21	0B	01	0C	00	00	00	00	00	à!
00000100		3C	00	00	00	00		00	B0	81		00	00	10	00	00	.<°
00000110		10	00	0.0		00	00	10		10		00		02		00	
00000120			01	00	00	d	9C(omp	re	SS	ed	D)	SP)	vice	00	00	à 0
00000130		00	10	00	00	10	00	00	00	00	10			10	40 00	00	.à
00000140		00	00	00	10	00	00	00	60	8F	04	00	45	00		00	
00000150		91	04	00	78	00	00	00	00	00	00	00	00	00	00	00	0'x
00000170		00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000170			04	00	0C	33	00	00	00	00	00	00	00	00	00	00	3
00000190		00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
000001A0		00	00	00	00	00	00	00	50	7A	00	00	40	00	00	00	
000001B0		00	00	00	00	00	00	00	00	90	04	00	30	01	00	00	
000001C0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
000001D0	00	00	00	00	00	00	00	00	2E	74	65	78	74	00	00	00	text
000001E0	A5	7 F	04	00	00	10	00	00	00	80	04	00	00	04	00	00	¥€
000001F0	00	00	00	00	00	00	00	00	00	00	00	00	60	00	00	60	
00000200	2E	69	64	61	74	61	00	00	D2	07	00	00	00	90	04	00	.idataÒ
00000210	00	08	00	00	00	84	04	00	00	00	00	00	00	00	00	00	, ,
00000220	00	00	00	00	40	00	00	40	2E	72	65	6C	6F	63	00	00	@@.reloc
00000230	0C	33	00	00	00	A0	04	00	00	34	00	00	00	8C	04	00	.34Œ

Fix the header information and check with <u>PE-bear</u>, this DII has the original name is **RFPmzNfQQFPXX** and only exports one function named **Main**:

255			,				
Offset	Name	V	alue	Meaning	9		
48360	Characterist	tics 0					
48364	TimeDateStar	пр 6	12C95CD	Monday	, 30.08.2021	08:24:45 UTC	
48368	MajorVersion	n 0					
4836A	MinorVersion	ח 0					
4836C	Name	4	8F92	RFPmzNf	FQQFPXX		
48370	Base	1					
48374	NumberOfFund	ctions 1					
48378	NumberOfName	es 1					
4837C	AddressOfFur	nctions 4	8F88				
48380	AddressOfNam	nes 4	8F8Cybe	r Secur	ity Services		
48384	AddressOfNa	neOrdinals 4	8F90				
Exported	Functions [1 entry]					
Offset	Ordinal	Function	RVA Na	ne RVA	Name	Forwarder	
48388	1	8190	488	FA0	Main		

Back to the shellcode, after decompressing the DII into memory, it will perform the task of a loader to map this DII into a new memory region. Then, call to the exported function (here is the **Main** function) to perform the the main task of malware:

```
plugx_decrypted_dll = plugx_mapped_dll;
   0070000 00 00 00 00 29 00 6C 02 A8 0A 03 00 92 15 6C 02 ....).l."....'.l.
0070010 52 E5 02 00 69 00 6C 02 0C 15 00 00 00 00 00 Rå..i.l......
plugx_mapped_dll→signature = 0;
plugx_decrypted_dll→ptr_shellcode_base = ptr_call_addr; // 00402029 E8 00 00 00 00
plugx_decrypted_dll -> shellcode_size = end_sc_offset;
plugx_decrypted_dll→ptr_encrypted_PlugX = ptr_enc_compressed_dll_addr;// 00403592 1C 9B ....
plugx_decrypted_dll→encrypted_PlugX_size = compressed_dll_size;// 0×2E
                                                               // 0×0402069 (offset 0×69 on disk)
// 0×0150C
plugx_decrypted_dll→config = config;
plugx_decrypted_dll > config_size = config_size;
plugx_decrypted_dll -> ptr_PlugX_entry_point = plugx_mapped_dll + payload_nt_headers -> OptionalHeader.AddressOfEntryPoint;
VirtualProtect(lpAddress, payload_raw_size,
                                                                         &fl0ldProtect);
if ( !(plugx_decrypted_dll→ptr_PlugX_entry_point)(plugx_mapped_dll, 1, 0) )
   return 0×15;
if ( ExportProc )
  ExportProc();
  ( !VirtualFree(compressed_buf, 0, M
return 0×16;
if ( VirtualFree(uncompressed_buf, 0, MEM_RELEASE) )
  return 0;
  eturn 0×17;
```

Note: At the time of analyzing this shellcode, we have not yet confirmed it is a variant of the PlugX malware, but only raised doubts about the relationship. It was only when we analyzed the above extracted DII, then we confirmed for sure that this was a variant of PlugX and renamed the fields in the struct for understandable reasons as screenshot above.

4. Analyze the extracted DII

We will not go into detailed analysis of this DII, but only provide the necessary information to prove that this is a PlugX variant as well as the process of decrypting the configuration information that the malware will be used.

4.1. How PlugX calls an API function

In this variant, information about API functions is stored in **xmmword**, then loaded into the **xmm0** (128-bit) register, the missing part of the function name will be loaded through the stack. The malicious code gets the handle of the DII corresponding to these API functions, then uses **GetProcAddress** function to retrieve the address of the specified API function to use later:

```
text:10027A90 000
                           push
                                   ebp
text:10027A91 004
                           mov
                                   ebp, esp
text:10027A93 004
                           sub
                                   esp, 84h
text:10027A99
                                   xmm0, xmmword_100078A0
                           movdqa
                                   eax, GetCurrent
text:10027AA1 088
                           mov
                                                   xmmword_100078A0 xmmword 'secorPtnerruCteG
text:10027AA6 088
                                   ebx
                           push
                                                                                              DATA XRE
text:10027AA7 08C
                           push
                                                                                               f_plugx
text:10027AA8 090
                           xor
text:10027AAA 090
                                   [ebp+lpName], ecx
                           mov
                                   [ebp+token_handle],
text:10027AAD 090
                           mov
                                   [ebp+var_60], 73h
text:10027AB0 090
                           moν
text:10027AB6 090
                                   edi
                           push
text:10027AB7 094
                                   edi, ds:GetProcAddress
                          mov
text:10027ABD 094
                                   xmmword ptr [ebp+ProcName], xmm0
                           movdqu
text:10027AC2 094
                                   eax, eax
text:10027AC4 094
                           jnz
                                   short loc_10027AD7
text:10027AC4
                                   eax, [ebp+ProcName]
text:10027AC6 094
                           lea
text:10027AC9 094
                                                              lpProcName
                           push
                                   f_retrieve_kernel32_handle
text:10027ACA 098
                           call
text:10027ACA
                                                            ; hModule
text:10027ACF
                           push
                                   edi ; GetProcAddress
text:10027AD0 09C
                           call
text:10027AD0
                                   GetCurrentProcess_0, eax
text:10027AD2 094
                           mov
text:10027AD2
text:10027AD7
                  loc_10027AD7:
                                                              CODE XREF: f_check_and_enable_privilege
text:10027AD7
                                   eax ; GetCurrentProcess_6
text:10027AD7 094
                           call
```

4.2. Create main thread to execute

The malware adjusts the **SeDebugPrivilege** and **SeTcbPrivilege** tokens of its own process in order to gain full access to system processes. Then it creates its main thread, which is named "**bootProc**":

```
f_create_unnamed_event(0)→dll_base = <mark>dll_base</mark>;
f_create_unnamed_event(0) → dll_base = dll_base;
f_create_unnamed_event(0) → dll_base = dll_base;
*wszSeDebugPrivilege =
*&wszSeDebugPrivilege[2] =
*&wszSeDebugPrivilege[4] =
*&wszSeDebugPrivilege[6] =
*&wszSeDebugPrivilege[8] =
*&wszSeDebugPrivilege[0×A] =
*&wszSeDebugPrivilege[0×C] =
*&wszSeDebugPrivilege[0×E] =
wszSeDebugPrivilege[0×10] = 0;
*wszSeTcbPrivilege =
*&wszSeTcbPrivilege[2] =
*&wszSeTcbPrivilege[4] =
*&wszSeTcbPrivilege[6] =
*&wszSeTcbPrivilege[8] =
*&wszSeTcbPrivilege[0×A] =
*&wszSeTcbPrivilege[0×C] =
V6 = 0:
                                                             // SeDebugPrivilege
// SeTcbPrivilege
f_check_and_enable_privilege(wszSeDebugPrivilege);
f_check_and_enable_privilege(wszSeTcbPrivilege);
strcpy(szbootProc, "bootProc");
critical_section = sub_10007E50(0);
return f_spawn_thread(critical_section, &p_thread_handle, szbootProc, f_main_thread_func
```

4.3. Communicating with C2

The malware can communicate with C2 via TCP, HTTP or UDP protocols:

```
*szProto_Host_Proxy_format_str = _mm_load_si128(&xmmword_10007120);
                                                        strcpy(v15, "%s:%s]\r\n");
port_num_hi = HIWORD(src > f_retrieve_ip_address);
                                                         port_num_lo = LOWORD(src > f_retrieve_ip_address);
                                                         v8 = a2[1];
strcpy(szTCP_proto, "TCP");
strcpy(szHTTP_proto, "HTTP"
sz_protocol_info = L"*";
                                                         v13 = _mm_load_si128(&xmmword_10007240);
strcpy(szUDP_proto, "UC
                                                         v14 = _mm_load_si128(&xmmword_19997189);
strcpy(szICMP_proto, "ICMP
switch ( choose_proto_flag
                                                         // Protocol:[%4s], Host: [%s:%d], Proxy: [%d:%s:%d:%s:%s]\r\n
                                                           szProto_Host_Proxy_full_str,
     se 2:
                                                           szProto_Host_Proxy_format_str,
    sz_protocol_info = szTCP_proto;
                                                           sz_protocol_info
                                                           a2 + 2,
     se 3:
                                                           v8
    sz_protocol_info = szHTTP_proto;
                                                           port_num_lo
                                                           &src→field_4,
                                                           port_num_hi.
    sz_protocol_info = szUDP_proto;
                                                           &src -> event_handle_1,
                                                           &src→field_84);
    se 5:
                                                         f_send_str_to_debugger(szProto_Host_Proxy_full_str);
    sz_protocol_info = szICMP_proto;
                                                             ch ( choose_proto_flag )
                                                             result = f_connect_c2_over_TCP(this, arg0, a2, src);
                                                              e 3:
                                                             result = f_connect_c2_over_HTTP(this, arg0, a2, src);
                                                                40
                                                             result = f_connect_c2_over_UDP(this, arg0, a2, src);
                                                             result = 0×32;
```

4.4. Implemented commands

The malware will receive commands from the attacker to execute the corresponding functions related to *Disk*, *Network*, *Process*, *Registry*, etc.

```
as_file_bef = f_aspping_file_and_retrun_buf();
if (asp_file_bef) = f_aspping_file_and_retrun_buf();
if (asp_file_bef) = f_aspping_file_and_retrun_buf();
if (asp_file_bef) = f_aspping_file_and_retrun_buf();
if (asp_file_bef) = f_aspping_file_bef) = f_aspping_file_b
```

The entire list of commands as shown in the table below that the attacker can execute through this malware sample:

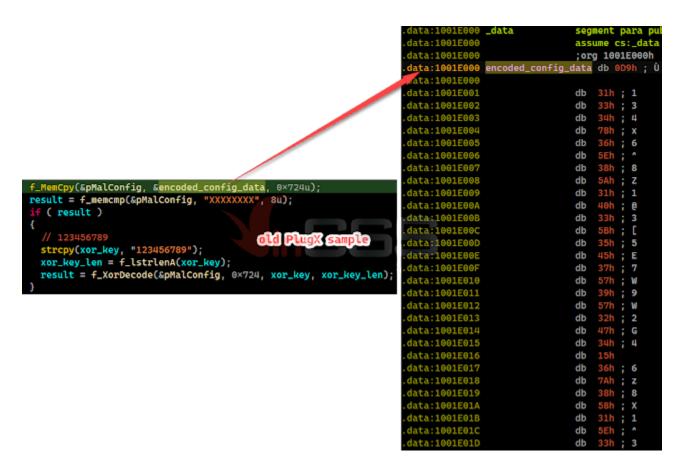
Command Group	Sub- command	Description
Disk	0x3000	Get information about the drives (type, free space)
	0x3001	Find file
	0x3002	Find file recursively
	0x3004	Read data from the specified file
	0x3007	Write data to the specified file
	0x300A	Create a new directory
	0x300C	Create a new process on hidden desktop
	0x300D	File action (file copy/rename/delete/move)
	0x300E	Expand environment-variable strings
Nethood	0xA000	Enumeration of network resources
Netstat	0xD000	Retrieve a table that contains a list of TCP endpoints
	0xD001	Retrieve a table that contains a list of UDP endpoints
	0xD002	Set the state of a TCP connection

Option	0x2000	Lock the workstation's display				
	0x2001	Force shut down the system				
	0x2002	Restart the system				
	0x2003	Shut down the system safety				
	0x2005	Display massage box				
PortMap	0xB000	Perform port mapping				
Process	0x5000	Retrieve processes info				
	0x5001	Retrieve modules info				
	0x5002	Terminate specified process				
RegEdit	0x9000	Enumerate registry				
	0x9001	Create registry				
	0x9002	Delete registry				
	0x9003	Copy registry				
	0x9004	Enumerates the values of the specified open registry key				
	0x9005	Sets the data and type of a specified value under a registry key				
	0x9006	Deletes a named value from the specified registry key				
	0x9007	Retrieves a registry value				
Service	0x6000	Retrieves the configuration parameters of the specified service				
	0x6001	Changes the configuration parameters of a service				
	0x6002	Starts a service				
	0x6003	Sends a control code to a service				
	0x6004	Delete service				
Shell	0x7002	Create pipe and execute command line				
SQL	0xC000	Get SQL data sources				
	0xC001	Lists SQL drivers				
	0xC002	Executes SQL statement				

Telnet	0x7100	Start telnet server
Screen	0x4000	simulate working over the RDP Protocol
	0x4100	Take screenshot
KeyLog	0xE000	Perform key logger function, log keystrokes to file "%allusersprofile%MSDN6.0USER.DAT"

4.5. Decrypt PlugX configuration

As analyzed above, the malware will connect to the C2 address via HTTP, TCP or UDP protocols depending on the specified configuration. So where is this config stored? With the old malware samples that we have analyzed $(\underline{1}, \underline{2}, \underline{3}, \underline{4})$, the PlugX configuration is usually stored in the .data section with the size of 0x724 (1828) bytes.



Going back to the sample we are analyzing, we see that before the step of checking the parameters passed when the malware executes, it will call the function that performs the task of decrypting the configuration:

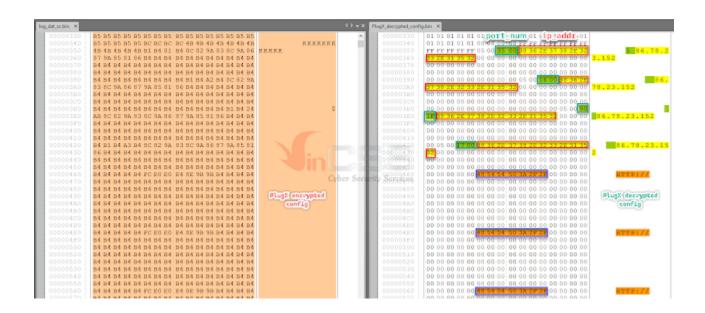
```
ptr_cmd_line = GetCommandLineW();
CommandLineToArgvW = ::CommandLineToArgvW;
strcpy(v46, "vW");
*v45 = _mm_load_si128(&xmmword_10007610);
if ( !::CommandLineToArgvW )
  shell32_handle = g_shell32_handle;
  strcpy(sz_shell32, "shell32");
  if ( !g_shell32_handle )
    shell32_handle = LoadLibraryA(sz_shell32);
    g_shell32_handle = shell32_handle;
  CommandLineToArgvW = GetProcAddress(shell32_handle, v45);
   ::CommandLineToArgvW = CommandLineToArgvW;
sz_arg_list = CommandLineToArgvW(ptr_cmd_line, &num_arguments);
sub_10007DC0(0);
f_decrypt_embedded_config_or_from_file_and_copy_to_mem();
if ( num_arguments = 1 )
  f_launch_process_or_create_service();
if ( num_arguments = 3 )
  lstrlenW = ::lstrlenW;
  arg_passed_1 = sz_arg_list[1];
  passed_arg1_info.buffer = 0;
  passed_arg1_info.buffer1 = 0;
```

Diving into this function, combined with additional debugging from shellcode, renaming the fields in the generated struct, we get the following information:

- PlugX's configuration is embedded in shellcode and starts at offset 0x69.
- The size of the configuration is **0x0150C** (**5388**) bytes.
- Decryption key is **0xB4**.

With all the complete information as above, it is possible to recover the configuration information easily:

IP	Port
86.78.23.152	53
86.78.23.152	22
86.78.23.152	8080
86.78.23.152	23



In addition to the list of C2 addresses above, there is additional information related to the directory created on the victim machine to contain malware files as well as the name of the service that can be created:

```
00 00 00 00 00 25 00 50 00 72 00 6F 00 67 00 72
                                                   00 61 00 6D 00 46 00 69 00 6C 00 65 00 73 00 25
wsz_bdreinit_exe[0] =
                                                   00 5c 00 42 00 69 00 74 00 44 00 65 00 66 00 65
                                                                                                  В
wsz bdreinit exe[1] =
                                                   00 6E 00 64 00 65 00 72 00 20 00 55 00 70 00 64
                                                                                                  d
                                                   wsz_bdreinit_exe[2] =
wsz_bdreinit_exe[3]
wsz_bdreinit_exe[0] =
                                                  00065-0068:00-64-00-65 00 72 00 20 00 43 00 72
wsz_bdreinit_exe[5] =
                                                  00 61 00 73 00 68 00 20 00 48 00 61 00 6E 00 64
                                                                                                  55
                                                                                                    h
LOWORD(wsz_bdreinit_exe[6]) = 0;
                                                  00 60 00 65 00 72 00 00 00 00 00 00 00 00 00 00
```

To make our life easier, I wrote a python script to automatically extract configuration information for this variant. The output after running the script is as follows:

```
$ python plugx_extract_config.py plugx_decrypted_config.bin

[+] Config file: plugx_decrypted_config.bin
[+] Config size: 5388 bytes
[+] Folder name: %ProgramFiles%\BitDefender Update
[+] Service name: BitDefender Crash Handler
[+] Proto info: HTTP://
[+] C2 servers:
      86.78.23.152:53
      86.78.23.152:22
      86.78.23.152:23
[+] Campaign ID: 1234
```

5. Conclusion

CrowdStrike researchers first published information on Mustang Panda in June 2018, after approximately one year of observing malicious activities that shared unique Tactics, Techniques, and Procedures (TTPs). However, according to research and collect from many different cybersecurity companies, this group of APTs has existed for more than a decade with different variants found around the world. Mustang Panda, believed to be a APT group based in China, is evaluated as one of the highly detrimental APT groups, applying sophisticated techniques to infect malware, aiming to gain as much long-term access as possible to conduct espionage and information theft.

In this blog we have analyzed the different steps the infamous PlugX RAT follows to start execution and avoid detection. Thereby, it can be seen that this APT group is still active and constantly looking for ways to improve their techniques. VinCSS will continue to search for additional samples and variants that may be associated with this PlugX variant that we analyzed in this article.

6. References

7. Indicators of Compromise

log.dll - db0c90da56ad338fa48c720d001f8ed240d545b032b2c2135b87eb9a56b07721

log.dll - 84893f36dac3bba6bf09ea04da5d7b9608b892f76a7c25143deebe50ecbbdc5d

log.dll - 3171285c4a846368937968bf53bc48ae5c980fe32b0de10cf0226b9122576f4e

log.dll – da28eb4f4a66c2561ce1b9e827cb7c0e4b10afe0ee3efd82e3cc2110178c9b7a

log.dat – 2de77804e2bd9b843a826f194389c2605cfc17fd2fafde1b8eb2f819fc6c0c84 Decrypted config:

[+] Folder name: %ProgramFiles%BitDefender Update

[+] Service name: BitDefender Crash Handler

[+] Proto info: HTTP://

[+] C2 servers:

86.78.23.152:53

86.78.23.152:22

86.78.23.152:8080

86.78.23.152:23

[+] Campaign ID: 1234

log.dat - 0e9e270244371a51fbb0991ee246ef34775787132822d85da0c99f10b17539c0 Decrypted config:

[+] Folder name: %ProgramFiles%BitDefender Update

[+] Service name: BitDefender Crash Handler

[+] Proto info: HTTP://

[+] C2 servers:

86.79.75.55:80

86.79.75.55:53

86.79.75.46:80

86.79.75.46:53

[+] Campaign ID: 1234

log.dat – 3268dc1cd5c629209df16b120e22f601a7642a85628b82c4715fe2b9fbc19eb0 Decrypted config:

[+] Folder name: %ProgramFiles%Common FilesARO 2012

[+] Service name: BitDefender Crash Handler

[+] Proto info: HTTP://

[+] C2 servers:

86.78.23.152:23

86.78.23.152:22

86.78.23.152:8080

86.78.23.152:53

[+] Campaign ID: 1234

log.dat – 02a9b3beaa34a75a4e2788e0f7038aaf2b9c633a6bdbfe771882b4b7330fa0c5 (THOR PlugX)
Decrypted config:

[+] Folder name: %ProgramFiles%BitDefender Handler

[+] Service name: BitDefender Update Handler

[+] Proto info: HTTP://

[+] C2 servers:

www.locvnpt.com:443

www.locvnpt.com:8080

www.locvnpt.com:80

www.locvnpt.com:53

[+] Campaign ID: 1234

Click here for Vietnamese version.

Dang Dinh Phuong – Threat Hunter

Tran Trung Kien (aka m4n0w4r) – Malware Analysis Expert

R&D Center – VinCSS (a member of Vingroup)

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25/04/2022

[RE026] A Deep Dive into Zloader - the Silent Night

Zloader, a notorious banking trojan also known as Terdot or Zbot. This trojan was first discovered in 2016, and over time its distribution number has also continuously increased. The Zloader's code is said to be built on the leaked source code of the famous ZeuS malware. In 2011, when source code of ZeuS was made public and since then, it has been used in various malicious code samples.



09/11/2021

[EX008] The exploit chain allows to take control of Zalo user accounts

While using the Zalo application, one of the popular chat applications in Vietnam today (According to statistics from Wikipedia, since May 2018, Zalo has reached 100 million users), the Threat Hunting team from VinCSS LLC discovered some security vulnerabilities that allow the attacker to form an exploit chain to take control of Zalo accounts.



27/10/2021

[RE025] TrickBot ... many tricks

1. Introduction First discovered in 2016, until now TrickBot (aka TrickLoader or Trickster) has become one of the most popular and dangerous malware in today's threat landscape. The gangs behind TrickBot are constantly evolving to add new features and tricks. Trickbot is multi-modular malware, with a main payload will be responsible for loading other plugins [...]



10/08/2021

[EX007] How playing CS: GO helped you bypass security products

Many of us love to play games, and as offensive security engineers, we also want to learn about how game studios are dealing with cheaters. We have observed that cheaters have used vulnerable graphic drivers to bypass anti-cheat mechanisms from several gaming cheating forums. In some cases, the cheaters tried to install vulnerable driver versions onto their computers, then exploited the vulnerability to read and write the game process's memory with the kernel privileges.



03/07/2021

[RE023] Quick analysis and removal tool of a series of new malware variant of Panda group that has recently targeted to Vietnam VGCA

Through continuous cyber security monitoring and hunting malware samples that were used in the attack on Vietnam Government Certification Authority, and they also have attacked a large corporation in Vietnam since 2019, we have discovered a series of new variants of the malware related to this group.