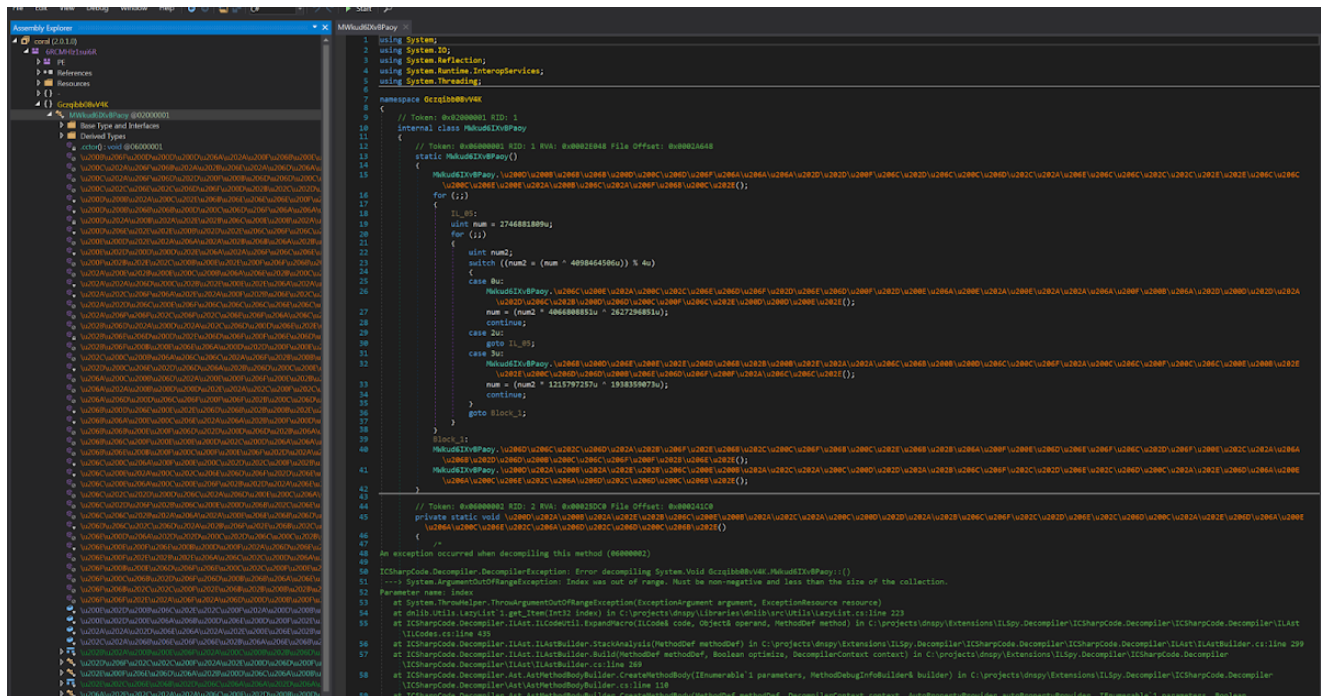


Recam Redux - DeConfusing ConfuserEx

blog.talosintelligence.com/2017/12/recam-redux-deconfusing-confuserex.html



This post is authored by [Holger Unterbrink](#) and [Christopher Marczewski](#)

Overview

This report shows how to deobfuscate a custom .NET ConfuserEx protected malware. We identified this recent malware campaign in our Advanced Malware Protection (AMP) telemetry. Initial infection is via a malicious Word document, the malware ultimately executes in memory an embedded payload from the Recam family. Recam is an information stealer. Although the malware has been around for the past few years, there's a reason you won't see a significant amount of documentation concerning its internals. The authors have gone the extra mile to delay analysis of the sample, including multiple layers of data encryption, string obfuscation, piecewise nulling, and data buffer constructors. It also relies on its own C2 binary protocol which is heavily encrypted along with any relevant data before transmission.

Technical Details

The Dropper

The word document (see above) uses common malware techniques, such as embedded VB code, to drop a .NET executable. We will not discuss these techniques further, but concentrate on the deobfuscation of the .NET malware dropper. The dropper is heavily

obfuscated with a custom version of ConfuserEx, a free .NET Framework protector. On opening the binary in a .NET decompiler like dnSpy it is initially unreadable (see Fig. 1).

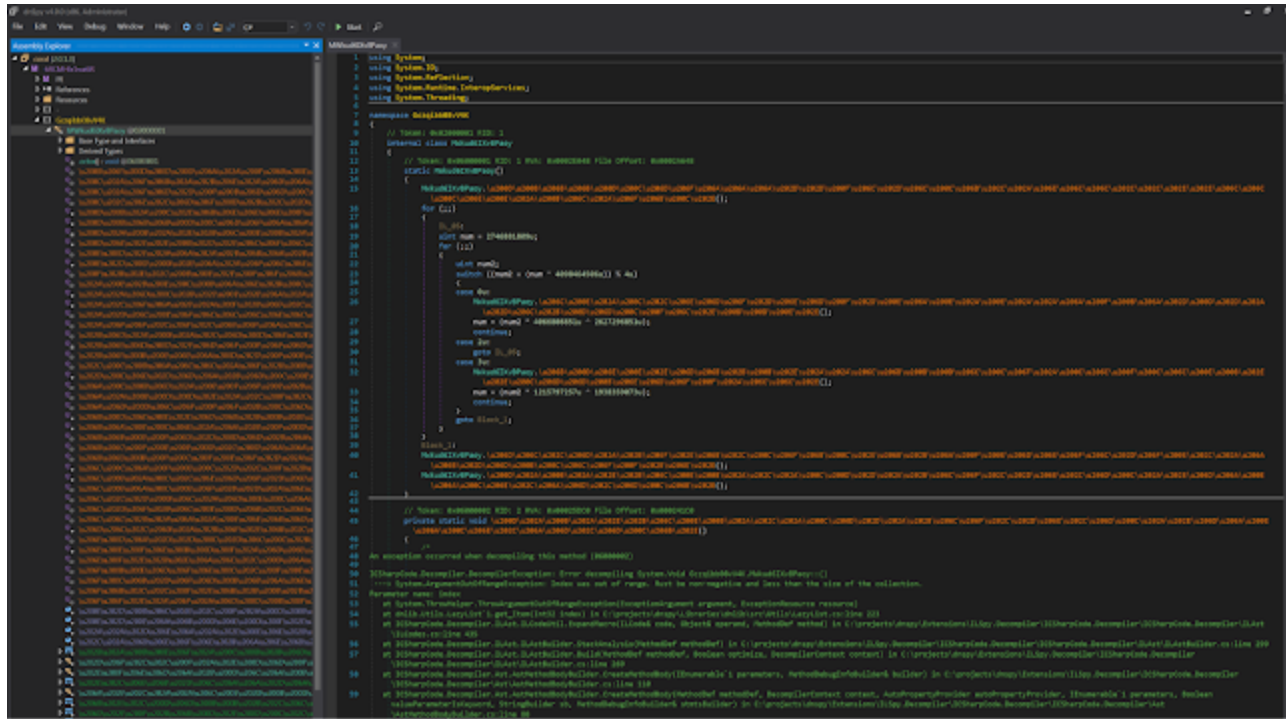


Fig. 1

There are a number of free deobfuscators available for ConfuserEx protected binaries; however, none of them are effective for this malware. Only some parts are able to be deobfuscated using these automated tools, leaving important sections of the binary unchanged, and breaking execution. This means we have no choice but to do it the hard way and deobfuscate it manually. There is documentation for manually unpacking ConfuserEx, but unfortunately, we hit bad luck again. The available documentation doesn't work with this version.

To get started, we first load the binary into dnSpy. We go to the <Module>. ctor and set a breakpoint on the last method (Fig 2). Now we can run the sample in our debugger and see that it has unpacked the first DLL ("ykMTM..." see Fig.2)

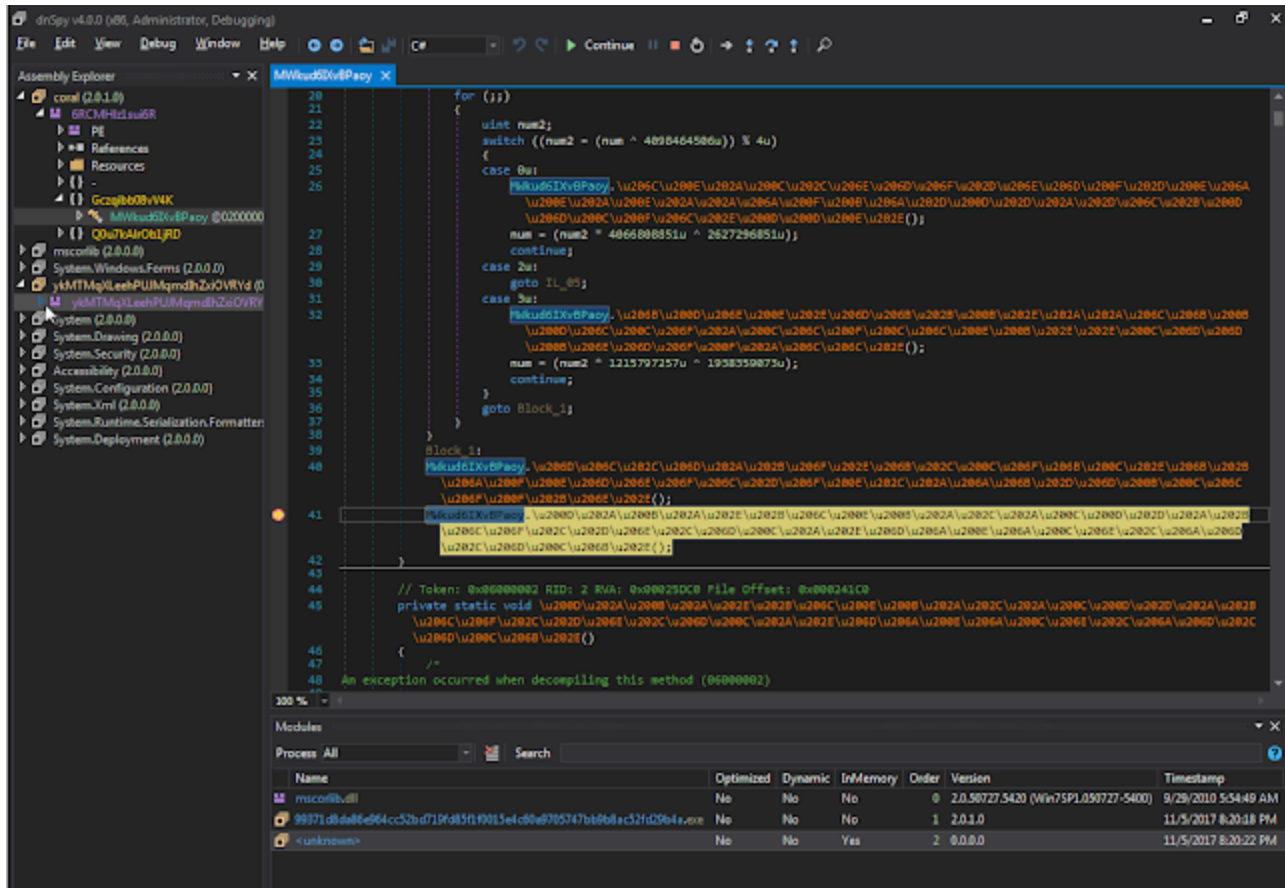


Fig. 2

We single step into the method where we hit the breakpoint and see in Fig. 3 that it has unpacked the next stage (coral).

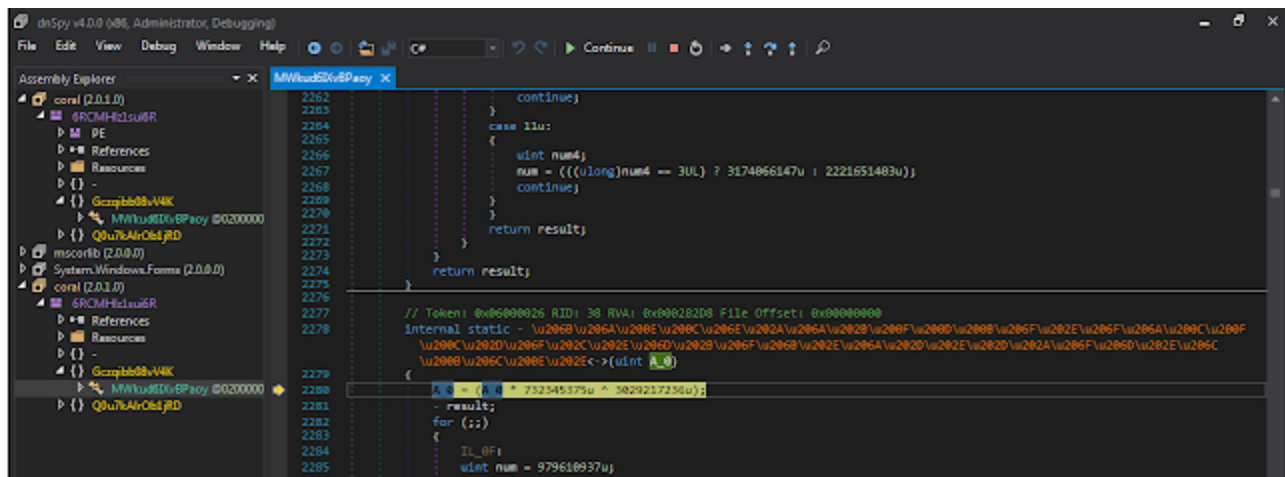


Fig. 3

We analysed this stage and found that we can set another breakpoint in the qMayiwZxj class

on line 113 (see Fig. 4)

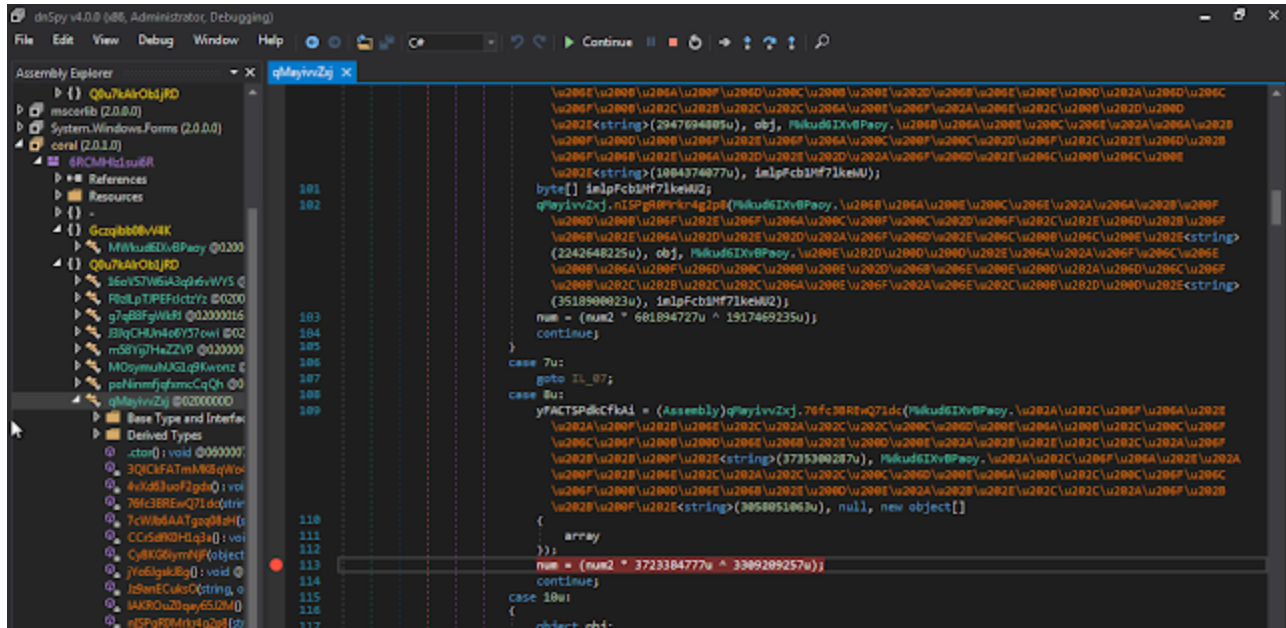


Fig. 4

This unpacks the next stage and we see the new unpacked stub.exe assembly (Fig. 5).

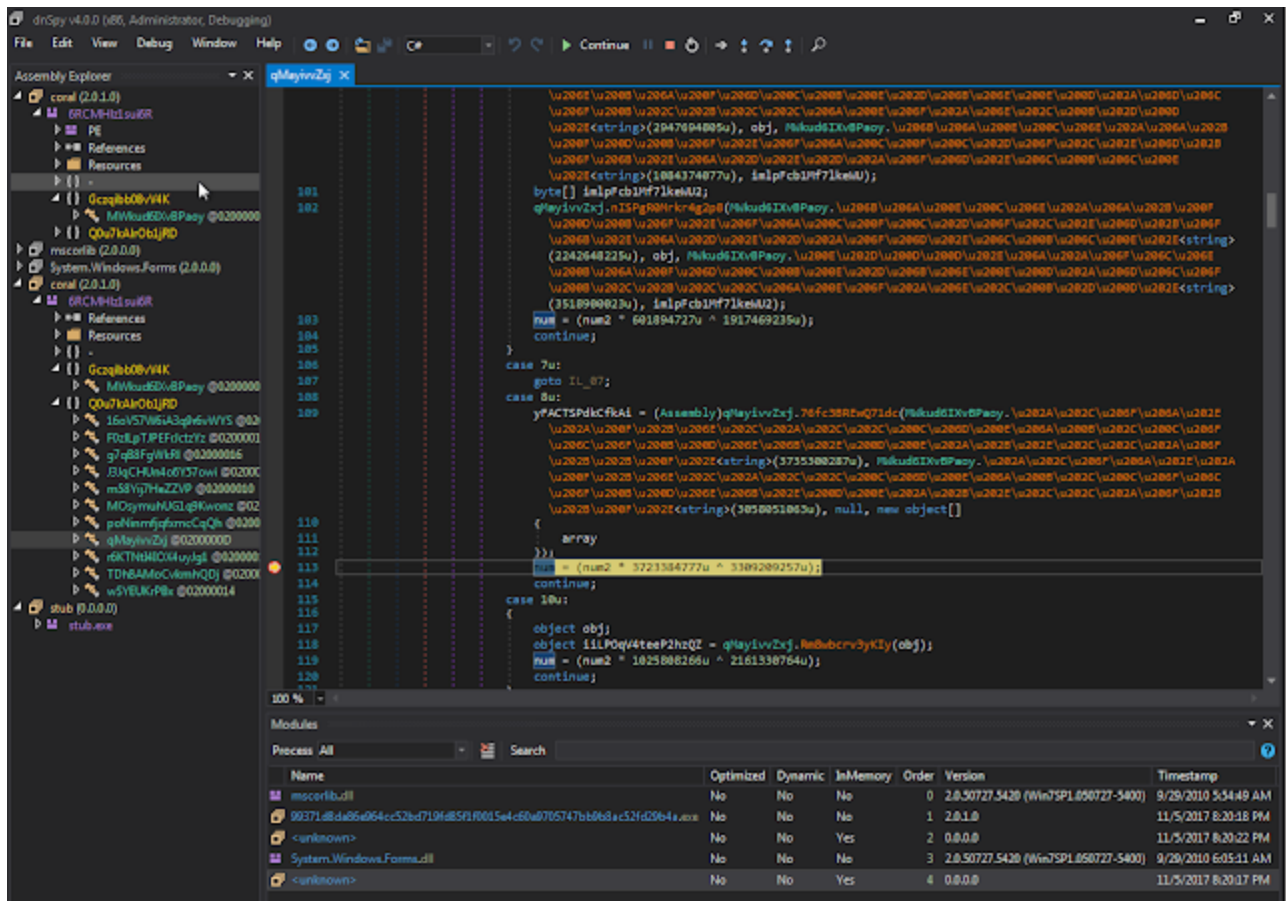


Fig. 5

If you have looked into other ConfuserEx obfuscated binaries, this looks familiar. Indeed, if you have a closer look, there is a well known friend, the `gchandle.free()` call on line 10082. This is our next breakpoint candidate. This call used to be the end of an unpacking stage in previous versions.

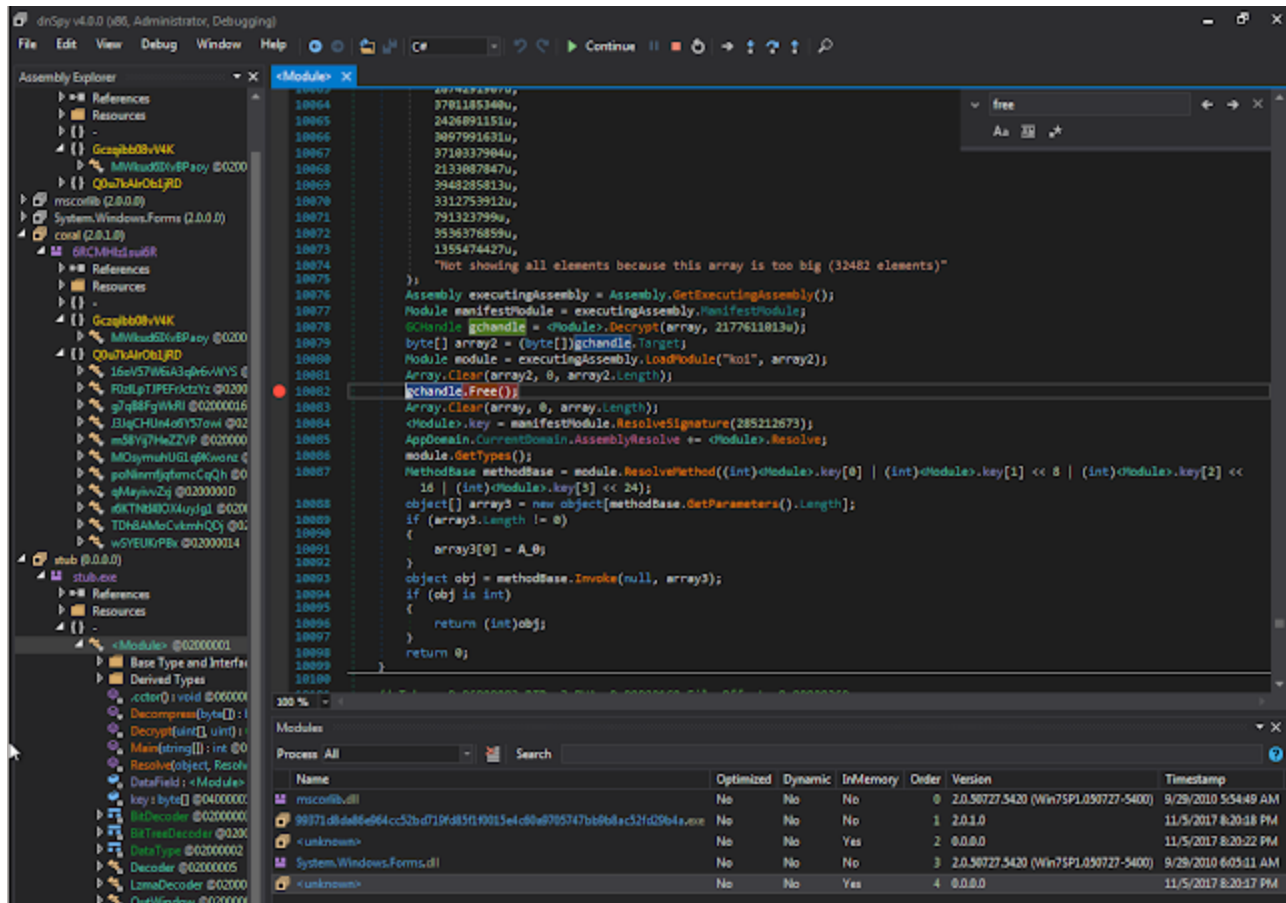


Fig. 6

As expected, this unpacks another module ConfuserEx is known for: `koi`.

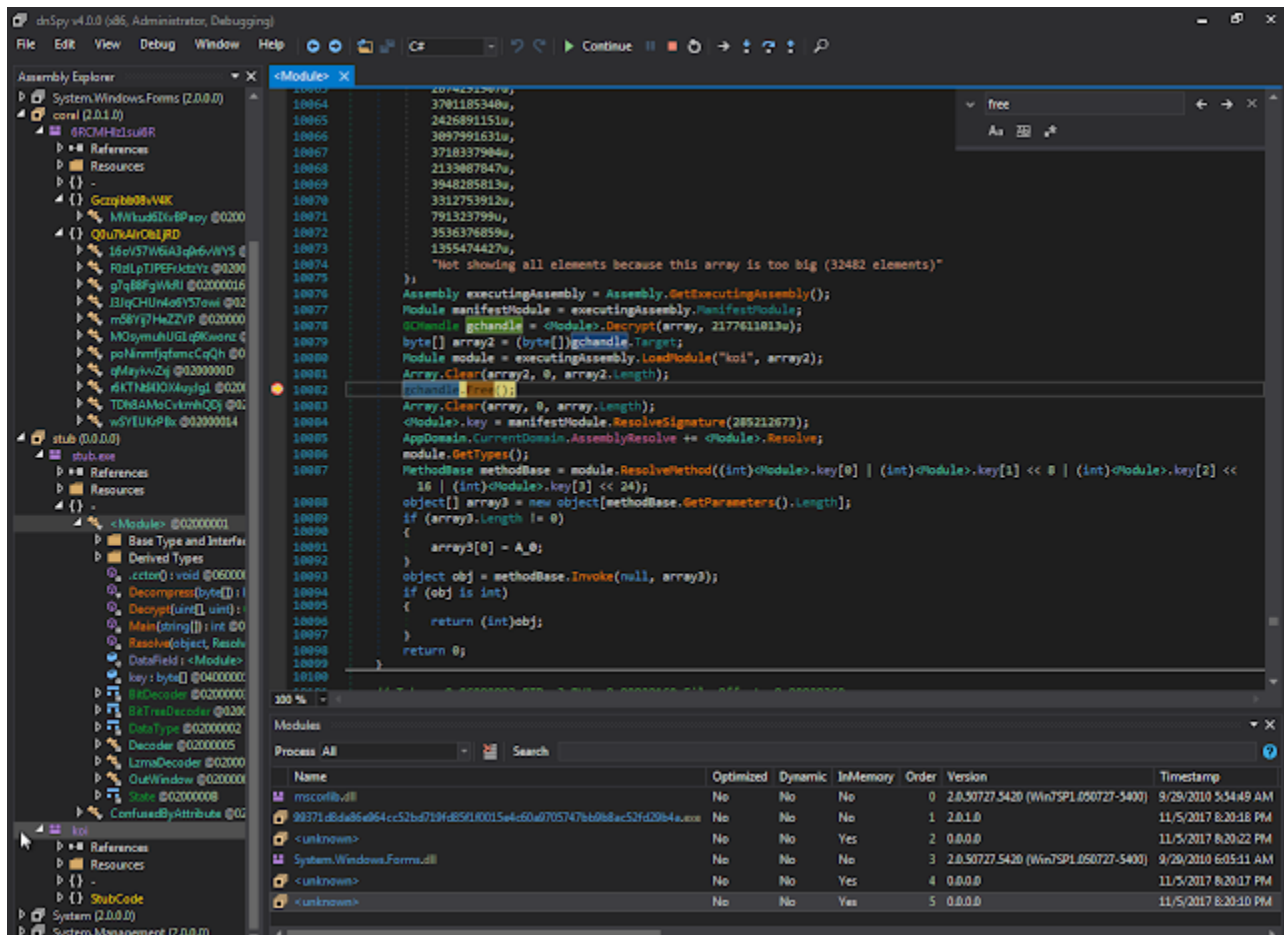


Fig. 7

We are getting closer, but the classes in koi are still empty and not yet filled with code:

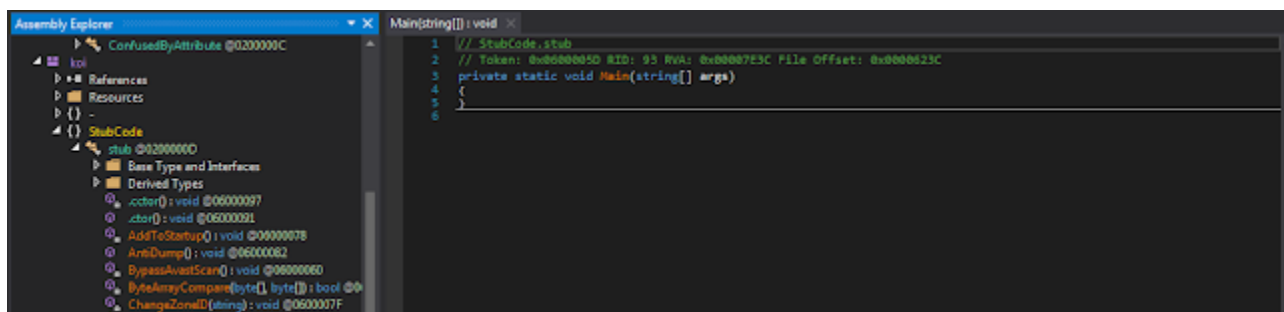


Fig. 8

Again, we set a breakpoint on the last method called in koi's ctor and proceed running the sample.

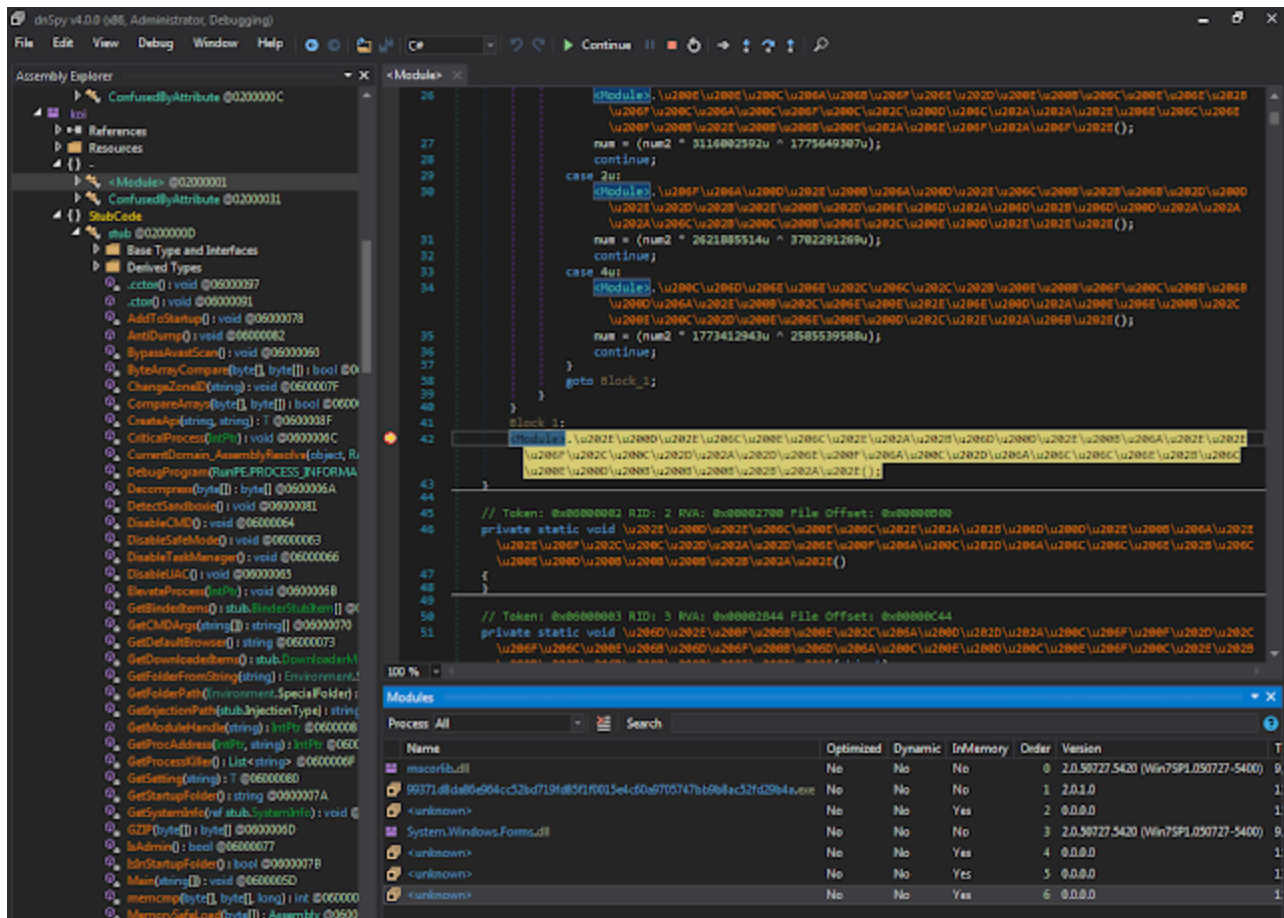


Fig. 9

Nice, another DLL is unpacked, unfortunately it is nothing important. Our Main class and most others in stub are still empty. Single stepping, brings us back into <module>. Once there, we analysed the methods and found out that we can set another breakpoint at line 92 for unpacking the next stage (see Fig. 10).

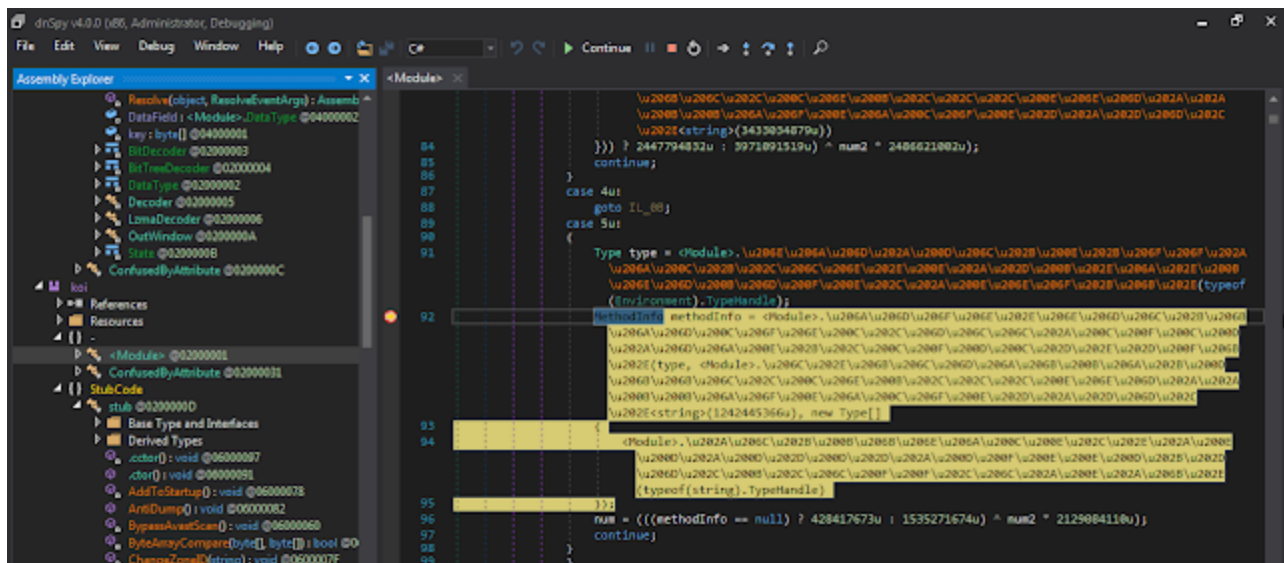


Fig. 10

Tada! If we now look in stub at the classes, they are filled with code. Now we can set a breakpoint on stub.Run() and start investigating what this malware loader is actually doing besides unpacking itself.

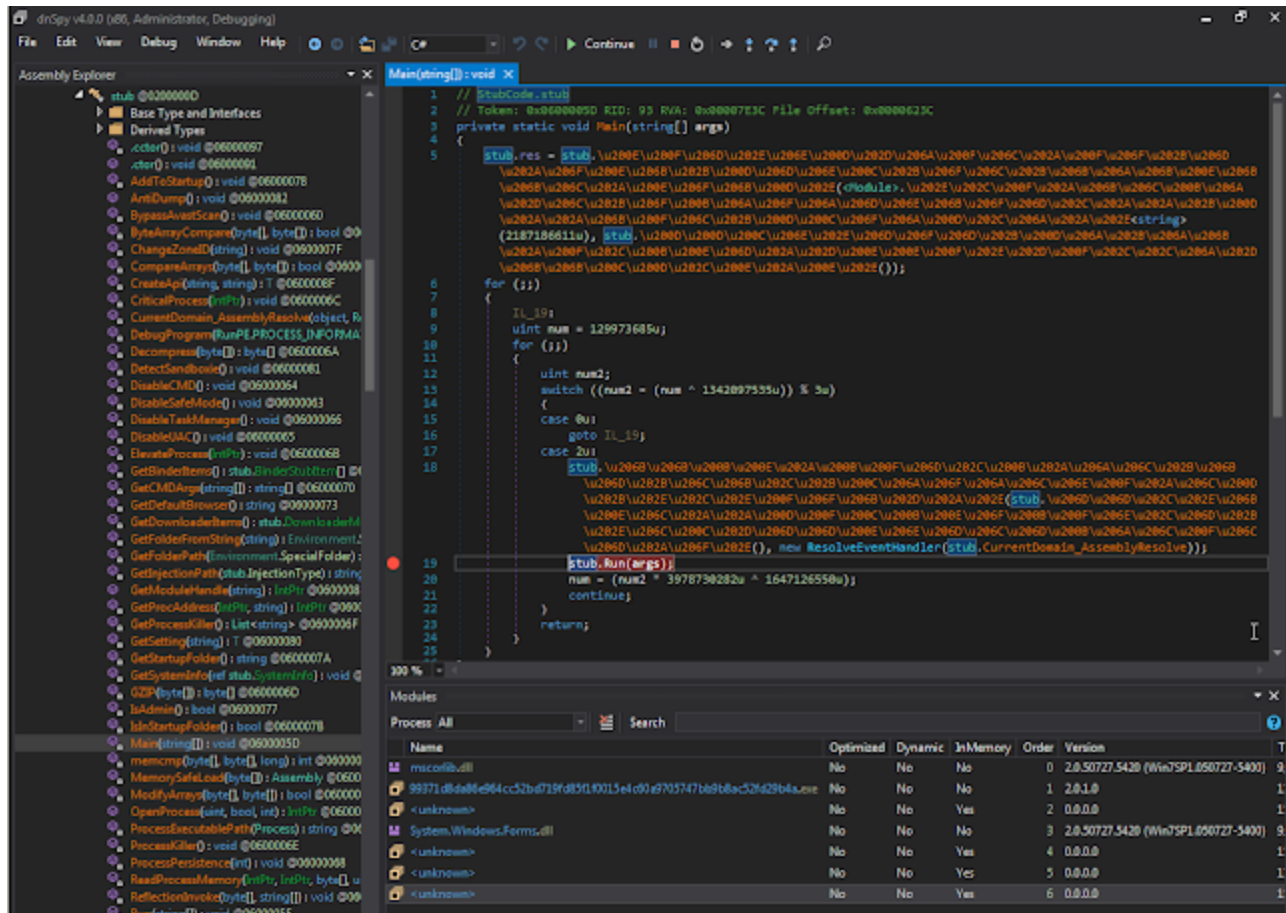


Fig. 11

We see that it is attempting to bypass some AV scans and reading several config parameters from the resource section. Below you can see the malware's configuration which was hidden encrypted in the resource section (Fig. 12) before unpacking.

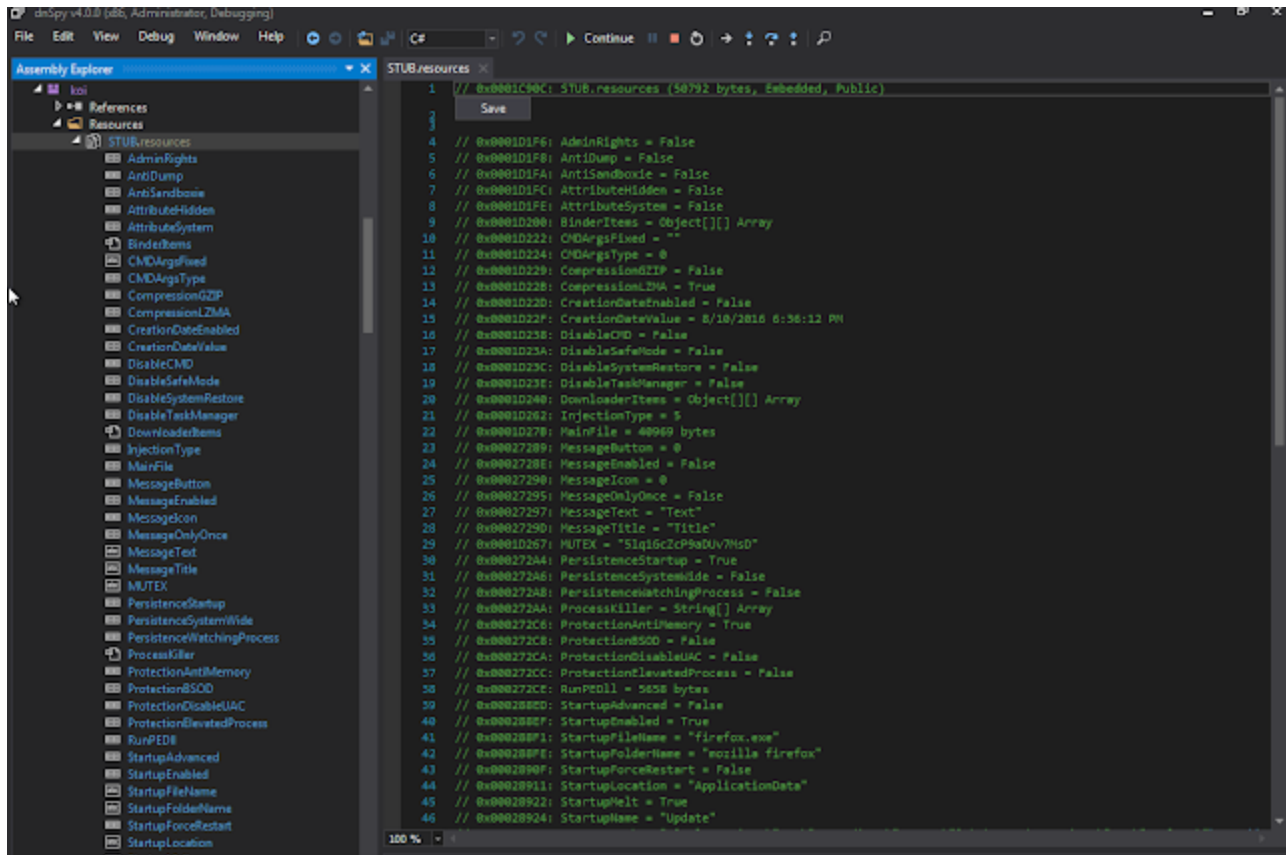


Fig. 12

It checks if it was executed from the Startup folder (e.g. %AppData%\mozilla firefox\firefox) as configured in the resource section. If not, it copies itself to the Startup folder and launches itself via cmd.exe. This means, we need to stop debugging and start again by loading the firefox.exe from %AppData%\mozilla firefox\firefox into dnSpy, following the unpacking again up to this point.

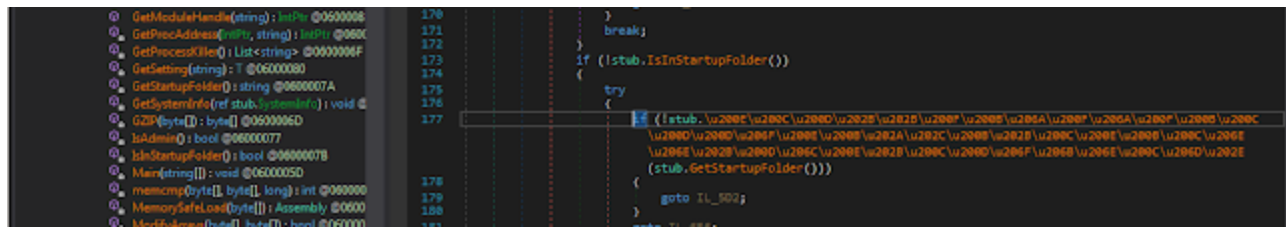


Fig. 13

Now we are in the "is executed from Startup location" branch. Here it gets interesting. First it makes itself persistent on the local machine. As you can see below, it writes a file called Update.txt with the following content to the %AppFolder%.

--- snip ---

C:\Users\ldex\AppData\Roaming\mozilla firefox\firefox.exe

exit

--- snip ---

```
252     }
253     IL_062:
254     stub.AddToStartup();
255     stub.SetAttributes();
256     byte[] array2 = null;
```

Fig. 14

```
2340     if (str0 == null && str1 == null && str2 == null)
2341     {
2342         return string.Empty;
2343     }
2344     if (str0 == null)
2345     {
2346         str0 = string.Empty;
2347     }
2348     if (str1 == null)
2349     {
2350         str1 = string.Empty;
2351     }
2352     if (str2 == null)
2353     {
2354         str2 = string.Empty;
2355     }
2356     int length = str0.Length + str1.Length + str2.Length;
2357     string text = string.FastAllocateString(length);
2358     string.FillStringChecked(text, 0, str0);
2359     string.FillStringChecked(text, str0.Length, str1);
2360     string.FillStringChecked(text, str0.Length + str1.Length, str2);
2361     return text;
```

Name	Value	Type
str0	"C:\Users\ldex\AppData\Local\Temp"	string
str1	"Update"	string
length	0x0000002A	int
text	"C:\Users\ldex\AppData\Local\Temp\Update.txt"	string

Fig. 15

Then it adds this file to the auto-run registry key by executing reg add in a cmd.exe to make sure the firefox.exe file gets executed at PC start up:

```
2144     public static string Format(string format, object arg0, object arg1)
2145     {
2146         return string.Format(null, format, new object[]
2147         {
2148             arg0,
2149             arg1
2150         });
2151     }
2152
2153     // Token: 0x060001C8 RID: 448 RVA: 0x00088E4 File Offset: 0x000878E4
2154     public static string Format(string format, object arg0, object arg1, object arg2)
2155     {
2156         return string.Format(null, format, new object[]
2157         {
2158             arg0,
```

Name	Value	Type
format	"reg add ""HKEY_CURRENT_USER\Software\Microsoft\Windows\CurrentVersion\Run"" /f /v ""[0]"" /d ""cmd /c type ""[1]"" cmd"" & exit"	string
arg0	"Update"	object (string)
arg1	"C:\Users\ldex\AppData\Local\Temp\Update.txt"	object (string)

Fig. 16a

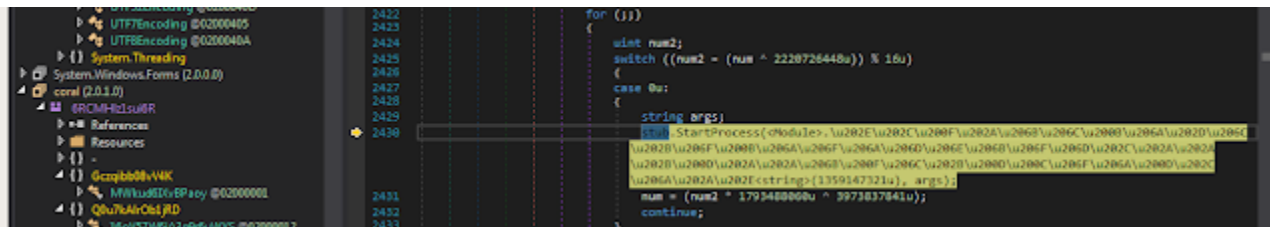


Fig. 16b

It executes a couple of other methods based on the configuration and then loads and decompresses the LZMA compressed malware payload file (Recam) from the resources MainFile section. After a couple of runtime fixes it loads RunPEDLL.dll and tries to inject the file into the user's browser. In case this fails (e.g. no browser is running), it injects the file into itself (firefox.exe). In both cases the RunPE.Run() method is used to do that.

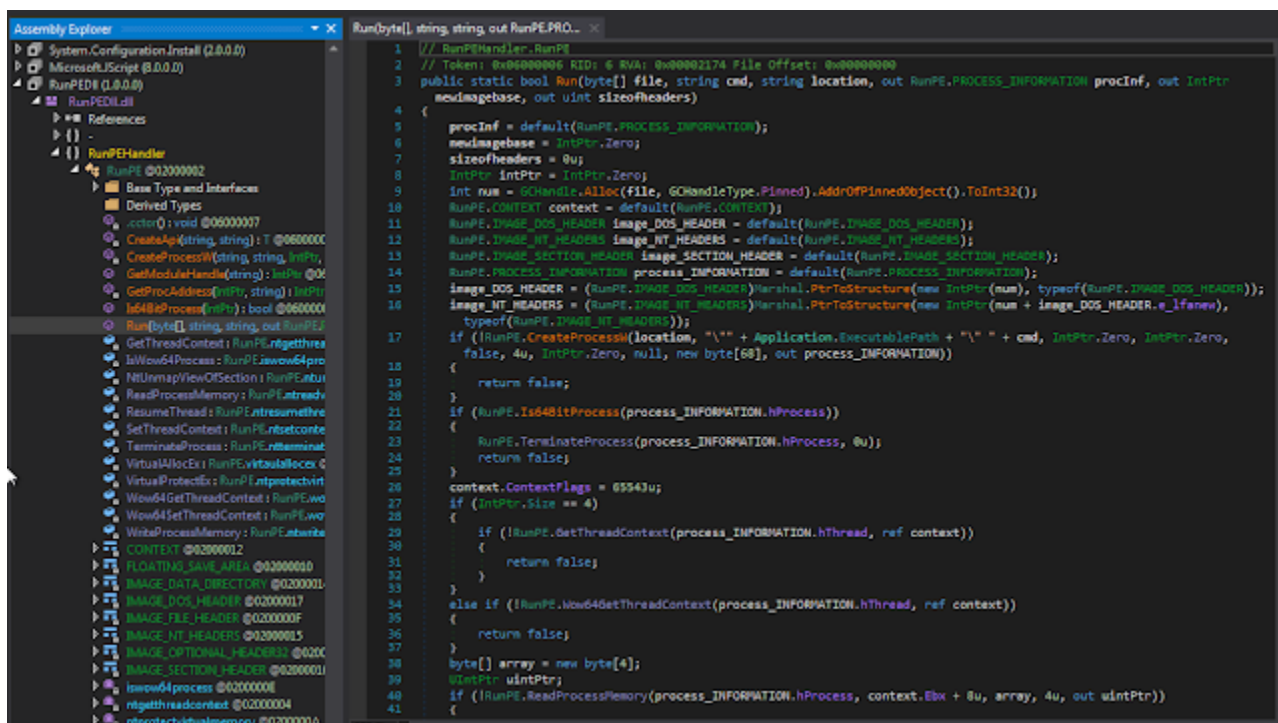


Fig. 17

From here on the work is done for the malware dropper and the loaded Recam binary takes over.

Payload

As mentioned in the introduction, the authors have gone the extra mile to frustrate analysis of the sample by using multiple obfuscation techniques, including multiple layers of data

encryption, string obfuscation, piecewise nulling, and data buffer constructors. It also relies on its own C2 binary protocol. All relevant data is heavily encrypted before transmission.

The dropped binary is packed with vanilla UPX. This part is easy to unpack; the tricky part comes in the next stage. After the original Entry Point (OEP) is restored, it begins with some homebrew cryptographic initialization for several values that get used consistently throughout runtime. Most remain constant following the initialization routine, but some change over time. Some preliminary string deobfuscation occurs shortly thereafter and includes a single hard-coded Command and Control server (C2) IP.

```
call Recam_string_decode2
mov [esp+12Ch+var_12C], ebx
mov [esp+12Ch+len], 0FFh
mov [esp+12Ch+ciphertext], offset decode2_var_len255
call Recam_string_decode2
mov [esp+12Ch+var_12C], ebx
mov [esp+12Ch+len], 20h
mov [esp+12Ch+ciphertext], offset aPassword ; "Password"
call Recam_string_decode2
mov [esp+12Ch+var_12C], ebx
mov [esp+12Ch+len], 10h
mov [esp+12Ch+ciphertext], offset HostID_plus_rand ; "HostId-%Rand%Ã"
call Recam_string_decode2
mov [esp+12Ch+var_12C], ebx
mov [esp+12Ch+len], 8
mov [esp+12Ch+ciphertext], offset mutex_name
```

This less frequently used deobfuscation routine is primarily based on a single-byte XOR loop. The other primary routine is JIT based and relies on a hard-coded decode key. Fortunately, IDA Pro's [Appcall](#) feature made short work of these obfuscations.

Direction	Type	Address	Text
Do...	p	sub_40B8C6+8CD	call Recam_string_JIT_decode; SMTP Password
Do...	p	sub_40B8C6+90A	call Recam_string_JIT_decode; EAS User
Do...	p	sub_40B8C6+93E	call Recam_string_JIT_decode; EAS Server URL
Do...	p	sub_40B8C6+971	call Recam_string_JIT_decode; EAS Password
Do...	p	sub_40C82F+6C	call Recam_string_JIT_decode
Do...	p	sub_40C82F+82	call Recam_string_JIT_decode; CryptUnprotectData
Do...	p	sub_40CB51+11	call Recam_string_JIT_decode; advapi32.dll
Do...	p	sub_40CB51+28	call Recam_string_JIT_decode; CredEnumerateA
Do...	p	sub_40CB51+44	call Recam_string_JIT_decode; CredFree
Do...	p	sub_40CB51+98	call Recam_string_JIT_decode
Do...	p	sub_40CB51+AE	call Recam_string_JIT_decode; CryptUnprotectData
Do...	p	sub_40D130+100	call Recam_string_JIT_decode; index.dat
Do...	p	sub_40D2FD+66	call Recam_string_JIT_decode; vaultcli.dll
Do...	p	sub_40D2FD+87	call Recam_string_JIT_decode; VaultOpenVault
Do...	p	sub_40D2FD+A3	call Recam_string_JIT_decode; VaultCloseVault
Do...	p	sub_40D2FD+C1	call Recam_string_JIT_decode; VaultEnumerateItems
Do...	p	sub_40D2FD+DD	call Recam_string_JIT_decode; VaultGetItem
Do...	p	sub_40D2FD+FB	call Recam_string_JIT_decode; VaultGetItem
Do...	p	sub_40D2FD+119	call Recam_string_JIT_decode; VaultFree
Do...	p	sub_40DFBA+1C	call Recam_string_JIT_decode; %s\Google\Chrome\User Data\Default>Login Data
Do...	p	sub_40E01F+1C	call Recam_string_JIT_decode; %s\Chromium\User Data\Default>Login Data
Do...	p	sub_40E084+1C	call Recam_string_JIT_decode; %s\Comodo\Dragon\User Data\Default>Login Data
Do...	p	sub_40E0E9+1C	call Recam_string_JIT_decode; %s\Yandex\YandexBrowser\User Data\Default>Login Data
Do...	p	sub_40E14E+1C	call Recam_string_JIT_decode; %s\Opera Software\Opera Stable>Login Data
Do...	p	sub_40E5C4+9B	call Recam_string_JIT_decode; GetModuleFileNameExA
Do...	p	sub_40E5C4+C7	call Recam_string_JIT_decode; GetModuleFileNameExA
Do...	p	sub_40E5C4+D5	call Recam_string_JIT_decode; DllPIWpr.iWWW
Do...	p	sub_40FAF8+61	call Recam_string_JIT_decode; %s\system32\cmd.exe
Do...	p	sub_4117A9+16	call Recam_string_JIT_decode; advapi32.dll
Do...	p	sub_4117A9+2C	call Recam_string_JIT_decode; GetUserNameA
Do...	p	sub_4117A9+46	call Recam_string_JIT_decode; USERNAME

Line 9 of 143

Fig. 18

Getting to the end of the preamble functions shortly following the PE Entry Point (EP), we get to an operation selection routine. The presence of unnecessary code and calculations disguises the fact that the jump to location 40849B will always be taken and the apparently interesting code that appears to involve file mangling and process creation is merely a decoy and always skipped in execution.

```

lea ebx, [esp+83Ch+var_618]
mov [esp+83Ch+lpValueName], 204h
mov [esp+83Ch+uExitCode], offset decode2_unk_len128
mov [esp+83Ch+Mode], ebx
call Recam_getenv
mov [esp+83Ch+uExitCode], 1
call Recam_arg0_AND_constant
test al, al
jz loc_40849B ; jmp taken (skip mangling & proc creation)

```

Recam_arg0_AND_constant proc near

```
var_1C= dword ptr -1Ch
arg_0= dword ptr 4

sub esp, 1Ch
mov [esp+1Ch+var_1C], offset flow_constant3
call Recam_base10_to_base16
and eax, [esp+1Ch+arg_0]
cmp eax, [esp+1Ch+arg_0]
setz al
add esp, 1Ch
retn
```

Moving forward, the malware sets a Run key for system persistence. Near the end of the operations function, an additional thread is created to start up a keylogger component, logging to %APPDATA%\Logs with <DAY>-<MONTH>-<YEAR> as the file name format. Logged input is stored in the commonly seen bracket delimiters. However, as one might expect by now, the final data is encrypted before written to the file on disk.

Next, the malware will create an ID file entitled .Identifier. If such a file already exists in the PWD of the sample (extracted via the GetModuleFilename API), it is simply read in instead of created from scratch.



Fig. 19

Data to be written to the file is generated piece by piece and results in the following format:

- (4 bytes) Static ID
- (13 bytes) HostId-<6 character alphanumeric rand, seeded from system time>
- (19 bytes) 19 null bytes
- (19 bytes) system time OR local time (19 byte format)
- (13 bytes) 13 null bytes

```

loc_4001E9:
lea   eax, [esp+4CCh+var_490]
mov   [esp+4CCh+Count], 7 ; int
mov   [esp+4CCh+Mode], 20h ; size_t
mov   [esp+4CCh+Filename], eax ; char *
call  Recam_get_time
mov   [esp+4CCh+Filename], offset a0000 ; "00000"
call  Recam_string_Jif_decode ; 2RAMd
mov   [esp+4CCh+Count], offset HostID_plus_rand ; "RR|+0+Vr\ue3(b)2W"
mov   [esp+4CCh+Mode], eax
mov   [esp+4CCh+Filename], 0
call  Recam_Find_char
test  eax, eax
js    short loc_40027D

add   eax, offset HostID_plus_rand ; "RR|+0+Vr\ue3(b)2W"
mov   [esp+4CCh+Mode], 6
mov   [esp+4CCh+Filename], eax
call  Recam_write_RNG
jmp   short loc_40027D

loc_40027D:
lea   eax, [ebx+4]
mov   dword ptr [ebx], 6198573Ah ; RECAM ID
mov   [esp+4CCh+Count], 32
mov   [esp+4CCh+Mode], offset HostID_plus_rand ; "RR|+0+Vr\ue3(b)2W"
mov   [esp+4CCh+Filename], eax
call  Recam_write_to_buffer ; Appends 'HostID-<6_alphanum_rand>'
lea   eax, [esp+4CCh+var_490]
mov   [esp+4CCh+Count], 32
mov   [esp+4CCh+Mode], eax
lea   eax, [ebx+2Ah]
mov   [esp+4CCh+Filename], eax
call  Recam_write_to_buffer ; Appends system time
mov   [esp+4CCh+File], 0 ; int
mov   [esp+4CCh+Mode], 0 ; int
mov   [esp+4CCh+Count], 0 ; int
mov   [esp+4CCh+Filename], edi ; FILE *
call  Recam_jump_to_file_pos
mov   [esp+4CCh+Count], 16
mov   [esp+4CCh+Mode], offset custom_binary_input
mov   [esp+4CCh+Filename], ebp
call  Recam_encrypt_with_divisor
lea   edx, [esp+4CCh+var_470]
mov   [esp+4CCh+var_40C], 60
mov   [esp+4CCh+Count], 60
mov   [esp+4CCh+Mode], ebx
mov   [esp+4CCh+Filename], ebp
mov   [esp+4CCh+File], edx
mov   [esp+4CCh+Str], edx
call  Recam_tbyte_XOR
mov   ebx, [esp+4CCh+Str]
mov   [esp+4CCh+File], edi ; File
mov   [esp+4CCh+Count], 60 ; Count
mov   [esp+4CCh+Mode], 1 ; Size
mov   [esp+4CCh+Filename], edx ; Str
call  fwrite
cmp   eax, 0Ah
setz  dl

; File
mov   [esp+4CCh+Filename], edi
mov   byte ptr [esp+4CCh+Str], dl
call  fclose
mov   [esp+4CCh+Filename], esi ; LPCSTR
call  Recam_GetFileAttributes
mov   dl, byte ptr [esp+4CCh+Str]
test  al, al
mov   bl, dl
jz    short loc_40030F

```

Fig. 20

Note that since a static ID is used, the first 4 bytes of the file always remain the same (the cryptography used for the C2 data is much more complex).

Offset(h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	
00000000	31	80	0F	80	76	15	FB	9B	15	A0	ED	BB	A3	3C	06	93	1€·Ev.û>. i»f<."
00000010	0B	6F	1D	D8	DA	E9	36	A9	F2	76	AB	25	DB	E7	65	79	.o.0Ûé6@òv«%Ûçey
00000020	95	0D	01	C9	CC	0E	1A	6A	7C	F9	34	F8	3E	6D	46	24	•..ÉÏ..j ù4ø>mF\$
00000030	CE	63	65	8D	D7	41	D2	2B	FE	31	35	C0	54	70	65	F2	Ïce.*Aò+p15ÀTpeò
00000040	52	73	76	EA													Rsvè

Fig. 21


```

mov     [esp+4CCh+Filename], ebp
call    Recam_string_decode
xor     edx, edx
cmp     dword ptr [ebx], 61985734h ; RECAM ID
jz      loc_40833F

```

Fig. 22a

```

lea     eax, [ebx+4]
mov     dword ptr [ebx], 61985734h ; RECAM ID
mov     [esp+4CCh+Count], 32
mov     [esp+4CCh+Mode], offset HostID_plus_rand ; "RÄ|+8+-Yr\nEè\b}%W"
mov     [esp+4CCh+Filename], eax
call    Recam_write_to_buffer ; Appends 'HostId-<6_alphanum_rand>'

```

Fig. 22b

As with the .Identifier file, the initial C2 beacon will also always be 68 bytes in length. Each C2 message (both client & server) will use the following format:

- (4 bytes) Length of data following these bytes
- (1 byte) C2 command
- (n bytes) Data relevant to the command

It's often easiest to break on a few instructions prior to deciphering the C2 beacon for many malware families these days. Whether it was intentional or not, the authors decided to opt of a homebrew crypto scheme allowing for randomized beacon data for each run (only the length bytes & C2 command for the beacon remain the same), or their homebrew crypto implementation is severely complex & broken.

```

loc_40204F:
    cmp     esi, [esp+20h+var_14]
    jge     loc_40208F

loc_40208F:
    mov     ecx, [ebx+10h]
    mov     eax, [esp+20h+var_14]
    mov     [edx-10h], ecx
    mov     ecx, [ebx+14h]
    mov     [edx-0Ch], ecx
    mov     ecx, [ebx+18h]
    mov     [edx-7], ecx
    mov     ecx, [ebx+1Ch]
    mov     [edx-3], ecx
    add     esp, 10h
    pop     ebx
    pop     esi
    pop     ebp
    retn

sub_40208A:
    cdd     ebx, 10h
    sub     ecx, [ebx]
    sub     edx, 10h
    inc     esi
    mov     edi, ecx
    mov     ebp, cl
    movzx  edi, ptr ds:dword_410600[ebp+4]
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    movzx  edi, ds:dword_410200[edi+4]
    xor     edi, ds:dword_410800[ebp+4]
    mov     ebp, ecx
    shr     ebp, 10h
    movzx  ecx, ch
    mov     eax, ebp
    movzx  ecx, byte ptr ds:dword_410600[ecx+4]
    mov     ebp, al
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    xor     edi, ds:dword_410E00[ebp+4]
    xor     edi, ds:dword_410200[ecx+4]
    mov     [edx], edi
    mov     ecx, [ebx+4]
    mov     edi, ecx
    movzx  ebp, cl
    shr     edi, 10h
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    movzx  edi, byte ptr ds:dword_410600[edi+4]
    mov     edi, ds:dword_410200[edi+4]
    xor     edi, ds:dword_410800[ebp+4]
    mov     ebp, ecx
    shr     ebp, 10h
    movzx  ecx, ch
    mov     eax, ebp
    movzx  ecx, byte ptr ds:dword_410600[ecx+4]
    mov     ebp, al
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    xor     edi, ds:dword_410E00[ebp+4]
    xor     edi, ds:dword_410200[ecx+4]
    mov     [edx+0], edi
    mov     ecx, [ebx+0Ch]
    mov     edi, ecx
    movzx  ebp, cl
    shr     edi, 10h
    movzx  edi, byte ptr ds:dword_410600[edi+4]
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    mov     edi, ds:dword_410200[edi+4]
    xor     edi, ds:dword_410800[ebp+4]
    mov     ebp, ecx
    shr     ebp, 10h
    movzx  ecx, ch
    mov     eax, ebp
    movzx  ecx, byte ptr ds:dword_410600[ecx+4]
    movzx  ebp, al
    movzx  ebp, byte ptr ds:dword_410600[ebp+4]
    xor     edi, ds:dword_410E00[ebp+4]
    xor     edi, ds:dword_410200[ecx+4]
    mov     [edx+0Ch], edi
    jmp     loc_40204F

```

Fig. 23

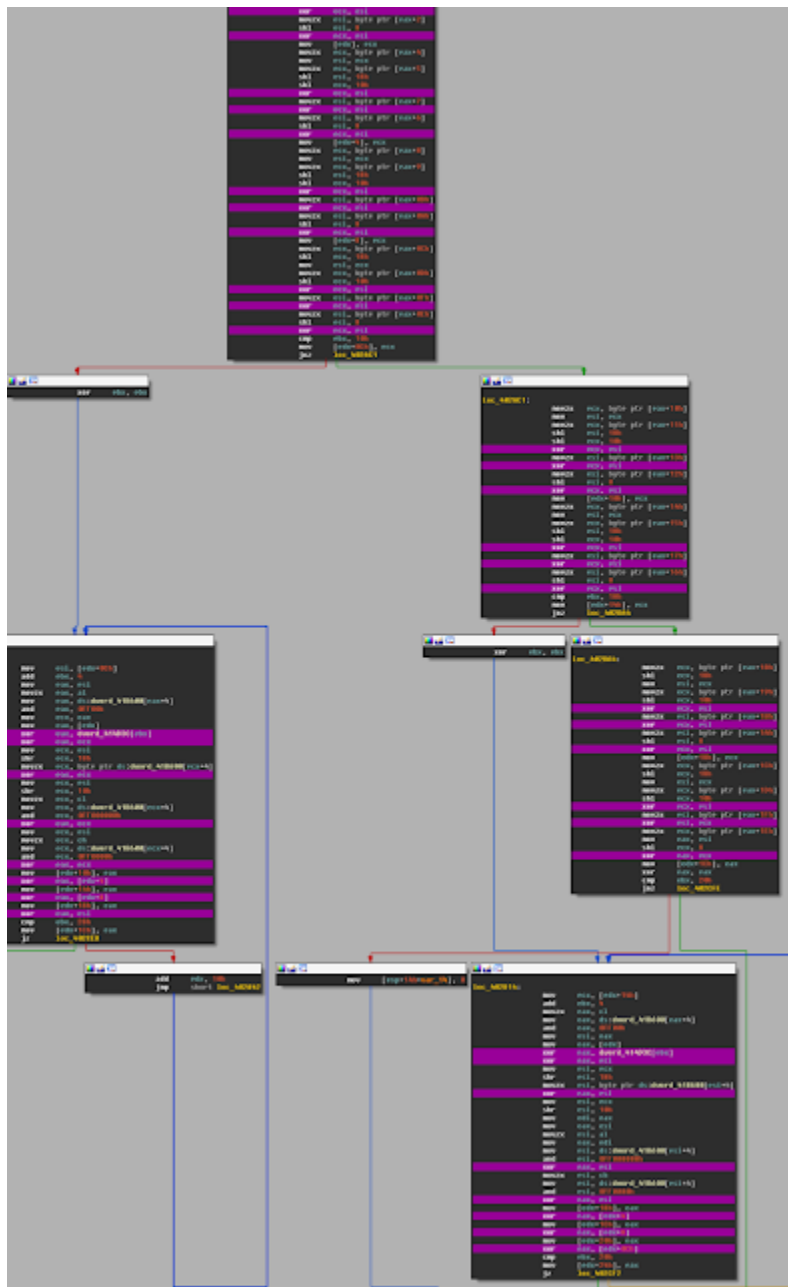


Fig. 24

The figure consists of four screenshots of network traffic analysis, likely from a packet capture tool like Wireshark. The top screenshot shows a sequence of beaconing messages. The second screenshot shows a server response to the beacon. The third and fourth screenshots show a client transmission and a series of 'keep alive' messages, respectively. The traffic is highlighted in purple, and the analysis is performed on a specific IP address.

Fig. 25

Once the beacon is sent, the sample waits for a server response. The C2 we encountered is now down and resetting connections, but pcaps captured in sandbox environments at an earlier date can give us a better idea of what to expect for the rest of the communications. The following example shows the beacon, the initial response, one additional client transmission and a series of "keep alive" messages consisting of the sole command byte.

```

00000000 41 00 00 00 83 92 98 5f b8 8e f8 48 e3 a8 4f 5f A....._ ...H..0_
00000010 b0 6c f8 ac 40 ad bb bd d9 86 c1 a0 48 d0 48 d2 .l..@... ..H.H.
00000020 84 80 86 a0 58 22 a1 66 fc 5f dd d6 ee 38 c4 b6 ....X".f ..8..
00000030 dc 4c af bb d3 a3 1b 73 09 c9 82 b5 85 85 f1 60 .L.....s .....`
00000040 91 29 79 dc 2a                                     .)y.*
    00000000 41 00 00 00 85 65 fe 8d 79 ac 99 49 c2 d2 9d 47 A....e.. y..I...G
    00000010 5c 59 b4 d4 3c 3d cb a9 6f 94 ad e7 7a 37 64 2d \Y..<=.. o...z7d-
    00000020 fa ed df a1 98 be cd d3 b4 4b 73 22 9b 12 4b c5 ..... .Ks"..K.
    00000030 2a ad 59 76 d9 6f 76 22 53 70 b2 ca bd 3b 6b 97 *.Yv.ov" Sp...;k.
    00000040 7f 53 64 28 b6                                     .Sd(.
00000045 3c 00 00 00 85 58 10 0a 85 88 a0 2e 8d 63 c3 40 <...X.. .....c.@
00000055 50 7f 48 73 0c 2b 3e 12 18 54 f5 4e 70 48 72 28 P.Hs.+>. .T.NpHr(
00000065 37 7e 91 f6 4d 16 53 b3 f1 f0 34 52 7d 49 6d c7 7~..M.S. ..4R}Im.
00000075 46 80 52 83 82 fb 15 c4 14 6c b7 45 9e 60 9b 38 F.R..... .l.E.`.8
    00000045 01 00 00 00 81                                     .....
00000085 01 00 00 00 81                                     .....
    0000004A 01 00 00 00 81                                     .....
0000008A 01 00 00 00 81                                     .....
    0000004F 01 00 00 00 81                                     .....
0000008F 01 00 00 00 81                                     .....
    00000054 01 00 00 00 81                                     .....
00000094 01 00 00 00 81                                     .....
    00000059 01 00 00 00 81                                     .....
00000099 01 00 00 00 81                                     .....
    0000005E 01 00 00 00 81                                     .....
0000009E 01 00 00 00 81                                     .....
    00000063 01 00 00 00 bd                                     .....
000000A3 01 00 00 00 bd                                     .....
    00000068 01 00 00 00 81                                     .....
000000A8 01 00 00 00 81                                     .....
    0000006D 01 00 00 00 c1                                     .....
000000AD 01 00 00 00 c1                                     .....
    00000072 01 00 00 00 c1                                     .....
000000B2 01 00 00 00 c1                                     .....
    00000077 01 00 00 00 81                                     .....
000000B7 01 00 00 00 81                                     .....
    0000007C 01 00 00 00 81                                     .....

```

Fig. 26

At this point, code execution depends on a flow state that is set only a few times throughout the binary (initially set to 0xFFFFFFFF). As far as the response length and C2 command are concerned, this state further dictates which each attribute must be. For example, the function responsible for checking the response length checks the flow state too. If state has changed, it checks if the message length exceeds 0x30000. If it's still in the default state, it checks if the length is 0x41 (length of the beacon message and its expected response). For the command byte itself, the default state checks if the command byte is set for the beacon phase of the communications (0x85). Once changed it will check to see if the command byte is less than or equal to 0xD2.



Fig. 27

```

Recam_verify_C2_command proc near
arg_0= dword ptr 4
arg_4= dword ptr 8

mov     edx, [esp+arg_0]
cmp     C2_flow_state, edx
mov     eax, [esp+arg_4]
jnz     short loc_4022C4 ; If flow state == 0xFFFFFFFF,
                        ; C2 command MUST be 0x85

loc_4022C4:
cmp     al, 85h ; If flow state == 0xFFFFFFFF,
               ; C2 command MUST be 0x85
setz   al
retn

Recam_verify_C2_command endp

```

Fig. 28

The response and subsequent data (if any) are relayed to a large jump table that is responsible for checking the command byte and proceeding from there with a particular action as issued by the server.

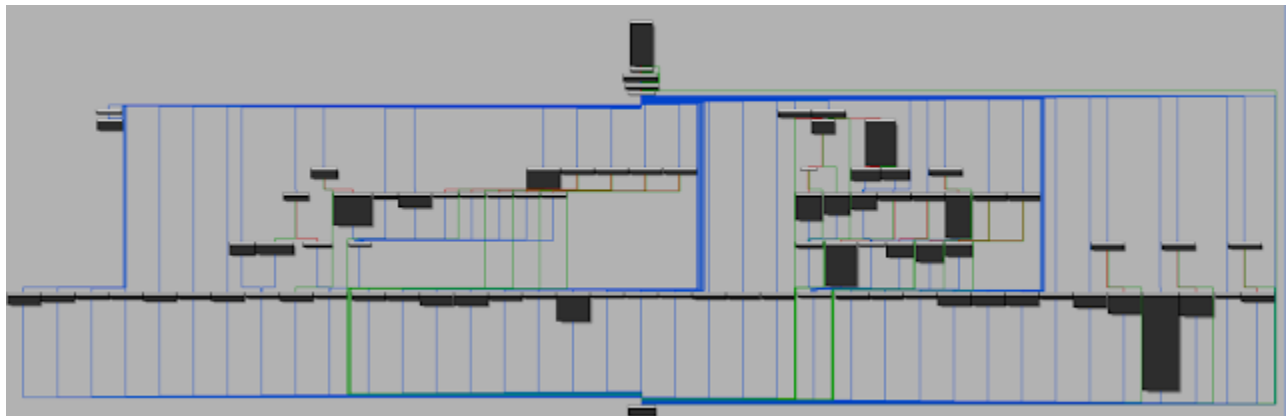


Fig. 29

The beginning of the previously mentioned function and jump table checks flow state again to see if the relevant parameter now equals a previously set state outside of the

0xFFFFFFFF. If this is the case, the data from the last server response is decrypted with the same routine used by the sample to encrypt data before transit. When in the default state, the command byte is passed to a LEA (Load Effective Address) where a calculated address is stored in EAX. In this case, there will be no calculated address due to the zero-extended command byte being referenced by the instruction. Instead, 0x7F gets added to the command byte. The single byte stored in AL taken from the DWORD stored in EAX is compared against 0x51. If equal, it proceeds to the function end and returns with no further action taken. Otherwise, the final byte stored in AL is zero-extended to EAX itself, multiplied by 4 and passed to the jump to determine the next action as requested by the server.



Fig. 30

As one might have gathered from the jump table, there are 82 possible commands that can be accepted from the server. However, not every command is unique. As we can see from the highlighted jump offsets below, many lead to the address shown earlier (RVA 0x227C) that is jumped to when completing no action.

```

data:00414094 C2_jump_table_offset dd offset loc_4010E9, offset loc_401106, offset loc_401122
data:00414094 ; DATA XREF: Reach_process_command+0E9
data:00414094 dd offset loc_401335, offset loc_401106, offset loc_401211 ; jump table for switch state
data:00414094 dd offset loc_401368, offset loc_401380, offset loc_40138C
data:00414094 dd offset loc_4013D1, offset loc_4015D7, offset loc_4015E8
data:00414094 dd offset loc_4015E8, offset loc_4015E8, offset loc_401C00
data:00414094 dd offset loc_401780, offset loc_401CA9, offset loc_401631
data:00414094 dd offset loc_401780, offset loc_401830, offset loc_4019F8
data:00414094 dd offset loc_401A81, offset loc_401B1F, offset loc_40189E
data:00414094 dd offset loc_401BC3, offset loc_401BE8, offset loc_4017D9
data:00414094 dd offset loc_401B3D, offset loc_401D73, offset loc_401DC8
data:00414094 dd offset loc_401D81, offset loc_401D81, offset loc_401D01
data:00414094 dd offset loc_401D81, offset loc_401DE5, offset loc_401D81
data:00414094 dd offset loc_401DF9, offset loc_401DF9, offset loc_401E00
data:00414094 dd offset loc_401E32, offset loc_401E46, offset loc_401E53
data:00414094 dd offset loc_401E32, offset loc_401868, offset loc_40187D
data:00414094 dd offset loc_401892, offset loc_4018CD, offset loc_40184D
data:00414094 dd offset loc_401892, offset loc_401892, offset loc_401908
data:00414094 dd offset loc_401892, offset loc_401984, offset loc_4019A7
data:00414094 dd offset loc_40198F, offset loc_40198F, offset loc_40198F
data:00414094 dd offset loc_401F79, offset loc_401F79, offset loc_401FA5
data:00414094 dd offset loc_401FA5, offset loc_401FD1, offset loc_401FD1
data:00414094 dd offset loc_401FFD, offset loc_401FFD, offset loc_401FFD
data:00414094 dd offset loc_401FFD, offset loc_401FFD, offset loc_401FFD
data:00414094 dd offset loc_401C88, offset loc_401C88, offset loc_401C88
data:00414094 dd offset loc_401D64, offset loc_402054, offset loc_402054
data:00414094 dd offset loc_402068, offset loc_402068, offset loc_4020A9
data:00414094 dd offset loc_40226D, offset loc_40226D, offset loc_40226D

```

Fig. 31

While time did not allow us to deeply examine each and every path, gathered sandbox pcaps along with our understanding of the command protocol allowed us to examine the commands sent by the server and calculate the jumps ourselves. Here are some examples of the functionality available in this variant of Recam, given the command.

- 0x85 (case 4) - Process initial server acknowledgement and set flow state
- 0x81 (case 0) - Keep-alive message
- 0xBD (case 12) - Download file to %TEMP% OR download file to %TEMP% and create new process
- 0x87 (case 6) - Create new process from argument
- 0x89 (case 8) - Close network socket, release mutex, call WSACleanup, & terminate process.

IOC

Malicious Word Document:

C3b1a98c6bc9709f964ded39b288aff66abc5c39b9662fdd28ddfcc178152d67

Dropper binary:

99371d8da86e964cc52bd719fd85f1f0015e4c60a9705747bb9b8ac52fd29b4a

Payload (Recam):

1fd8520246c75702c000f4fac3f209d611c21bfdb81df054c9558d5e002a85ce

Command and Control Server:

185.140.53.212

Domain registered to the Command and Control Server

U811696.nvpn.so

Domain Owner

meridoncharles@yahoo.com

Directories and Files

%AppData%\mozilla firefox\firefox.exe

%AppData%\mozilla firefox\Identifier

%Temp%\Update.txt

Registry:

Key: HKEY_CURRENT_USER\Software\Microsoft\Windows\CurrentVersion\Run\Update

Value: cmd /c type C:\Users\dex\AppData\Local\Temp\Update.txt | cmd

Similar files:*Dropper related:*

006583023242bf4a8dcd0190aef32500dcacfacf1dda7b24409133dfccbfd186
0086bb92bd34b41e180bec90dd15d4b0d0eb9c7384a68b66354d603ad8e14706
01226da791e32f8cc907f88b2b672068b78b86b1a0d154bd22274234a7d9b5e6
031046538b60f9b243aa74bdec2a13ab2aee4b941a136daca12de78c3419dd6e
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03fe0caa0f3e21f1975bdceaa5e38e00d725879322c8815e7b3b60b1c0cfec20
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14d341bd2c2b7ea90b3b15fcb8154caa036b629156199d955047fb081de6669d
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6e74c3a1f5bc17f38da7305068f9203943090067b2dfa67fb9361fa1821a674b
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7c127f9f22a346d4369405f540c4e467abac02bcf4589dbd6d3e97f18fbe0f98
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c064c70d887ea1ccf37e3edd19dc717b53f2c0513424416fc783654ab9f61134
c13d2ce1b561ed02c02c9a2899a7d7e9c97400ef8f81eb48d6ba02a90d76c689
c2fef90a5ad97a030712b9245fccffb26f226370bd6f67c989bbf95e02ca03b8
c8d98b7f1cf8811c3e75b0ec8c1011d0ed4c3951420f04e3b743da564869e02b
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d7d94b3aea52f0a9bda593806d56f962c09f4d389820cd9016ceed73767feb55
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f7ac275c2971820fb311f41049c5e460251486625696caa0fc60d0ffc683fb70
fc0bd596957f70e87954ef42acc0bfe0123c11be6879c7c910a4a7e434e36c46

RTF Document related:

27bf1851e64f5e6d6e33b2b3bc89b82dda2da2fd9a747c847c148909dda028d3
7a63150ebd09fd4f8c8f1b7485a0139108f22aea043c9759239d9976948b5c75
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aa3f5e91fb4f447bfff93123c404a62b8ebe2a9790ebbd02c57f5ab61d2d2882
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Payload related:

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Conclusion

Malware is a moving target, it is constantly evolving in an arms race between the malware authors and the security researchers. This analysis shows the level of sophistication employed by threat actors in order to attempt to escape detection.

Obfuscation is an art form. Techniques can range from frequently changed packers to the multiple techniques employed in malware such as this. Often malware packers are modified by their authors very soon after deobfuscation tools or reports are publically released. In

many cases it is enough for them to change minor parts of the obfuscator to confuse the deobfuscation tools. Hence, malware researchers can't rely on these tools and must resort to be able to manually deobfuscate code when necessary.

Understanding the steps that threat actors will go to to hide from detection and analysis is vital when it comes to protecting systems from malware. It is by applying lessons learnt from analyses such as this, that we are able to detect advanced malware with tools such as Advanced Malware Protection ([AMP](#)) and [Threatgrid](#).

Coverage

Additional ways our customers can detect and block this threat are listed below.

PRODUCT	PROTECTION
AMP	✓
CloudLock	N/A
CWS	✓
Email Security	✓
Network Security	✓
Threat Grid	✓
Umbrella	✓
WSA	✓

Advanced Malware Protection ([AMP](#)) is ideally suited to prevent the execution of the malware used by these threat actors.

[CWS](#) or [WSA](#) web scanning prevents access to malicious websites and detects malware used in these attacks.

[Email Security](#) can block malicious emails sent by threat actors as part of their campaign.

Network Security appliances such as [NGFW](#), [NGIPS](#), and [Meraki MX](#) can detect malicious activity associated with this threat.

[AMP Threat Grid](#) helps identify malicious binaries and build protection into all Cisco Security products.

[Umbrella](#), our secure internet gateway (SIG), blocks users from connecting to malicious domains, IPs, and URLs, whether users are on or off the corporate network.

Open Source Snort Subscriber Rule Set customers can stay up to date by downloading the latest rule pack available for purchase on [Snort.org](#).

