

Zloader | ThreatLabz

zscaler.com/blogs/security-research/zloader-learns-old-tricks

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Introduction

Zloader (a.k.a. Terdot, DELoader, or Silent Night) is a modular trojan based on leaked ZeuS source code. As detailed in our [previous blog](#), Zloader reemerged following an almost two-year hiatus with a new iteration that included modifications to its obfuscation techniques, domain generation algorithm (DGA), and network communication.

Most recently, Zloader has reintroduced an anti-analysis feature similar to one that was present in the original ZeuS 2.x code. The feature restricts Zloader's binary execution to the infected machine. This characteristic of ZeuS was abandoned by many malware variants derived from the leaked source code including Zloader, until now. In this blog post, we explain how this anti-analysis feature works and how it differs from the original ZeuS implementation.

Key Takeaways

- Zloader (a.k.a. Terdot, DELoader, or Silent Night) is a modular trojan based on the leaked ZeuS source code dating back to 2015.
- Zloader has continued to evolve since its resurrection around September 2023 after an almost two-year hiatus.
- The latest version, 2.4.1.0, introduces a feature to prevent execution on machines that differ from the original infection. A similar anti-analysis feature was present in the leaked ZeuS 2.X source code, but implemented differently.

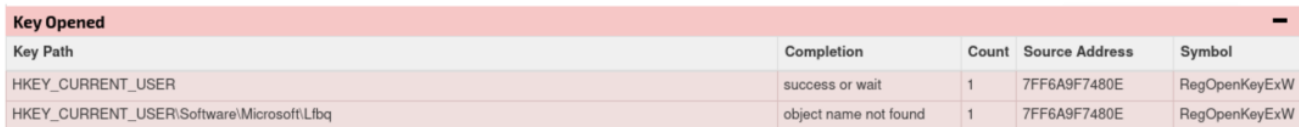
Technical Analysis

In the upcoming sections, we explore the technical intricacies of Zloader's latest anti-analysis feature introduced in versions 2.4.1.0 and 2.5.1.0. We also draw comparisons to ZeuS to provide a comprehensive understanding of their respective approaches.

Registry check

Zloader samples with versions greater than 2.4.1.0 will abruptly terminate if they are copied and executed on another system after the initial infection. This is due to a Windows registry check for the presence of a specific key and value.

The screenshot below shows the Windows Registry check failing in a malware sandbox.



Key Opened				
Key Path	Completion	Count	Source Address	Symbol
HKEY_CURRENT_USER	success or wait	1	7FF6A9F7480E	RegOpenKeyExW
HKEY_CURRENT_USER\Software\Microsoft\Libq	object name not found	1	7FF6A9F7480E	RegOpenKeyExW




Figure 1: Registry key check performed in a sandbox.

The registry key and value are generated based on a hardcoded seed that is different for each sample.

The Python code below replicates the algorithm to generate the registry key.

```
#!/usr/bin/env python3
SEED = 0x1C5EE76F0FE82329
def calculate_registry_key(seed):
    key = ""
    key_length = 1 + seed % 8

    if key_length < 4:
        key_length = 4

    for i in range(key_length):
        key += chr(seed % 0x1A + 0x61)
        seed = (((seed << 8) | (seed >> (64 - 8))) & 0xffffffffffffffff) + 1

    key = key.capitalize()
    return key
print(calculate_registry_key(SEED))
```

If the registry key/value pair is manually created (or this check is patched), Zloader will successfully inject itself into a new process. However, it will terminate again after executing only a few instructions. This is due to a secondary check in Zloader’s MZ header.

MZ header check

A bit further in the code, there is an additional check that involves a **DWORD** present in the MZ header at the offset 0x30, which is only executed after being injected into a new process. The **DWORD** used in the check of the analyzed sample can be seen in the image below.

Offset (h)	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F	
00000000	4D	5A	90	00	03	00	00	00	04	00	00	00	FF	FF	00	00	MZ.....ÿÿ..
00000010	B8	00	00	00	00	00	00	00	40	00	00	00	00	00	00	00@.....
00000020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
00000030	44	12	D0	AA	00	00	00	00	00	00	00	00	E8	00	00	00	D.D.....è..
00000040	0E	1F	BA	0E	00	B4	09	CD	21	B8	01	4C	CD	21	54	68	..°..'í!..Lí!Th
00000050	69	73	20	70	72	6F	67	72	61	6D	20	63	61	6E	6E	6F	is program cannot
00000060	74	20	62	65	20	72	75	6E	20	69	6E	20	44	4F	53	20	be run in DOS
00000070	6D	6F	64	65	2E	0D	0D	0A	24	00	00	00	00	00	00	00	mode....\$.....
00000080	2D	23	CB	40	69	42	A5	13	69	42	A5	13	69	42	A5	13	-#Ë@iB¥.iB¥.iB¥.
00000090	7D	29	A3	12	68	42	A5	13	7D	29	A4	12	6E	42	A5	13	})£.hB¥.})¤.nB¥.

Figure 2: MZ header with random **DWORD** at 0x30 offset.

The **DWORD** at the 0x30 offset is part of the ten reserved **WORDS** that go from offset 0x28 to offset 0x3C of the MZ header. These bytes are usually **null**. However, in the example above, the malware contained an integer value (0xAAD01244), which is compared with the file size (0x29A00). Since this integer is a very large number, the check fails. The decompiled code of the file size check is shown in the figure below.

```

loc_7FF7721E89A9:
lea rcx, [rsp+8038h+var_C98]
call Get_MZ_Offset_sub_7FF772208000
add rax, 30h ; 0
mov [rsp+8038h+Pointer_to_MZHeader_DWORD], rax
mov rcx, [rsp+8038h+Pointer_to_MZHeader_DWORD]
mov esi, [rax]
mov rcx, 467E9CA5DE43798Fh
call XOR_Return_13_sub_7FF7721F1050
add rsi, rax
lea rcx, [rsp+8038h+var_C98]
call FileSize_sub_7FF772201220
cmp rsi, rax

```

```

169 if ( (Call_GetFileAttributes_GetFileSize_sub_7FF7721F89F0(v67, (__int64)v71) & 1) != 0 )
170 {
171     Pointer_to_MZHeader_DWORD = (_DWORD *) (Get_MZ_Offset_sub_7FF772208000((__int64)v71) + 0x30);
172     Seed_Offset = (unsigned int) *Pointer_to_MZHeader_DWORD;
173     MZHeader_DWORD_Final = XOR_Return_13_sub_7FF7721F1050(0x467E9CA5DE43798F164) + Seed_Offset;
174     if ( MZHeader_DWORD_Final < FileSize_sub_7FF772201220((__int64)v71) )
175     {
176         v103 = (_QWORD *) ((unsigned int) *Pointer_to_MZHeader_DWORD + Get_MZ_Offset_sub_7FF772208000((__int64)v71));
177         v6 = FileSize_sub_7FF772201220((__int64)v71);
178         v7 = Get_MZ_Offset_sub_7FF772208000((__int64)v71);
179         v102 = (unsigned int) sub_7FF7721E3C70(v7, v6);
180         v8 = (unsigned int *) (void) Call_Resolve_API_sub_7FF7721E1C70(0, GetTickCount);
181         v101 = v8();

```

Figure 3: Decompiled code of the file size check against the MZ **DWORD**.

What the malware developers are doing here is utilizing the additional MZ header **DWORD** as a pointer to the seed's offset, which explains the purpose of the check. This is due to the **DWORD** being overwritten after the initial execution. If the pointer points beyond the binary, it indicates that the seed has already been written, eliminating the need for reinitialization.

This suggests that the initial binary for system infection must include a **null** seed, with the MZ **DWORD** at 0x30 holding the seed's offset. Subsequently, this offset is initialized with a pseudo-random **QWORD** generated via the Mersenne Twister algorithm, leaving a hardcoded seed that differs per infected sample.

The figure below shows the decompiled code where the seed is being generated and written.

```

call VirtualProtect_Seed_sub_7FF7721E2460
call MersenneTwister_sub_7FF7721E4330
mov rcx, [rsp+8038h+Pointer_to_MZHeader_DWORD]
mov [rcx], eax
cmp cs:Seed, 0

```

```

184 VirtualProtect_Seed_sub_7FF7721E2460(v77);
185 *Pointer_to_MZHeader_DWORD = MersenneTwister_sub_7FF7721E4330();
186 if ( Seed )
187 {
188     sub_7FF7721FB4E0(v79);

```

Figure 4: Decompiled code where the seed is first created.

Without the seed and MZ header values set correctly, the Zloader sample won't run or install on a different machine, unless it is patched or if the environment is replicated with all the registry and disk paths/names, alongside all the original artifacts from the original victim's machine.

Registry value content

In previous versions of Zloader, there was a single registry key and value containing some machine information (install path, computer/bot ID, victim-specific RC4 key, etc.), similar to the ZeuS **PeSettings** we will examine in the next section. The key/value pair was encrypted with the ZeuS **VisualEncrypt** algorithm and RC4, using the RSA key present in the static configuration as the key, but it wasn't used to avoid infecting a new machine, as it was created again when executed in a different environment.

Now, there is an additional value created using the seed previously mentioned.

The figure below shows the registry keys and values added to the victim's system during the infection process.

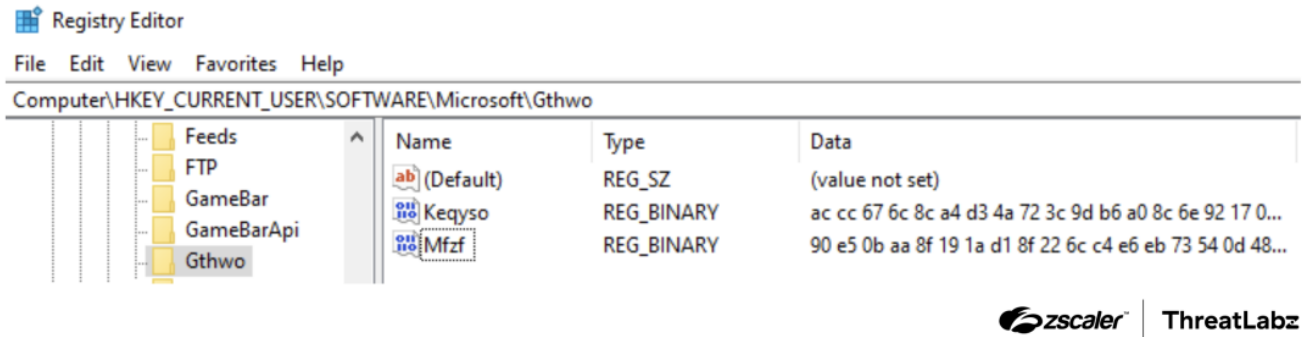


Figure 5: Registry keys and values added when infecting the machine.

The content has a fixed length of 1,418 bytes and is encrypted with RC4, but without the additional **VisualEncrypt** layer. The RC4 key is also based on the seed generated while performing the infection, which is then used to create the names of the registry key and value.

The decrypted format and content are as follows:

```

00000000 41 00 64 00 6f 00 62 00 65 00 5c 00 49 00 6e 00 |A.d.o.b.e.\.I.n.|
00000010 66 00 72 00 61 00 42 00 61 00 73 00 65 00 2e 00 |f.r.a.B.a.s.e...|
00000020 65 00 78 00 65 00 00 00 00 00 00 00 00 00 00 00 |e.x.e.....|
00000030 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000040 57 00 61 00 62 00 75 00 75 00 5c 00 45 00 66 00 |W.a.b.u.u.\.E.f.|
00000050 79 00 63 00 79 00 64 00 6d 00 61 00 00 00 00 00 |y.c.y.d.m.a....|
00000060 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000070 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000080 57 00 61 00 62 00 75 00 75 00 5c 00 47 00 65 00 |W.a.b.u.u.\.G.e.|
00000090 78 00 61 00 6e 00 69 00 00 00 00 00 00 00 00 00 |x.a.n.i.....|
000000a0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000000b0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000000c0 57 00 61 00 62 00 75 00 75 00 5c 00 4c 00 6f 00 |W.a.b.u.u.\.L.o.|
000000d0 6b 00 61 00 79 00 6c 00 62 00 6f 00 00 00 00 00 |k.a.y.l.b.o....|
000000e0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000000f0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000100 57 00 61 00 62 00 75 00 75 00 5c 00 47 00 79 00 |W.a.b.u.u.\.G.y.|
00000110 79 00 70 00 6b 00 00 00 00 00 00 00 00 00 00 00 |y.p.k.....|
00000120 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000130 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000140 57 00 61 00 62 00 75 00 75 00 5c 00 45 00 71 00 |W.a.b.u.u.\.E.q.|
00000150 71 00 61 00 00 00 00 00 00 00 00 00 00 00 00 00 |q.a.....|
00000160 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000170 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000180 57 00 61 00 62 00 75 00 75 00 5c 00 59 00 77 00 |W.a.b.u.u.\.Y.w.|
00000190 77 00 75 00 00 00 00 00 00 00 00 00 00 00 00 00 |w.u.....|
000001a0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000001b0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000001c0 57 00 61 00 62 00 75 00 75 00 5c 00 49 00 76 00 |W.a.b.u.u.\.I.v.|
000001d0 76 00 65 00 64 00 00 00 00 00 00 00 00 00 00 00 |v.e.d.....|
000001e0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000001f0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000200 57 00 61 00 62 00 75 00 75 00 5c 00 48 00 61 00 |W.a.b.u.u.\.H.a.|
00000210 6b 00 6f 00 67 00 69 00 00 00 00 00 00 00 00 00 |k.o.g.i.....|
00000220 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000230 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000240 59 00 66 00 6f 00 77 00 76 00 6f 00 5c 00 46 00 |Y.f.o.w.v.o.\.F.|
00000250 75 00 76 00 61 00 61 00 71 00 00 00 00 00 00 00 |u.v.a.a.q.....|
00000260 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000270 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
00000280 59 00 66 00 6f 00 77 00 76 00 6f 00 5c 00 4d 00 |Y.f.o.w.v.o.\.M.|
00000290 79 00 6c 00 75 00 6b 00 00 00 00 00 00 00 00 00 |y.l.u.k.....|
000002a0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000002b0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
000002c0 59 00 66 00 6f 00 77 00 76 00 6f 00 5c 00 45 00 |Y.f.o.w.v.o.\.E.|
000002d0 73 00 6e 00 6f 00 00 00 00 00 00 00 00 00 00 00 |s.n.o.....|
000002e0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|

```

*

0000058A

The structure is divided into 64 bytes for each entry. The first structure is the binary path inside **%APPDATA%**, and the following are the Zloader modules.

Zeus implementation

It's been thirteen years since the Zeus 2.0.8 source code was leaked, but it is still widely leveraged by threat actors and some of its concepts are still relevant. The technique described in the section above, and used by Zloader to store the installation information and avoid being run on a different system, was also performed by Zeus v2, but implemented in a different way.

In Zeus, the binary had an overlay section called **PeSettings**, where the installation information was stored instead of in the registry. The encrypted Zeus overlay section is shown in the figure below.

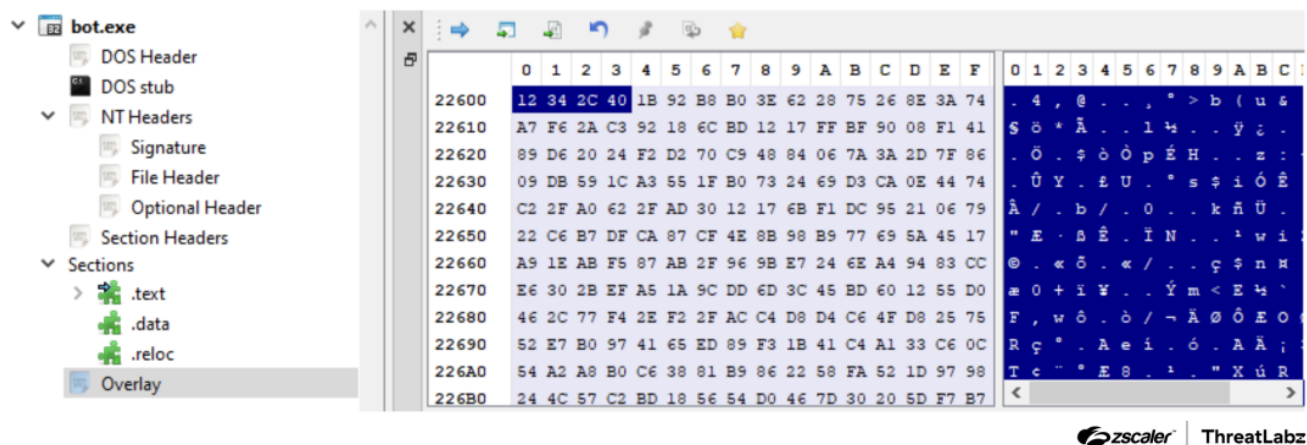


Figure 6: The encrypted Zeus overlay section.

The header of this section is decrypted with the RC4 key present in the static config. The figure below shows the Zeus section header.

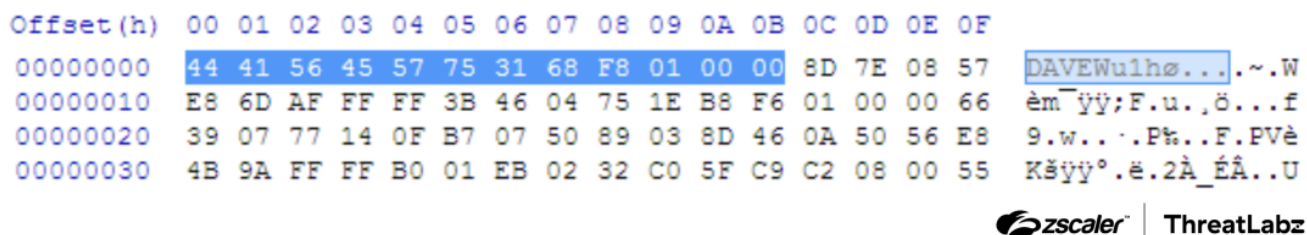


Figure 7: Zeus overlay section header.

The decrypted header is composed of three **DWORDS**:

- Magic word (**DAVE**)
- CRC32 of the data
- Size of the data

If the size of the data is equal to 0xC, it means the trojan is not installed and will proceed with the infection to generate all the required information, such as the computer/bot ID, install paths, and machine-specific RC4 key, which is generated per install and stored as an

initialized RC4 S-box.

Then, ZeuS will encrypt the **PeSettings** again and replace the overlay data with it, while changing the header CRC and data size **DWORDS**.

Below you can see the **PeSettings** structure in its decrypted form:

```
00000000 e6 01 00 00 41 00 44 00 4d 00 49 00 4e 00 2d 00 |...A.D.M.I.N.-.|
00000010 50 00 43 00 5f 00 45 00 35 00 33 00 32 00 36 00 |P.C._.E.5.3.2.6.|
00000020 34 00 38 00 41 00 34 00 34 00 43 00 43 00 37 00 |4.8.A.4.4.C.C.7.|
00000030 46 00 31 00 43 00 00 00 00 00 00 00 00 00 00 00 |F.1.C.....|
00000040 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |.....|
*
00000070 00 00 00 00 00 00 00 00 00 00 00 00 e4 50 d2 69 |.....P.i|
00000080 18 6c e3 11 b3 bc 80 6e 6f 6e 69 63 01 89 b5 78 |.l.....nonic...x|
00000090 79 63 ae 4b f3 14 94 9a ab db c2 be 09 32 df 16 |yc.K.....2..|
000000a0 bc a3 0a 33 57 6f 49 e5 21 62 c6 5f 12 e2 97 25 |...3WoI.!b._...%|
000000b0 87 55 b7 a0 da a8 67 36 29 dc 08 f1 8a 6d c9 e8 |.U....g6)....m..|
000000c0 91 13 90 54 6b 8f 2b 5e 68 46 9b 9e 69 80 e4 76 |...Tk.^hF..i..v|
000000d0 88 85 cc bd bb 40 ce 10 6a 71 75 5d 93 dd 4d 07 |[email protected]...M.|
000000e0 92 7e ba 61 ad 1d 34 f6 ac 98 a5 af 59 86 3d 27 |.~.a..4....Y.='|
000000f0 5c 38 b6 c7 aa c0 9c 52 d0 64 77 5a 3e 8e fe 0d |\8....R.dwZ>...|
00000100 7f bf 1b 20 f8 00 a4 6c 45 3b 41 8d 81 05 e6 d4 |... ..lE;A....|
00000110 f9 e3 9f 02 37 b1 d9 60 ef 83 1f e9 cd a2 17 8c |....7..`.....|
00000120 2c c4 c1 15 65 4c d5 8b ca 3c 26 1e ec 6e 30 d8 |,...eL...<&.n0.|
00000130 a9 4a 2f 7d 18 a7 7b 56 0f f7 ea 39 1a 96 c8 4e |.J/}...{V...9...N|
00000140 73 b3 d2 f5 cb d3 74 e0 5b 51 50 eb 84 0c b4 b2 |s.....t.[QP.....|
00000150 3a ee 4f fb 58 1c 28 70 a6 43 82 66 7c 04 22 0b |:0.X.(p.C.f|".|
00000160 cf 3f f4 42 44 c5 23 47 53 19 0e 35 11 7a 95 48 |?.BD.#GS...5.z.H|
00000170 ed 2a f2 c3 99 b8 2e 06 24 ff e7 fc 9d fd d7 b0 |.*.....$.....|
00000180 b9 d6 31 e1 d1 fa f0 de a1 2d 72 03 00 00 55 76 |..1.....-r...Uv|
00000190 71 69 63 75 5c 79 70 77 75 66 2e 65 78 65 00 00 |qicu\ypwuf.exe..|
000001a0 00 00 47 61 75 6c 5c 75 6d 70 75 68 2e 62 79 67 |..Gaul\umpuh.byg|
000001b0 00 00 00 00 00 00 4f 74 68 65 79 6e 00 00 00 00 |.....Otheyn....|
000001c0 55 71 63 75 73 00 00 00 00 00 50 69 67 6f 63 6f |Uqcus.....Pigoco|
000001d0 00 00 00 00 43 61 73 75 73 61 00 00 00 00 8a 2d |....Casusa.....-|
000001e0 48 10 30 a0 77 68 15 00 00 83 |H.0.wh....|
```

When trying to run a sample that's already installed, it will generate the computer/bot ID, and if it doesn't match with the one stored in the **PeSettings**, ZeuS will exit. The same thing occurs if the install paths don't match.

Conclusion

In recent versions, Zloader has adopted a stealthy approach to system infections. This new anti-analysis technique makes Zloader even more challenging to detect and analyze. The samples analyzed by ThreatLabz have all been pre-initialized, suggesting a more targeted distribution strategy.

Zscaler ThreatLabz continues to track this threat and add detections to protect our customers.

Zscaler Coverage



SANDBOX DETAIL REPORT

Report ID (MD5): **350CECBEE06BC8FD89820D2095A0FB02**
File Type: **exe64**
Analysis Performed: **4/17/2024 1:12:40 PM**

CLASSIFICATION

Class Type
Malicious
Category
Malware & Botnet

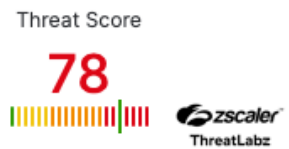


Figure 8: Zscaler Cloud Sandbox report

In addition to sandbox detections, Zscaler’s multilayered cloud security platform detects indicators related to Zloader at various levels with the following threat names:

[Win64.Downloader.Zloader](#)

Indicators Of Compromise (IOCs)

SHA256	Description
cba9578875a3e222d502bb6a85898939bb9e8e247d30fcc0d44d83a64919f448	Zloader sample
85962530c71cd31c102853d64a8829f93b63bd1406bdec537b9d8c200f8f0bcc	Zloader sample
b1a6bf93d4ee659db03e51a3765d4d3c2ee3f1b56bd9b701ab5939d63f57d9ee	Zloader sample
85b1a980eb8ced59f87cb5dd7702e15d6ca38441c4848698d140ffd37d2b55e6	Zloader sample

URLs

URL	Description
https://eingangfurkunden[.]digital/	Zloader C2
https://citscale[.]com/api.php	Zloader C2
https://adslsdfdsfmo[.]world/	Zloader C2
https://gycltda[.]cl/home/wp-api.php	Zloader C2

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