# **From AMSI to Reflection 0x0**

**rxored.github.io**[/post/csharploader/bypassing-amsi-with-csharp](https://rxored.github.io/post/csharploader/bypassing-amsi-with-csharp/)

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#### **Introduction**

In Windows environments, in both initial access and post-exploitation phases, script-based malware plays a major role. Often, hackers utilize microsoft office suite to gain initial access (using droppers, loaders) to the victim and Windows powershell to explore internal network, perform scans… basically to do the post exploitation stuff. (well of course, there are powershell based droppers.)

There is something that is common to both of these tools. Windows scripting engine.

And as a result, Microsoft and antimalware vendors have developed many security mechanisms to deal with those threats that utilize script-based malware. For example, modern anti-malware solutions can statically analyze scripts, binaries and detect whether they are malicious or not using signatures such as strings.

And because of that, malware authors use various techniques to bypass those defense mechanisms. One of the major techniques is code obfuscation.

consider the following example, that I took from MSDN.

```
function displayEvilString
{
    Write-Host 'pwnd!'
}
```
Assuming the above PowerShell snippet is malicious, we can write a signature to detect the malware. this signature can be Write-Host 'pwnd!' or simply 'pwnd!' .

So to avoid signature-based detection, the above snippet can be obfuscated like shown below.

```
function obfuscatedDisplayEvilString
{
   $xorKey = 123
   $code = "LHsJexJ7D3see1Z7M3sUewh7D3tbe1x7C3sMexV7H3tae1x7"
   $byte = [Convert]::FromBase64String($code)
   $newBytes = foreach($byte in $bytes) {
        $byte -bxor $xorKey
   }
   $newCode = [System.Text.Encoding]::Unicode.GetString($newBytes)
}
```
And this is a win for malware authors since this is beyond what anti-malware solutions can emulate or detect until AMSI joins the conversation.

### **Antimalware Scan Interface**

Antimalware Scan Interface, AMSI for short is a standard interface that allows applications to interact with anti-malware products installed on the system. This means is that it provides an API for Application developers. Application developers can use the API to implement security features to make sure that the end-user is safe.

AMSI also enables anti malware vendors to defend againts script based malware.

According to Microsoft, AMSI provides the following features by default.

- User Account Control
- PowerShell
- Windows Script Host
- JScript && VBScript
- Office VBA macros

As it is clear from those default features, AMSI specifically provides anti-malware security mechanisms to defend against script-based malware.

## **AMSI in action**

So let's take Safetykatz as our example.

When we run the binary, the result we get is.



See, as we expected, PowerShell stops the execution of the program once it has detected the program is suspicious using AMSI. So, how can we bypass this?, well before that, we have to dive deep into AMSI internals to understand how things work.

## **AMSI internals**

As I previously mentioned, AMSI enables anti malware vendors to defend againts script based malware. This is done by using AMSI providers. An AMSI provider is basically a COM object that implements IAntimalwareProvider COM interface. An anti malware vendor who's willing to implement AMSI interface should then register the COM object by creating a CLSID entry in HKLM\CLSID and registering the same CLSID under HKLM\Software\Microsoft\AMSI\Providers\ .



As it is shown in the above diagram, AMSI provides a dll called amsi.dll for application developers to interfere with AMSI providers indirectly.

Let's examine PowerShell from process hacker to check whether amsi.dll is loaded.



as we can see, amsi.dll has been loaded into powershell.exe. Now, let's take a look at this dll in-depth and see if we can find anything interesting. Even without looking at the dll, it is possible to think of some techniques to bypass AMSI, Anyway, its time to dig deep.

Before start reading disassembly, let's examine the export table of amsi.dll.



Out of the above exported functions, only two are important to us.

- AmsiInitialize
- AmsiScanBuffer
- AmsiScanString

Of course there are some other important exports. To name a few, pllReqisterClass, DllGetClassObject and AmsiUacScan .

First we'll go through AmsiScanBuffer.

### **AmsiScanString**

Microsoft documentation does not tell us much about AmsiScanString function. However it gives some basic information about it. Such as,

it's prototype,

```
HRESULT AmsiScanString(
 [in] HAMSICONTEXT amsiContext,
 [in] LPCWSTR string,
 [in] LPCWSTR contentName,
 [in, optional] HAMSISESSION amsiSession,
 [out] AMSI_RESULT *result
);
```
and parameter information.

According to the documentation, The first parameter this function accepts is amsiContext , which is a handle of type HAMSICONTEXT that was initially received from AmsiInitialize.

Second and third parameters hold pointers to wide character strings. first one for the string that should be scanned and the latter for the contentName.

contentName can be either filename, script id, url or similar of the content being scanned.

Fourth parameter is marked optional, however if multiple scan requests are to be correlated within a session, this parameter should be set to the handle returned by AmsiOpenSession function.

Fifth parameter is an output parameter and this is the one that indicates whether the input string is malicous or not.

As MSDN says, this function (and AmsiScanBuffer) returns S\_0K if the call is successful. However, the return value does not indicate whether the buffer is malicious. instead, the function uses fifth parameter of type AMSI\_RESULT to send the scan results to caller.

```
typedef enum AMSI_RESULT {
    AMSI_RESULT_CLEAN,
    AMSI_RESULT_NOT_DETECTED,
    AMSI_RESULT_BLOCKED_BY_ADMIN_START,
    AMSI_RESULT_BLOCKED_BY_ADMIN_END,
    AMSI_RESULT_DETECTED
} ;
```
Let's a take a look at AmsiScanString in disassembly.



Function allocates some space in the stack and checks if the string is empty or not. If string turns out to be empty, it simply returns after loading 0x80070057 into rax.



if string to be scanned is not null,



function checks if result is null pointer. if so, well the same thing as above, it returns with bad value loaded into rax.

else, result is valid, it loops through each wide character of the string to get the length of it.



After getting the string length, it calls AmsiScanBuffer function.

It is clear that this is just a simple wrapper function around AmsiScanBuffer.

#### **AmsiScanBuffer**

According to the MSDN and as well as the name suggests, the AmsiScanBuffer function scans a buffer for malicous content.

here is the function prototype [msdn](https://docs.microsoft.com/en-us/windows/win32/api/amsi/nf-amsi-amsiscanbuffer)



Function takes 6 parameters. One of which is the pointer to the AMSI\_RESULT enum which i explained above - \*result . According to MSDN, others include a buffer, which will be scanned by the anti-malware vendor - buffer , length of the buffer - length , filename, URL, unique script ID - contentName and a handler to the session - HAMSISESSION structure.

And here's how this function looks like in disassembly.



here we can see stack pointer is stored in  $r11$  register and since this is x64 \_stdcall, the first four parameters are stored in rcx, rdx, r8 and r9 registers. Rest are stored in the stack. With that information, we can assume a pointer to the AMSI\_RESULT enum is stored in the stack.

then we can see few comparisons around global data. if the comparisons turns out to be successful, it calls WPP\_SF\_qqDqq function. (windows sofware trace preprocessor).



then there is a pretty huge if condition, which is essentially checks if any of the above parameters are invalid



by looking at the comparison, the function won't successfully return if **[rbp]**, which is the first qword of amsiContext is not equal to 0x49534d41.



And if parameters invalid, it returns 0x80070057 (which i think is the bad return value)



else, as we can see in the above snippet, buffer (rdx register) is now loaded with address of CAmsiBufferStream::vftable and stored the value in the stack. This may sound familiar to anyone who has done some C++ reverse engineering since this is a one way to represent constructor calls in assembly (setting vtable to the object's first bytes).

to confirm that we can take a look at CAmsiBufferStream::vftable.



as we can see, CAmsiBufferStream::vftable is indeed, a virtual function table and what those two instructions doing is creating an object of type CAmsiBufferStream. It is also possible to see some member variable intializations too.

My assumption is that amsiContext->thirdMember is somekind of a class that antimalware vendor has registered to perform scans.

To make sure our assumptions so far are correct, we'll go over this function using windbg.

Since we already know interesting parts of the function, it is easy to place breakpoints.



First few breakpoints are placed at locations in assembly where **amsiContext's** member variables are being referenced. Reason being this handle is still unknown to us. Therefore it could be useful to extract every possible information about it. Last breakpoint is placed at the address where **CAmsiBufferStream:vftable** is referenced.



So from the above image, we can assume that the first member of the amsiContext is a QWORD but it compares it with a DWORD and second and third members are also QWORDs (8 bytes).

0:018> dq /c1 0x000002347f5d44d8 L1 000002347f5d44d8 000002347e90cce0 0:018> dq /c1 0x000002347f5d44e0 L1 000002347f5d44e0 000002347eb5d120

We can refer to the memory map to get more information about what those QWORDs are.



Now it is clear those two pointers are from heap segment 1. However, we still have no idea about the type of those pointers.

However we already know those are pointers to objects thanks to our previous static analysis.



Above screenshot shows the virtual function table of CAmsiBufferStream.

Then the next address where we can find some more information regarding **amsiContext members** is,

00007ff9455033d6 488b01 mov rax, qword ptr [rcx] ds:000002347eb5d120= {amsi!ATL::CComObject<CAmsiAntimalware>::vftable' (00007ff94550bb48)} 00007ff9455033d9 488b4018 mov rax, qword ptr [rax+18h] 00007ff9455033dd ff15cd8d0000 call qword ptr [amsi!\_guard\_dispatch\_icall\_fptr (00007ff9`4550c1b0)]

in the above snippet, rcx holds one of those pointers we just discussed,

000002347eb5d120 (thirdMember). In the first instruction, 64 bit value at that address is loaded into rax register, which, according to the above snippet, is 00007ff94550bb48. It also specifies that this is a vtable located in .rodata section of the asmi.dll's memory image.



next two instructions retreives address **0x18** offset from the vtable into rax register and calls the address stored in rax



This proves that our assumption on function pointer extracted from the HAMSICONTEXT being a anti-malware vendor's registered function is false and it is a pointer to amsi!CAmsiAntimalware::Scan method.

We have uncovered some important details about HAMSICONETXT so far. We already know that the first member is a DWORD, and it should be equal to **0x49534d41** in order for scan to be successful. Third member is a pointer to an object of class CAmsiAntimalware , which has a virtual function called amsi!CAmsiAntimalware::Scan.

And by moving its 0x0 offset rax register, we can access it's virtual function table where we can find Scan at the 0x18.

The whole thing can be roughly decompiled down into below C code.

```
class CAmsiAntimalware {
        private:
            [...]
        public:
           virtual Scan(CAmsiBufferStream *, AMSI_RESULT, DWORD);
            [...]
    }
    typedef HAMSICONTEXT {
        QWORD unk1;
        QWORD *secondMember;
        CAmsiAntimalware *antimalware;
       [...]
    };
    HRESULT __stdcall AmsiScanBuffer
    (
           HAMSICONTEXT amsiContext,
           PVOID buffer,
           ULONG length,
           LPCWSTR contentName,
           HAMSISESSION amsiSession,
           AMSI_RESULT *result
    )
    {
       auto var;
       if ((WPP_GLOBAL_Control != &WPP_GLOBAL_Control) && (*(WPP_GLOBAL_Control +
0x1c)) != 4))
        {
           WPP_SF_qqDqq(
                *((BYTE*)WPP_GLOBAL_Control + 0x10),
               buffer,
               length,
               amsiContext,
               buffer,
               amsiSession,
               result
           );
        }
        if (
               buffer == NULL ||
               result == NULL ||
               amsiContext == NULL ||
                *((DWORD *)amsiContext) != 0x49534D41 ||
                *((QWORD *)amsiContext + 1) == 0x0 ||
               *((QWORD *)amsiContext+2) == 0x0
            )
```

```
{
        return 0x80070057;
    }
    else
    {
        CAmsiBufferStream bufferStream = CAmsiBufferStream(
            buffer,
            length,
            amsiContext->secondMember,
            contentName,
            session
        );
        return amsiContext->antimalware->Scan(
            amsiContext->antimalware, // this
            &bufferStream, // CAmsiBufferStream *
            result,
            \Theta);
    }
}
```
We are not done yet. Goal here is to understand how AMSI works. Therefore, our next target is amsi!CAmsiAntimalware::Scan.

But before drill down into it, we need to construct the HAMSICONTEXT structure out of the knowlegde we have.



now we can see decompiler output is much more accurate and readable.



We can also try constructing a CAmsiAntimalware class but we dont have enough information to populate member variables.

#### **CAmsiAntimalware::Scan**



So ghidra has created a nice view of the stack frame for us. And by looking at the parameters, we see the function expects a pointer to an IAmsiBuffer object and a pointer to a pointer of IAntimalwareProvider object.

We saw that in the AmsiScanBuffer that this value is set to zero.



Then continues to setup all those memory curruption protection machanisms and to check the validity of the input parameters. First it checks if third parameter, result is null (remember, result is a pointer to AMSI\_RESULT enum).



if it is not, it jumps to label result\_valid . else, it sets eax to 0x80070057 and return. In the result\_valid label, it sets \*result to AMSI\_RESULT\_CLEAN (0x0). So it looks like the function is clearing the \*result to not detected state. Which means we can expect value of result to change.

It also checks if provider is null. If not, it sets value of it to null and continue execution from LAB\_7ff94550565c . else, it continues the execution from the same location but without setting \*provider to null.



LAB\_7ff94550565c does the same thing as AmsiScanBuffer did at the block 0x7ffxxxxx335d . However instead of calling WPP\_SF\_qqDqq it calls WPP\_SF\_q . Also note that above snippet sets rdx to either address of [WPP\_GLOBAL\_CONTROL] or 0x1e.

LAB\_7ff94550568d looks interesting.



First it calls rand() function. In case you dont know, it's pretty common C library function and it generates a psuedo random number and return it. In the next line, it stores a member of CAmsiAntimalware class at offset 0x1c0 in r13 register. Then there are some multipications around the generated value value.

ghidra being ghidra, has renamed registers with the variable names (this is good if we are doing x86 reversing becuase most of calling conventions pass parameters through stack, However, in our case, since parameters are passed through registers, renaming those can cause confusion), So to make it clear, we'll use listing view.



It assigns the return value from rand() to ecx register and loads eax with **0x51eb851f**. then it multiplies random value stored in ecx with the value loaded in eax . Note that this instruction is capable of changing the value at edx register.

Then there's a shift right instruction, which shifts 5 bits from edx register, then it multiplies shifted edx with 0x64 and stores the value in eax .

sub instruction substracts  $\cos x$ , by  $\cos x$ . what this whole thing does is similar to below expression

```
rand() % 0x64;
```
value of ecx is then stored in a local variable loc\_rand and function checks if r13, which holds the value of this- $>0 \times 100$  is 0/null. If yes, it jumps to LAB 7ff9455058c4. else, it continues exection from next address.

Now we got two control paths to follow. but first, I'm not gonna take the jump.

#### **Control flow path 1**



0x7ffxxxxx56bb , address of this->0x40 gets loaded into r14 , which then gets stored in a local variable. Next instruction loads this->0xc0 into r12 register.

Then there's an unconditional jump and this one jumps directly into a loop. so Im gonna save that part for a debugging session and continue with the other control flow path.

#### **Control flow path 2**



LAB\_7ff9455058c4 starts with a comparison of  $r13$  (this->0x1c0 but as a local variable) with this->0x1c0. The comparison checks if r13 is less than this->0x1c0. if it is, control flow is directed to address 0x7ffxxxxx58cd . else, control flow is directed to label LAB 7ff9455058f7.

First instruction at  $\theta$ x7ffxxxxx58cd sets r14 to zero (rbx is xored by itself at the begining of the function). Next two instructions checks if  $r12$  is null.



if not, value at address r12 is set to  $[RSI + r13*0x8 + 0x40]$ . Then it checks if rcx is null. If we assume the jump to LAB\_7ff9455058c4 taken from 0x7ffxxxxx56b5, then rcx would be the remainder of rand() % 0x64 thing. if rcx is null, jump is taken to label LAB 7ff9455058fd . else, it loads value at  $({*}(rcx) + 0x8)$  to rax and calls it through \_guard\_dispatch\_icall .

if r12 is null, jump is also taken to label LAB\_7ff9455058fd .



on the other hand, LAB\_7ff9455058f7 also jumps to LAB\_7ff9455058fd after moving 0x1 into [rdi] . We already know that rdi is pointing to AMSI\_RESULT enum. Constant 1 means AMSI\_RESULT\_NOT\_DETECTED .



this simply checks if this->0x1c0 is null, if it is, it jumps to label LAB\_7ff94550590e else, it continues exection from address  $0x7f$ fxxxxx5906.

block starting at  $0 \times 7$  f f x x x x 5906 basically checks if R14 is null. it sets b1 if previous comparison has caused sign flag to be 1. The operation may look like this in pseudocode.

 $b1 = (r14 < 0) + 1;$ 

as you can see in the above control flow graph, code is finally directed towards

LAB\_7ff94550590e . What this snippet does is, call

CAmsiAntimalware::GenerateEtwEvent method. it passes this and amsiStream and bl through rcx, rdx and r9 registers as first three arguments. fourth and the last one is passed through r9 and this is basically the AMSI\_RESULT .

Now Im going to find where AMSI\_RESULT is being modified. We already know rdi is a pointer to the enum.



In the above snippet, rdi (result) is assigned to value of eax if we go up in the control flow, we can see eax is assigned with local\_108 .

Now we know some interesting places to place breakpoints and analyze, it is time to get into a windbg session.

First, Im gonna place a break point at address at place where provider is checked.

```
0:018> bp 0x7ffxxxxx5654
0.018> g
[...]
0.018> r r9
r9=0000000000000000
```
As it is clear from the above snippet, rg register which holds a pointer to a pointer of IAntimalwareProvider class is set to zero. We saw this earlier in AmsiScanBuffer function.

Even if some value is passed down through this register, CAmsiAntimalware::Scan will set it to zero.

the next important piece for us is where this is being accessed.



above diagram shows exection has been stopped just after the instruction where function accessess this->0x1c0 .

And the value at that address is set to 0x1. This gives us a hint that this member might be numerical value rather than a pointer.



A little below that, we can the random number generated by rand() being stored in ecx register and that value is 0x2ea6 .

Since we already know what this snippet does, we can perform the calculation by ourself.

```
>>> hex(0x2ea6 % 0x64)
'0x2a'
```


#### Above diagram conludes that.



Above diagram shows where the function retreives address of this->0x40 into r14 register.

When this- $>0 \times 40$  is printed, it also looks like an address that pointed at heap.

Value at \*this->0x40 looks like a function pointer and when disasseble that address, windbg prints disassembly of MpOav!DllRegisterServer (another dll ? we'll see)but disassembly starts from the middle of the function. This might not be a function pointer after all.



here is another place where a member of CAmsiAntimalware class has been referenced. this time as we've discussed when doing static analysis, stores address this->0xc0.

It doesnt provide us with imformation about type of data even if we take a look at the data at that address,

#### **Control flow path 1 continued**

Now we are at the instruction in disassembly where that loop begins.



We see that in the above image, first instruction loads address of  $rsp+0x48$  into rcx register and calls GetSystemTimePreciseAsFileTime , which is used to retrieve the current system date and time with the highest possible level of precision in UTC format.

before the call instruction it also initialize rsp+0x40 and rsp+0x48 with 0x0.

Then value at address r14 gets stored in rcx register. if you remember, r14 register stores &this->0x40 so rcx would be value of this->0x40 .

Then can see some manipulations around that value.



mov rax, qword ptr [rcx] stores value at \*this->0x40 in rax register. Next instruction takes 0x18 th offset of it and stores it back in rax register. Then that address is called using a gaurd\_dispatch\_icall\_fptr .

With that information it is clear that  $\frac{t}{t}$  this- $>0 \times 40$  is a pointer to an object of an unknown class. rcx now points to that object and rax holds one of function pointers in the object's vftable. Well my guess is that this is the windows defender's AMSI COM interface.

The first argument passed to the function is this->0x40. Second, third and fourth are passed through rdx and r8 registers. we can see that in the disassembly rdx being set to rsp+0x70 (amsiBuffer) and r8 being initialized to the address of rsp +0x40 (who's value is 0).

Weird thing is, the function is jumping into the middle of a function.

Let's try following it.



Well this makes it bit clear. First of all we not jumping into the middle of a function, See that ret instruction up there? What this tells us is, we jumped into a function but it is not labelled correctly.

However if you try to goto this address from a disassembler, it will fail. Indicating that this a function from another dll.

here's the memory map.



See? It seems like this dll is the COM dll that implements IAmsiAntimalware interface for windows defender.

To confirm that, let's check the registry.

// registry

Now it is confirmed, let's go through this function.



First it does some work on the stack frame and moves 0x80070057 to rax register if third parameter is null (pointer to a stack variable of CAmsiAntimalware::Scan method), And we know this is E\_INVALIDARG . And then function jumps to the epilogue. So this is basically a small sanity check.



then it moves 1 or AMSI\_RESULT\_NOT\_DETECTED into third parameter and checks if first parameter (rcx) + 200 is 0. We know that first parameter (rcx) passed down to this function is CAmsiAntimalware->0x40 . (yes doesnt make much sense.)



In our case, comparison turns out to be true.

A little below that, there's a call to another fuction from this dll.



it seems to take only one argument and it is &rcx+0x70 .



if we step into it, windbg indentifies function as RtlEnterCriticalSection from ntdll. According to [msdn](https://docs.microsoft.com/en-us/windows/win32/api/synchapi/nf-synchapi-entercriticalsection), EnterCriticalState function waits for ownership of the specified critical section object. The function returns when the calling thread is granted ownership. function accepts a single parameter and it is of LPCRITICAL\_SECTION .

In this case, critical section that this function waits for is  $rcx+0x70$ .



next few instructions compare rsi+0x98 with 0 (both rsi and rcx pointed to same address but since rcx now points to rcx+0x70, rsi is used). if comparison fails, it jumps to another location disassembly where LeaveCriticalState is being called.





as shown in the above diagram, function loads rbp, which points to the critical section (rsi->0x70) into rcx . Then LeaveCriticalState function is called.

then two local variables, rsp+0x54 and rsp+0x50, get initialized to 0x0 and 0x1, following a mov instruction which loads a global variable into rcx . then it does a comparison of rcx+0x60 with 0.

In our case, comparision fails and for that reason, jump will be taken.

```
MpOav!DllRegisterServer+0x1198:
00007fffae7b3928 488b4138 mov rax, qword ptr [rcx+38h]
ds:00000250a357c158=00000250a356b1f0
00007fffae7b392c 4c8d442450 lea r8, [rsp+50h]
00007fffae7b3931 498bd7 mov rdx, r15
00007fffae7b3934 488b4948 mov rcx, qword ptr [rcx+48h]
```
here we can see another call.

PID: 11152 - WinDbg 1.2111.9001.0 (Administrator

rcx is set to  $\lceil$ rcx+0x48] and rdx is loaded with amsiBuffer meanwhile r8, third argument is loaded with address rsp+0x50 .



as we can see in the above diagram, this call is to MPCLIENT! MpAmsiScan function. This is basically a function exported by windows defender's MPCLIENT.dll. So this means we have reached our destination.



Let's step over this function and inspect the return value since it is out of scope of this article to reverse engineer windows defender internals.



According the above diagram, the return value we get is 0x0. And there's no way to determine whether this is a indication of detection or not because windows documentation does not provide imformation about MpAmsiScan

Therefore we have to try some tricky methods to identify it.

First, im going to continue the exection.

as expected the result is,



Then we can place a breakpoint at the address where MpAmsiScan return and send some non-malicous input.

Weirdly enough, return value is same. So this function must be using an output parameter to pass the result of the scan, just like AmsiScanBuffer.

Can you remember that the third parameter to MpAmsiScan is a pointer to a local variable? Just in case, keep it's address in mind.

Somewhere down below, before the program generates an event saying safetykatz is malicious, return value or output parameter of MpAmsiScan must be accessed in order determine whether it's detected by windows defender or not.

Back to where we left off,

return value of MpAmsiScan is stored in edi register and function compares it with 0 after moving some value to rcx register.

00007fffae7b3945 8bf8 mov edi, eax 00007fffae7b3947 488b0d32f90300 mov rcx, qword ptr [MpOav!DllRegisterServer+0x40af0 (00007fffae7f3280)] MpOav!DllRegisterServer+0x11be: 00007fffae7b394e 85ff test edi, edi 00007fffae7b3950 7925 jns MpOav!DllRegisterServer+0x11e7 (00007fffae7b3977) [br=1]

if return value (edi) is greater than or equal to zero,



it sets value of third parameter (pointed by r14) to 1 and simply returns. Also note that return value is set to edi .

else if return value of MpAmsiScan (edi) is less than 0,



it checks validity of some data and calls a function and then returns after setting return value to that of MpAmsiScan stored in edi register, just like the previous one.



Because the return value we got from MpAmsiScan is 0x0, execution path will be the first one we've discussed above.

There is something interesting that we havent discussed about that control flow path. There is a comparison of rsp+0x54 and 1. if that comparison is able to set zero flag, next instruction sets al register to 1.

in our case, rsp+0x54 is not equal to 1.

0:018> dd @rsp+0x54 L1 00000015`8864e564 00000000

which means, al wont be set to 1. If you can remember,  $rsp+0\times54$  is only accessed once, just after the call to LeaveCriticalState and that that is the only instruction that sets rsp+0x54 to 0x0. My guess is that this checks if function has entered the

LeaveCriticalSection block. It then sets  $[s]$   $[rsi+0C8h]$   $(rsi == first parameter)$  to the value of al. Note that rsi+0xc8 should be set to zero in order for this function to be sucessful. We discussed rest of this block earlier.

after the function returns, we'll end up back at CAmsiAntimalware:: Scan. Good news is, we dont need to read every instruction since we already know what we are looking for.



Above image shows how the call looks in decompiled pseudo code. return value of the callee is stored in local variable uVar2 . However, we know this is not accurate because caller need to pass three args to the callee (we see none). That's not important to us though.



Here, the if confition only evaluate true when loc\_rand() is equal to zero and a global variable is less than 5. loc\_rand is basically the local variable where the random number was stored. Therefore this block is not going to execute.



Above if condition checks if return value (stored in r14) is zero. In our case it is. we know that the third argument passed to the collee is the address of  $rsp+0\times40$  and was passed through r8 .

below image shows disassembly of the above snippet



As shown above, mov r8d, dword ptr[rdi] moves value at address stored in rdi into r8 register. rdi stores the address of AMSI\_RESULT enum passed down to CAmsiAntimalware::Scan method. it then moves rsp+0x40 , output paramater we discussed earlier into eax register.



comparison instruction and jump instruction checks if value in r8 (result) is greater than that of in eax (output parameter). jump wont be taken and execution will directed to the next mov instruction.

This is basically checking if current scan's result is greater than that of previous one.



In the above snippet it loads eax into  $\lceil$  rdi  $\rceil$ , and value of r15 into r13 and compare some global variables related to WPP .

According to the decompiled snippet, this checks some global variables related to WPP tracer and if checks are valid, it jumps to a location in disassembly after setting rdx register to the address  $r14 + 1f$ . Well this has nothing to do with addresses eventhough the instruction is lea . r14 is 0x0. therefore what this does is, it loads 0x1f into rdx register.

However, if we step through each instruction, cmp rcx, rdx will evaluate to oxo and the jump will be taken.



in the above snippet, dword value at address stored in rdi is compared to hex 0x8000, decimal 32768. Aand this is exactly the same value [msdn](https://docs.microsoft.com/en-us/windows/win32/api/amsi/ne-amsi-amsi_result) specifies in their documentation for AMSI\_RESULT enum. quoting msdn,

'Any return result equal to or larger than 32768 is considered malware, and the content

should be blocked. An app should use AmsiResultIsMalware to determine if this is the case.'

next instruction is a *jge* and it essentially takes the jump if dword at address stored in rdi (AMSI\_RESULT) is greater than or equal to 0x8000. if it is, it breaks from the loop.

In our case, value at address stored in rdi is less than 0x8000 so the jump won't be taken. Instead control flow will be redirected to



r15 is incremented by 1 and it is then compared to  $this \rightarrow 0 \times 100$ , whose value is 1. if r15 is below that value, it will jump to the address where the loop begins.

Possibly, the loop is going through every registered anti-malware vendor's COM interface. Since I dont have any anti malware services installed in the VM, its going to loop only once. This also uncovers some details about CAmsiAntimalware class members. The loop terminates after loop iterator veriable being compared to this->0x1c0 . Therefore this- >0x1c0 is the value that indicates number of registered anti malware services or AMSI providers.



Now the question is, we just executed a malicous program and it just got flagged as AMSI\_RESULT\_NOT\_DETECTED . But we still see powershell produces that red ugly output saying that it detected a malicous program.

And suprisingly, there's no call to AmsiResultIsMalware.



First if condition checks if r13 register is less than the number of providers (this->0x1c0). We saw that r15, which acts as the counter loaded into r13 previously. What this is checking is that if anything malicous detected before going through all the providers.

Now it is time to conclude our assumptions on AmsiInitialize.

### **AmsiInitialize**

### **The End**

So yeah that's it for now… we explored AMSI in-depth in this article. In the next one, We will go through some common AMSI bypass techniques.

#Spread Anarchy!