Knockin' on Heaven's Gate – Dynamic Processor Mode Switching

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George Nicolaou

Abstract

This post presents the research conducted under the domain of dynamic processor mode (or context) switching that takes place prior to the invocation of kernel mode functions in 32bit processes running under a 64bit Windows kernel. Processes that are designed and compiled to execute under a 32bit environment get loaded inside the Windows-on-Windows64 (WoW64) subsystem and are assigned threads running in IA-32e compatibility mode (32bit mode). When a kernel request is being made through the standard WoW64 libraries, at some point, the thread switches to 64bit mode, the request is executed, the thread switches back to compatibility mode and execution is passed back to the caller.

The switch from 32bit compatibility mode to 64bit mode is made through a *specific segment call gate* referred to as the *Heaven's Gate*, thus the title of this topic. All threads executing under the WoW64 environment can execute a FAR CALL through this segment gate and switch to the 64bit mode.

The feature of mode switch can also be viewed from the security and maliciousness point of view. It can be used as an anti reverse engineering technique for protecting software up to the malicious (or not) intends of cross process generic library injection or antivirus and sandbox evasion. The result of this research is a library named W64oWoW64 which stands for Windows64 On Windows On Windows64.

Introduction

Within the WoW64 environment, threads that wish to switch between compatibility mode (32bit mode) to 64bit mode, in order to request the invocation of kernel mode functions, have to go through the *Heaven Gate* located at code segment selector *0×0033* that identifies the *call gate* inside the GDT. The process of context switching occurs multiple times throughout the lifespan of a WoW64 process and is essential for their compatibility with the Windows 64bit kernel. However, this feature creates a number of minor security issues or inconsistencies to security or software analysis products. Over the next paragraphs, we will explore the methodology used by the operating system when user mode applications engage in the invocation of kernel functions as well as the differences between the two modes from the perspective of a thread. Next we shall explore how context switching can be used by 32bit malicious processes to communicate, control and inject libraries in 64bit mode applications using a library created just for this purpose.

Research Laboratory

This research was conducted on a Windows 7 64bit Operating System with all updates and patches installed as of (see post date). The tools used are:

Tracing to Heaven

Before we get our hands dirty, we need to briefly dive into the Windows mechanisms of calling the heaven gate. To do this we need to understand how the gate is used, the reasons, as well as what happens when we switch to 64bit mode.

To begin with, we shall trace a call to ZwTestAlert. You can do this by loading any 32bit application on a 64bit Windows operating system and issuing the following command on the first *LdrpDoDebuggerBreak* breakpoint.

bp ntdll32!ZwTestAlert

Note that the breakpoint is set for ntdll32 which is the 32bit WoW64 version of the ntdll library. If by any chance you've set a breakpoint using the *bp ZwTestAlert* command then you'd be setting it to the 64bit ntdll version of the library. Before we go any further let us check the modules currently loaded in memory. For this example the POC executable *HeavenInjector.exe* was used and the lm (list modules) command was executed with the result illustrated in *Figure 1*.

0:000>1m				
start	end	module name		
	000000000`00400000 000000000`00405000	HeavenInjector (private pdb symb		
	000000000`74f80000 000000000`74f88000	wow64cpu	(pdb svmbols)	\mathbf{r}
	000000000'74f90000000000000'74fec0000	wow64win	(pdb symbols)	$\mathcal{L}_{\mathcal{A}}$
	00000000`74ff0000 00000000`7502f000	wow64	(pdb symbols)	\cdot
	000000000`77790000 000000000`77939000	ntdll	(pdb symbols)	$\overline{}$
	00000000`77970000 00000000`77af0000	ntd1132	(pdb symbols)	\blacksquare

Figure 1: Default Loaded Modules

As you can see there are two versions of the ntdll library, one being the 32bit one and the other the 64bit one. Along side the executable we can list three other libraries which depend on the processor architecture currently being used by your system. It is worth noting that the test OS is running on AMD64.

Execute the program using the g command or the *F5* key until you reach the breakpoint we've just set. You can view the disassembly code of the ZwTestAlert function by hitting *u* (which disassembles 8 instructions from the current address or 9 instructions if your platform runs on an Itanium processor.) [1] or by bringing up the Windbg's Disassembly window from View > Disassembly. The code is shown in *Figure 2*.

$0:000:x86>$ q Breakpoint 0 hit		
ntdl132!ZwTestAlert:		
77991da0 b87e010000	mov	eax,17Eh
$0:000:\times86>u$		
ntdll32!ZwTestAlert:		
77991da0 b87e010000	mov	eax, 17Eh
77991da5 b902000000	mov	ecx.2
77991daa 8d542404	lea	edx, [esp+4]
77991dae 64ff15c0000000	call	dword ptr fs:[0C0h]
77991db5 83c404	add	$\texttt{esp.4}$
77991db8 c3	ret	
77991db9 8d4900	lea	ecx, [ecx]
ntdll32!ZwThawRegistry:		
77991dbc b87f010000	mov	eax,17Fh

Figure 2: ZwTestAlert Disassembled Code

On a 32bit Windows 7 version the above function looks slightly different (See *Figure 3*).

$0:000$ a ZwTestAlert		
ntdll!ZwTestAlert:		
770d5d30 b874010000	$_{\text{mov}}$	eax.174h
770d5d35 ba0003fe7f	mov	edx.offset SharedUserData!SvstemCallStub
770d5d3a ff12	call	dword ptr [edx]
770d5d3c c3	ret	
770d5d3d 8d4900	lea	ecx, [ecx]

Figure 3: ZwTestAlert 32bit Disassembled Code

An obvious difference between the two Operating Systems is the SYSENTER instruction located at *SharedUserData!SystemCallStub* which is not present in the WoW64 ZwTestAlert function (Figure 2). That instruction is replaced with a CALL instruction to a pointer located in *fs:[0C0h]*. On Windows 32bit processes the *fs* segment holds the address of the TEB for the current thread and the 0C0h value signifies the offset from that address to the value that is being read. To view the current TEB address we need to issue the *!wow64exts.info* command as shown in Figure 4 below. Note that the Windbg pseudo register *\$teb* holds the address of the 64bit TEB for this thread.

> $0:000:x86> 1$ wow64exts.info PEB32: 0x7efde000 PEB64: 0x7efdf000 Wow64 information for current thread: TEB32: 0x7efdd000 TEB64: 0x7efdb000 $: 0x190000$ 32 bit, StackBase StackLimit : 0x18e000 Deallocation: 0x90000 $: 0 \times 8f d20$ 64 bit, StackBase StackLimit : 0x8c000 Deallocation: 0x50000 Wow64 TLS slots: $0x000000000008ec10$ WOW64_TLS_STACKPTR64: WOW64_TLS_CPURESERVED: 0x000000000008fd20 WOW64 TES INCRUSINULATION:
WOW64 TLS INCRUSINULATION:
WOW64 TLS EXCEPTIONADDR:
WOW64 TLS USERCALLBACKDATA:
WOW64 TLS EXTENDED FLOAT: 0x0000000000000000 0x0000000000000000 0x0000000000000000 0x00000000000000000 0x0000000000000000 WOW64_TLS_APCLIST:
WOW64_TLS_RPCLIST:
WOW64_TLS_FILESYSREDIR:
WOW64_TLS_WOW64INFO:
WOW64_TLS_WOW64INFO: 0x00000000000000000 0x0000000000000000 0x00000000000000000 0x000000007efde248

Figure 4: WoW64 Information

The value TEB32 contains the address of this thread's TEB. We note that value and issue the *dt* command to dump the TEB structure, along with any values, into the command window. To do this we execute the following command:

dt teb 7efdd000

Where *_teb* is the symbol of the 32bit TEB and *7efdd000* is the address of the 32bit TEB. The resulting output should be similar to the one shown in *Figure 5*.

 $\begin{array}{l} 0:000:x86 > \mathrm{dt} _\mathrm{teb} \ \text{7efdd} 000 \\ \mathrm{HeavenInjector} \text{!}_\mathrm{TEB} \end{array}$ vening

+0x000 NtTib = NT_TIB

+0x01c EnvironmentPointer : (null

+0x020 ClientId : _CLIENT $(null)$: _CLIEN
: (null)
agePointer CLIENT_ID +0x028 ActiveRpcHandle +0x02c ThreadLocalStoragePointer 0x7efdd02c Void +0x030 ProcessEnvironmentBlock : 0x7efde000 _PEB +0x034 LastErrorValue D. +0x038 CountOfOwnedCriticalSections : 0 +0x03c CsrClientThread \therefore (null) +0x040 Win32ThreadInfo $(null)$ $\begin{bmatrix} 26 \\ 0 \\ 15 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$ +0x044 User32Reserved +0x0ac UserReserved 0x74f82320 Void +0x0c0 WOW32Reserved +0x0c4 CurrentLocale 0x809 +0x0c8 FpSoftwareStatusRegister : 0 +0x0cc SystemReserved1 : [54] (null) +0x1a4 ExceptionCode $_{0n0}$ +0x1a8 ActivationContextStack : _ACTIVATION_CONTEXT_STACK +0x1bc SpareBytes1 [24] ""
_GDI_TEB_BATCH
_CLIENT_ID +0x1d4 GdiTebBatch +0x6b4 RealClientId +0x6bc GdiCachedProcessHandle : (null) +0x6c0 GdiClientPID $\ddot{}$ -n +0x6c4 GdiClientTID n +Ox6c8 GdiThreadLocalInfo : (null) Figure 5: TEB32 Of Main Thread

As you can see the offset *+0x0c0* points the the *WOW32Reserved* field which contains the address *0x74f82320*. All WoW64 calls to the kernel are being redirected to this address. If we disassemble any other ntdll32 functions such as *ZwOpenProcess*, *NtLoadDriver*, etc we can see that the same CALL instruction with the same address is called.

Continuing the execution of the program and tracing into the *call dword ptr fs:[0C0h]* instruction by hitting *F11* or typing *t* into the command window we end up at the address pointed to by the *WOW32Reserved* field which lands inside the *wow64cpu* library at function X86SwitchTo64BitMode as shown in *Figure 6.*

> $0:000: x86 > t$ wow64cpu!X86SwitchTo64BitMode: 0033:74F8271E 74f82320 ea1e27f8743300 jmp

> > Figure 6: wow64cpu!X86SwitchTo64BitMode

The above instruction jumps to the given address of the code segment through a specified segment selector call gate. Intel's specification ^[2] refers to this instruction as a FAR Jump instruction which if it's segment selector (in this case 0×0033) is a call gate then then the code jumps to the code segment specified in the call gate descriptor (which is located in the GDT) and executes the code pointed to by the gate, if the segment selector is for a code segment then a far jump to the segment is performed. which in this case handles the switch from 32bit to 64bit.

When we trace the JMP instruction we end up being in 64bit mode at the address pointed to by the instruction. The address contains the entry point of the *wow64cpu*!CpupReturnFromSimulatedCode function (as shown in Figure 7) which in short, sets up the environment for the current system call and executes the *SYSCALL* instruction. Once finished, all results are normalized for 32bit mode and the function returns back to the initial 32bit system call shown in *Figure 2*.

wow64cpu!CpupReturnFromSimulatedCode:		
00000000'74b2271e 67448b0424	mov	r8d, dword ptr [esp] ds:00000000
000000000`74b22723 458985bc000000	M ₀	dword ptr [r13+0BCh], r8d
000000000`74b2272a 4189a5c8000000	mov	dword ptr [r13+0C8h].esp
00000000'74b22731 498ba42480140000 mov		$rsp, qword ptr [r12+1480h]$
00000000'74b22739 4983a4248014000000 and		qword ptr [r12+1480h].0
00000000`74b22742_448bda	$m \cap V$	r11d.edx
wow64cpu!TurboDispatchJumpAddressStart:		
00000000'74b22745 41ff24cf	imp	qword ptr [r15+rcx*8]
wow64cpu!ServiceNoTurbo:		
000000000`74b22749 4189b5a4000000	$_{\text{mov}}$	dword ptr [r13+0A4h], esi
000000000`74b22750 4189bda00000000	mov	dword ptr [r13+0A0h], edi
00000000`74b22757 41899da8000000	mov	dword ptr [r13+0A8h], ebx
000000000`74b2275e 4189adb8000000	M ₀	dword ptr [r13+0B8h],ebp
00000000`74Ъ22765 9∈	pushfq	
AAAAAAAAN 7.4699766		

Figure 7: WoW64cpu!CpupReturnFromSimulatedCode Entry

For the purposes of this research a short assembly algorithm was devised to understand the effects of a FAR CALL instruction through the heaven gate. The algorithm is shown below:

When this algorithm was executed the value pushed by the far CALL instruction within the stack revealed an additional segment selector call gate 0×0023 (which is actually the 32bit code segment we just came from) who's purpose is to switch from the current 64bit mode to the compatibility 32bit mode. *Figure 8.1* Illustrates the top of the stack right after the CALL. As you can see the last four (4) bytes 0x004011c3 contain the return address whereas the preceding two (2) bytes 0×0023 contain the segment selector eall gate number.

In conclusion, the process of switching modes is required for the communication between the WoW64 processes and the windows kernel. *Figure 8.2* below illustrates the process discussed in the above paragraphs.

Figure 8.1: Heaven Gate After-CALL Return Address

WoW64 ZwTestAlert Call x64 Switch Illustrated

After-Switch Environment

Before we begin abusing the heaven gate, we need to understand the post-switch environment of the thread including which libraries are loaded and how we can reconstruct it in such a way allowing us to execute any 64bit compiled code or libraries.

To begin with, we need to locate all 64bit libraries loaded along side the executable. We can identify them using the *lm* command as shown in *Figure 1* then using the *!dh* command in conjunction with the address of a library to dump it's headers. *Figure 9* illustrates this process for a single library *wow64*.

 $0:000>1m$ module name $\begin{array}{ll} \texttt{start} & \texttt{end} \\ 000000000^\circ\ 004000000^\circ\ 000000000^\circ\ 00405000 \\ 000000000^\circ\ 73310000^\circ\ 000000000^\circ\ 733c\texttt{f}000 \\ 000000000^\circ\ 74\texttt{h}20000^\circ\ 000000000^\circ\ 74\texttt{h}28000 \\ 000000000^\circ\ 74\texttt{h}30000^\circ\ 00000000$ start end HeavenInjector (p: MSVCR100 (pdb sym) wow64cpu (pdb sym) wow64win (pdb sym) $w \cap w \in \mathcal{A}$ (pdb sym kernel32 (pdb sym (pdb s KERNELBASE (pdb sym) ntdll $ntd1132$ (pdb sym) $0:000$ > !dh 00000000'74b90000 File Type: DLL FILE HEADER VALUES 8664 machine (X64) 5 number of sections 4E212272 time date stamp Sat Jul 16 08:32:34 2011 0 file pointer to symbol table 0 number of symbols F0 size of optional header 2022 characteristics Executable App can handle >2gb addresses DĪĪ Figure 9: Retrieving wow64 Headers

Switching from 32bit to 64bit does not cause any other libraries to be loaded, therefore we reach to the conclusion that only the following libraries are accessible and loaded for the 64bit mode of this process. We can verify this by first dumping the 64bit PEB structure using the following command:

dt ntdll! peb $@$ \$peb -r

Next we locate the Ldr.InLoadOrderModuleList.Flink address and issue a *!list* command listing all libraries currently loaded for this process. *Figure 10* shows the PEB structure.

To issue the !list command we need the address of the first InLoadOrderModuleList entry which is the Flink entry located at address *0×00000000`005f3400* in the above figure. Next we issue the following command to dump the linked entries in the InLoadOrderModuleList chain:

!list -x "dt _LDR_DATA_TABLE_ENTRY" 0×00000000`005f3400

The resulting output should list the libraries we are seeking. Therefore, when switching to 64bit mode the current process environment is running with the following:

- 1. In 64bit Mode
- 2. Has the following libraries loaded ntdll.dll, wow64.dll, wow64win.dll, wow64cpu.dll
- 3. Has separate TEB and PEB structures than the 32bit process

Unfortunately, our initial goal to execute any code or libraries within that environment lacks one key element. When a process is loaded, the loader first loads the ntdll.dll library and right after that the kernel32.dll library. In our case kernel32.dll is never loaded within the 64bit environment. Therefore, we need to load the library using the *LdrLoadDll* function located within ntdll.

Issue 1: Aligning the stack for 64bit mode

Another issue that might come up with the execution of certain functions is the issue of maintaining the stack alignment between modes. When switching to 64bit mode the stack register ESP, or in this case RSP, retains it's original value aligned on a 32bit boundary. In order to overcome this issue all we have to do is "waste" enough bytes to align the stack for 64bit execution then realign it before switching back to 32bit mode.

Issue 2: Identifying and calling ntdll API functions

After crossing the heaven's gate, the environment we come across is no different than the unknown environment of a simple shellcode environment. This means that our code has no prior knowledge of function pointers or environment variables. In order to overcome this issue we would have to walk the PEB table, identify the address of the ntdll library and then locate the necessary functions for the successful execution of our payload.

Issue 3: Loading Kernel32.dll – Understanding The Constraints and Protections

Any attempts to load kernel32.dll using the *LdrLoadDll* function would result to the error code *0xC0000018* (STATUS_CONFLICTING_ADDRESSES). This is due to the fact that the default memory location of kernel32 is already mapped as private. Therefore, when *LdrpFindOrMapDll* attempts to map the section of the image using *LdrpMapViewOfSection* a process of walking the VAD tree is initialized resulting to a conflicting address between the library's preferred base and a privately allocated page at the same address. That page is located at the original kernel32.dll base address and is placed there to prevent loading the library from a WoW64 environment. *LdrpMapViewOfSection* ends up loading the library at a different base and returns *STATUS_IMAGE_NOT_AT_BASE*. This triggers an algorithm within *LdrpFindOrMapDll* function that ends up comparing the library string provided by our call to *LdrLoadDll* with the string located at *ntdll!LdrpKernel32DllName*, which contains the unicode string "kernel32.dll". As a side note, it is worth mentioning that the exact same processes occurs when loading the user32.dll library. The algorithm's purpose is to identify whether the system library kernel32.dll has not been loaded at it's preferred base address and if so unload it and return the conflicting addresses error.

In order to solve this issue, one could employ a simple hooking technique to redirect execution from the string comparison function RtlEqualUnicodeString to a stub function that would force *RtlEqualUnicodeString* to return a negative answer which would in turn result to the OS loading the kernel32.dll library at any base address. This however is not a complete solution since certain functions contained within ntdll require numerous structures from the library that are referenced using their absolute address. In addition the kernel32 library's initialization function *KernelBaseDllInitialize* (which is also the EP of the library) would fail to execute and raise an unhandled exception in the process. Therefore, loading kernel32.dll at any base address except the one specified by the operating system is a bad idea.

Loading kernel32 at its original base address requires an understanding of the methodology used to load a 32bit executable within the WoW64 environment. It is essential for us to identify the protections placed by the loader so that they can be overcome.

When running the 64bit version of WinDbg the first breakpoint you come across is hit by the 64bit *ntdll!LdrpDoDebuggerBreak* function prior to the invocation of any wow64 processing. If you hit *F5* or type *g* in the command line you view an output similar to the one shown in *Figure 11*.

> $0:000>g$ WOW64_IMAGE_SECTION
WOW64_IMAGE_SECTION
NOT_AN_IMAGE
NOT_AN_IMAGE
C:\Windows\syswow64\kernel32.dll
C:\Windows\syswow64\KERNELBASE.dll
C:\Windows\syswow64\USER32.dll
C:\Windows\syswow64\GDI32.dll Modload: 00000000'75650000 00000000'756e0000
Modload: 00000000'775-0000 00000000'756-0000 C:\Windows\syswow64\GDI32.dll

Figure 11: WoW64 Initialization

If you compare the addresses of the *NOT_AN_IMAGE* and the first *WOW_IMAGE_SECTION* with the modules loaded in a 64bit application such as calc.exe (as shown in *Figure 12*) you will immediately identify that those locations are actually the system wide base addresses for the *kernel32.dll* and *user32.dll* libraries. However, if you execute the lm command at the 32bit executable's EP those modules are no longer registered within the loader data table entry inside the PEB.

ModLoad: 000000000`fff40000 00000001`00023000	calc.exe
ModLoad: 000000000`773f0000 00000000'77599000	ntdll.dll
ModLoad: 000000000`770c0000 00000000'771df000	C:\Windows\system32\kernel32.dll
ModLoad: 000007fe`fdb70000 000007fe`fdbdc000	C:\Windows\svstem32\KERNELBASE.dll
ModLoad: 000007fe`fdfb0000 000007fe`fed38000	C:\Windows\svstem32\SHELL32.dll
ModLoad: 000007fe`feec0000 000007fe`fef5f000	C:\Windows\system32\msvcrt.dll
ModLoad: 000007fe'ff350000 000007fe'ff3c1000	C:\Windows\system32\SHLWAPI.dll
ModLoad: 000007fe`ff240000 000007fe`ff2a7000	C:\Windows\system32\GDI32.dll
ModLoad: 00000000'76d10000 00000000'76e0a000	C:\Windows\system32\USER32.dll
ModLoad: 000007fe`feeb0000 000007fe`feebe000	C:\Windows\svstem32\LPK.dll

Figure 12: calc.exe Initial Loaded Modules

In order to identify how these pages are allocated and assigned at the loader table we need to trace the execution following up the first breakpoint (in 64bit *ntdll!LdrpDoDebuggerBreak*) we come across. Therefore, we hit *F10* or p until we reach the CALL instruction pointing to *ntdll!Wow64LdrpInitialize* as shown in *Figure 13*.

Figure 13: Call to ntdll!Wow64LdrpInitialize (wow64!Wow64LdrpInitialize)

This function, located within wow64.dll, is responsible for initializing the 32bit Wow64 subsystem such as the 32bit ntdll.dll, calling the function that initializes filesystem redirections and so on. Since we are only interested at the initialization of the page, we trace through its code until we reach the CALL to *wow64!ProcessInit* as shown inFigure 14.

	00000000`74c5c1a7 0f8503010000 00000000`74c5c1ad 488d0d3ce80200	ine lea	wow64!Wow64LdrpInitialize+0x190 (000000 rcx, [wow64!CpuThreadSize (00000000'74c{
	00000000`74c5c1b4_e867cdffff	call	wow64! $\texttt{ProcessInit}$ (00000000'74c58f20)
00000000`74c5c1Ъ9 8Ъd8		$m \cap V$	ebx.eax
	00000000`74c5c1bb ff153750ffff	call	gword ptr [wow64! imp LdrProcessInitial
000000000`74c5c1c1 413bdd		CMD	ebx.r13d

Figure 14: Call to wow64!ProcessInit

This function's responsibility, amongst other, is to load the debug wow64log.dll (which does not exist on production systems) for debugging purposes, initialize filesystem redirection and most importantly, make our work slightly more difficult by mapping the addresses of the libraries we wish to load. The *wow64!InitializeContextMapper* function call (shown in Figure 15) is responsible for mapping the firstWOW64 IMAGE_SECTION for kernel32.dll and looking up the export table of the library (which falls outside the scope of this research).

000000000 74c5915d 4533c9 r9d.r9d xor AAAAAAAAAN AILEEN CONCADARA ndx $f(x,y)$ $70k1$ $1 - -$	00000000'74c59145 4885c9 00000000 74c59148 7406 00000000'74c5914a ff15507fffff 00000000`74c59150 e89b020000 000000000 74c59155 85c0 00000000`74c59157 0f88a5000000		test пe call call test 1s.	rox.rox wow64!ProcessInit+0x230 (00000 qword ptr [wow64!_imp_LdrUnloa wow64!InitializeContextMapper eax.eax wow64!ProcessInit+0x2e2 (00000
--	---	--	---	---

Figure 15: Call to wow64!InitializeContextMapper

We Trace over *wow64!InitializeContextMapper* until the point where we come across wow64!Map64BitDlls (as shown in *Figure 16*) who's purpose is to setup the environment is such a way thus denying the mapping of our libraries to their original system-wide default base address.

00000000'74c591c5 8905ad1d0300	$_{\text{mov}}$	dword ptr [wow64!Ntdll32KiUserCallbackDisp
000000000'74c591cb8b04255003fe7f	mov	eax.dword ptr [SharedUserData+0x350 (00000
00000000'74c591d2 8905801d0300	M ₀	$dword$ ptr [wow64!Ntdll6432Map+0x8 (0000000)
000000000074c591d8 e8570a0000	call	$www64!Map64BitD11s (00000000'74c59c34)$
00000000`74c591dd 85c0	test	eax.eax
000000000`74c591df 7821	i js l	wow64!ProcessInit+0x2e2 (00000000'74c59202

Figure 16: Call to wow64!Map64BitDlls

Tracing through that function we hit the first interesting CALL instruction which points to *ntdll!LdrGetKnownDllSectionHandle* (as shown in *Figure 17*). Unfortunately, looking up this function in your favorite search engine does not produce any significant results (at least to the eyes of the author). Therefore, the function prototype and purpose is explained in the followup paragraph.

$0:000$ > p			
wow64!Map64BitDlls+0x57: 00000000074c59c8b 488b4d00	mov		rcx, qword ptr [rbp] ss:00000000'74c8a1c8={wow64!'s
$0:000$ b			
wow64!Map64BitDlls+0x5b:			
00000000`74c59c8f 4c8d442460	lea	$r8$, [r sp+60h]	
$0:000$ > p			
wow64!Map64BitDlls+0x60:			
00000000 74c59c94 33d2 $0:000$ b	xor	edx, edx	
wow64!Map64BitDlls+0x62:			
00000000'74c59c96 ff152474ffff			
$0:000$ de rex	call		qword ptr [wow64! imp LdrGetKnownDllSectionHandle
00000000`74c53ef0 0065006b 006e0072 006c0065 00320033			k.e.r.n.e. 1.3.2.
00000000 74c53f00 0064002e 006c006c 00000000 00000000			d.1.1
00000000 74c53f10 00730025 00730025 00000000 00000000			$\%$. S . $\%$. S
00000000 74c53f20 74696e49		696c6169 6f43657a 7865746e	InitializeContex
00000000 74c53f30 70614d74 3a726570 69616620 2064656c			tMapper: failed
00000000 74c53f40 20657564 66206f74 6e656c69 20656d61			due to filename
00000000 74c53f50 676e656c 65206874 65656378 676e6964			length exceeding
00000000'74c53f60 534f4420 58414d5f 5441505f 454c5f48			DOS MAX PATH LE

Figure 17: Call to ntdll!LdrGetKnownDllSectionHandle

LdrGetKnownDllSectionHandle receives three arguments, first a pointer to the variable which holds the Unicode name of the library (or section name), a Boolean flag which dictates which directory handle should be used when loading the section. TRUE for the Wow64 directory containing the 32bit versions of the library requested and FALSE for the 64bit version directory. Finally the last argument is a pointer to a variable which would receive the handle of the section. Note that this function is essentially a wrapper for the *NtOpenSection* routine. The function's prototype is defined below:

```
NTSTATUS LdrGetKnownDllSectionHandle(
LPCSTR lpwzLibraryName,
BOOL bIs32BitSection,
HANDLE * lphSection
);
```
Given what we know now and what is illustrated at *Figure 17* we can safely deduct that the *wow64!Map64BitDll* function retrieves the handle of the section used to map kernel32 throughout all processes. Next, that handle is used to call the *ntdl!!NtMapViewOfSection* ^[3] function as shown in *Figure 18*. Note that at the same figure, at address 00000000`74c59ded the pointer to a unicode string which reads "NOT_AN_IMAGE" is moved to the quad word pointed to by *r13+28h*, where at this specific location *r13* holds the address of the 64bit TEB and offset *+28* contains the *ArbitraryUserPointer* within the TIB structure.

00000000'74c59db7 498b5d28		mov	rbx, qword ptr [r13+28h]
	00000000`74c59dbb 488b4c2460	mov	rex.gword ptr [rsp+60h]
	00000000 74c59dc0 488364245000	and	gword ptr [rsp+50h].0
	00000000`74c59dc6_488364245800	and	qword ptr [rsp+58h],0
	00000000'74c59dcc c744244801000000 mov		dword ptr [rsp+48h], 1
	00000000`74c59dd4 8364244000	and	dword ptr [rsp+40h],0
	00000000`74c59dd9 c744243802000000	$_{\text{mov}}$	dword ptr [rsp+38h], 2
	00000000`74c59de1 488d0530a5ffff	lea	rax, [wow64!'string' (00000000'74c54318)]
	00000000`74c59de8 4c8d442450	lea	$r8$, [$rsp+50h1$]
00000000`74c59ded 49894528		mov	gword ptr [r13+28h], rax
	00000000'74c59df1 488d442458	lea	rax, [rsp+58h]
00000000`74c59df6 4533c9		xor	r9d.r9d
	00000000`74c59df9 4889442430	mov	gword ptr [rsp+30h],rax
	00000000`74c59dfe 488364242800	and	qword ptr [rsp+28h],0
	00000000`74c59e04 488364242000	and	qword ptr [rsp+20h],0
00000000'74c59e0a 4883caff		or.	rdx.0FFFFFFFFFFFFFFFFFh
	AAAAAAAA`74a59⊖A⊖ ff15aa72ffff	call	oward ntr [waw64] imm NtManViewOfSection.

Figure 18: Call to ntdll!NtMapViewOfSection

So far we have deducted that wow64 loads the section of kernel32 into memory at it's correct base address (also verified by the return value of *NtMapViewOfSection*) and named after the string put in *ArbitraryUserPointer* prior to the invocation of NtMapViewOfSection. Immediately, and right after the call to the mapping function, the section's handle is closed and the freshly mapped section is unmapped using the *NtUnmapViewOfSection* function. The reason for mapping and unmapping the section becomes obvious over the next few instructions which are the cause of our original problem with kernel32. The proceeding function call to NtAllocateVirtualMemory ^[4] (as shown in Figure 19) receives as arguments, you've guessed it, a pointer to the original base of kernel32.dll, an AllocationType of MEM_RESERVE and Protect of PAGE_EXECUTE_READWRITE.

00000000'74c59e4b4c8d4c2458	lea	r9.frsp+58hl
00000000`74c59e50 488d542450	lea	rdx, [rsp+50h]
$000000000'74c59e55$ $4533c0$	xor	r8d.r8d
$000000000'74c59e58$ 488bcb	mov	rox.rbx
$000000000'74c59e5b c744242840000000 m$ ov		dword ptr [rsp+28h], 40h
$000000000'74c59e63c744242000200000m$ ov		dword ptr [rsp+20h],2000h
000000000`74c59e6b ff159772ffff call		oward ntr [waw64] imm NtAllocateVirtualMemory

Figure 19: Call to ntdll!NtAllocateVirtualMemory

Next, the function re-iterates using "user32.dll" as known section handle and proceeds with executing the same algorithm of allocating the memory page. Once finished, the function returns towow64!ProcessInit and the initialization process continues.

The above paragraphs have walked us through the process of protecting the pages where crucial libraries are supposed to be loaded at. However, since the protection placed is just a memory allocation, it can be overcome by simply freeing that memory page. \bigcirc

Constructing The Payload

Over the next few paragraphs, we shall describe the payload stub's code which is responsible for overcoming the issues identified above. Henceforth, whenever we refer to the *payload*, we refer to the piece of 64bit code that executes right after passing through the heaven's gate.

Solving Issue 1

Before we begin doing any library loading we need to align the stack on a 64bit boundary. This means that the stack needs to be aligned on an 8byte boundary from it's base. For example, if the stack is allocated at address *0×00000000'00100000* then all proceeding elements should be referenced at multiples of 8 resulting to a stack which resembles the following structure:

Additionally, any modifications to the stack need to be backed up so the stack can be realigned back to it's original 32bit boundary. The solution is rather simple, since the 64bit alignment allows only the first 61 bits of the stack pointer value to be set then the last 3 bits have to be discarded. In order to understand this, let us have a look at some binary values along with their hexadecimal representations as well as what happens when we take away the last three bits of that value:

As you can see when taking away the last 3 bits of each value, it becomes 64bit aligned. In order to code this in assembly all we require is the AND and SUB instructions as follows:

Then at the end of the function we add the subtracted bytes from the local variable *bStackAlignment* to realign the stack pointer back to it's initial value as shown below:

However, for the purposes of the W64oWoW64 library the lower two (2) bytes of the ESP register are ANDed with the value 0xFFF8 (since the Visual Studio MASM compiler doesn't like referencing the lower byte of the stack register SPL).

Solving Issue 2

In order to prepare the environment for the execution of arbitrary code or library we need to retrieve a number of API functions located in the 64bit loaded ntdll library.

- **LdrLoadDll** In order to load kernel32.dll and a payload library compiled for x64.
- **LdrGetKnownDllSectionHandle** Which will be used to retrieve the section handle of kernel32.dll and user32.dll in order to retrieve their original base address.
- **NtFreeVirtualMemory** Which will be used to free the original library base address memory page that was allocated by *wow64!Map64BitDlls*.
- **NtMapViewOfSection** Which is used to map the section retrieved by *LdrGetKnownDllSectionHandle* at a random base address so we can retrieve the library's original base address from the PE Header.
- **NtUnmapViewOfSection** To unmap and clean up the memory after the original base address of the library has been retrieved.

In addition to that, the addresses of the following libraries are required:

- **ntdll.dll Base Address** Which can be retrieved from the PEB.
- **kernel32.dll Base Address** Which is returned by the *LdrLoadDll* call or can be accessed through the PEB.

For the purpose of retrieving the base address of the ntdll library, a function named *GetModuleBase64* was devised and implemented. You can find this function in *w64wow64.c* which is attached to this post. It receives the library name and returns base address of the library. Implementation details for this function can be found in the source code. For the purposes of this post all you need to know is that the function retrieves the PEB of the current thread and walks the *Ldr.InLoadOrderModuleList* chain to retrieve the library base.

Additionally, retrieving function pointers from libraries is made possible with the use of another function in the same file named GetProcAddress64. This function receives the module base as first argument and the API function name as its second argument. The function walks through the PE header of the provided module, loads up the *IMAGE_EXPORT_DIRECTORY* and identifies the function.

The following code within the *InitializeW64WoW64()* function is self explanatory, its purpose is to resolve the address of ntdll.dll and the first required

functions LdrGetKnownDllSectionHandle, *NtFreeVirtualMemory*, *NtMapViewOfSection*, *NtUnmapViewOfSection*.

```
void * lvpNtdll = GetModuleBase64( L"ntdll.dll" );
UNICODE_STRING64 sUnicodeString;
__int8 * lvpKernelBaseBase;
__int8 * lvpKernel32Base;
PLDR_DATA_TABLE_ENTRY64 lpsKernel32Ldr;
PLDR_DATA_TABLE_ENTRY64 lpsKernelBaseLdr;
sFunctions.LdrGetKnownDllSectionHandle = GetProcAddress64( lvpNtdll,
        "LdrGetKnownDllSectionHandle" );
sFunctions.NtFreeVirtualMemory = GetProcAddress64( lvpNtdll,
        "NtFreeVirtualMemory" );
sFunctions.NtMapViewOfSection = GetProcAddress64( lvpNtdll,
        "NtMapViewOfSection" );
sFunctions.NtUnmapViewOfSection = GetProcAddress64( lvpNtdll,
        "NtUnmapViewOfSection" );
```
Solving Issue 3

The next issue we come across, is the issue of properly loading the 64bit kernel32.dll library into the address space of the process. One solution is to patch theRtlEqualUnicodeString function to return false when it's two arguments are equal to "kernel32.dll". However, care must be taken to uninstall the hook right after kernel32.dll is loaded since the next time a library is loaded it would reload it, resulting in a rather weirdly looking address space with more than two kernel32.dll libraries loaded.

Our approach is much simpler and requires us to free the memory location of kernel32 and user32 library default base addresses then load them independently using the *LdrLoadDll* function. In order to do that, the function *FreeKnownDllPage*() was constructed which receives the section name Unicode string (the one used in *LdrGetKnownDllSectionHandle*) and frees up the memory page by loading the section, walking through the PE header to get the original base address and executesNtFreeVirtualMemory to free up the memory location. Following is the prototype of this function

BOOL FreeKnownDllPage(wchar_t * lpczKnownDllName)

This function is called with the following arguments within the *InitializeW64WoW64()* init function as shown below:

if(FreeKnownDllPage(L"kernel32.dll") == FALSE) return FALSE; if(FreeKnownDllPage(L"user32.dll") == FALSE) return FALSE;

The function implementation is given below:

```
BOOL FreeKnownDllPage( wchar t * lpwzKnownDllName )
{
        DWORD64 hSection = 0;
        DWORD64 lvpBaseAddress = 0;
        DWORD64 lvpRealBaseAddress = 0;
        DWORD64 stViewSize = 0;
        DWORD64 stRegionSize = 0;
        PTEB64 psTeb;
        /*
        ** X64Call of WOW64Ext Library - http://blog.rewolf.pl/
        ** (Copyright (c) 2012 ReWolf)
        */
        X64Call( sFunctions.LdrGetKnownDllSectionHandle, 3,
                (DWORD64)lpwzKnownDllName,
                (DWORD64)0,
                (DWORD64)&hSection );
        psTeb = NtTeb64();
        psTeb->NtTib.ArbitraryUserPointer = (DWORD64)lpwzKnownDllName;
        X64Call( sFunctions.NtMapViewOfSection, 10,
                (DWORD64)hSection,
                (DWORD64)-1,
                (DWORD64)&lvpBaseAddress,
                (DWORD64)0,
                (DWORD64)0,
                (DWORD64)0,
                (DWORD64)&stViewSize,
                (DWORD64)ViewUnmap,
                (DWORD64)0,
                (DWORD64)PAGE_READONLY );
        lvpRealBaseAddress =
                (DWORD64)GetModule64PEBaseAddress( (void *)lvpBaseAddress );
        if( X64Call( sFunctions.NtFreeVirtualMemory, 4,
                (DWORD64)-1,
                (DWORD64)&lvpRealBaseAddress,
                (DWORD64)&stRegionSize,
                (DWORD64)MEM RELEASE ) != NULL ) {
                        PrintLastError(); //XXX doesnt work
                        return FALSE;
        }
        X64Call( sFunctions.NtUnmapViewOfSection, 2, (DWORD64)-1,
                (DWORD64)lvpBaseAddress );
        return TRUE;
}
```
For now, all you need to know is that this function calls *LdrGetKnownDllSectionHandle* with the following arguments:

- 1. **lpwzLibraryName -** The name of the known section contained withing lpwzKnownDllName.
- 2. **bIs32BitSection** False, since we are loading the 64bit version of the library.
- 3. **lphSection** A pointer to a local variable which shall receive the section handle

Next, once successfully executed the function calls *NtMapViewOfSection* to map the library at a random base address (since the original is already allocated) with the following arguments:

- 1. **SectionHandle** The handle contained withing the *hSection* local variable that was just retrieved.
- 2. **ProcessHandle** Current process handle which is equal to -1.
- 3. **BaseAddress** A pointer to a local variable *lvpBaseAddress* which will receive the base address this section will be loaded at.
- 4. **ZeroBits** Not required and set to 0.
- 5. **CommitSize** Not required since it is already set.
- 6. **SectionOffset** Not required and set to 0.
- 7. **ViewSize** A pointer to the local variable *stViewSize* which is set to 0 and will receive the section size.
- 8. **InheritDisposition** Set to *ViewUnmap* since we don't plan on creating any child processes.
- 9. **AllocationType** Not used and set to 0.
- 10. **Win32Protect** Protection set to *PAGE_READONLY* since we only wish to read from it.

Once executed, the *NtMapViewOfSection* function will place the base address of the newly loaded section in *lvpBaseAddress* local variable which is then used as an argument to *GetModule64PEBaseAddress()* function within *w64wow64.c* to retrieve the BaseAddress field of the PE Header's optional header. This address can then be fed to NtFreeVirtualMemory ^[5] with the following arguments to free up the memory page:

- 1. **ProcessHandle** Current process -1.
- 2. **BaseAddress** Pointer to the real base address of the module retrieved through the *GetModule64PEBaseAddress* function.
- 3. **RegionSize** A pointer to a local variable which is set to 0.
- 4. **FreeType** We wish to free up the memory therefore *MEM_RELEASE* is provided.

Finally, the *NtUnmapViewOfSection* is called in order to unmap the section that was just loaded just for the sake of keeping the memory clean.

Once *FreeKnownDllPage* finishes executing, all you have to do now is load kernel32 using *LdrLoadDll* which will load it at its original base. The code for that is shown below:

```
sUnicodeString.Length = 0x18;
sUnicodeString.MaximumLength = 0x1a;
sUnicodeString.Buffer = (DWORD64)L"kernel32.dll";
if( X64Call( GetProcAddress64( lvpNtdll, "LdrLoadDll" ), 4,
        (DWORD64)0,
        (DWORD64)0,
        (DWORD64)&sUnicodeString,
        (DWORD64)&lvpKernel32Base ) != NULL ) {
                PrintLastError();
                return FALSE;
}
```
Once kernel32.dll and it's static dependency KERNELBASE.dll are loaded, we need to call their initialization functions *Dllmain* located at their EP. To do that all we have to do is retrieve the EntryPoint field from their PE Header and call the Dllmain function with the standard arguments and the *DLL_PROCESS_ATTACH* flag as shown below:

```
lvpKernelBaseBase = (int8 *)GetModulBeBase64( L"KERNELBASE.dll");X64Call( ( lvpKernelBaseBase + (int)GetModule64EntryRVA( lvpKernelBaseBase ) ),
        3,
        (DWORD64)lvpKernelBaseBase,
        (DWORD64)DLL_PROCESS_ATTACH,
        (DWORD64)0 );
X64Call( ( lvpKernel32Base + (int)GetModule64EntryRVA( lvpKernel32Base ) ),
        3,
        (DWORD64)lvpKernel32Base,
        (DWORD64)DLL_PROCESS_ATTACH,
        (DWORD64)0 );
```
Finally, once the libraries are loaded and in order to make the functional, there is one small detail that needs to be taken care of. Each library's Ldr data table entry contains two fields that are modified by the loader right after a library is loaded. However, in the case of kernel32 and KERNELBASE those are not set. The fields which we are referring to are the *LoadCount* field which needs to be set to -1 in order to lock in the library, and the *Flags* field which needs to have the *LDRP_ENTRY_PROCESSED* and

LDRP_PROCESS_ATTACH_CALLED flags set. To do that, we make use of the *GetModule64LdrTable*() function within *w64wow64.c* which receives a Unicode string of the library name (the one matched in the *BaseDllName* entry within the *LDR_DATA_TABLE_ENTRY* structure) and returns a pointer to the data table entry. Next, all we have to do is apply the mentioned modifications as shown below:

```
lpsKernel32Ldr = GetModule64LdrTable( L"kernel32.dll" );
lpsKernel32Ldr->LoadCount = 0xffff;
lpsKernel32Ldr->Flags += LDRP_ENTRY_PROCESSED | LDRP_PROCESS_ATTACH_CALLED;
lpsKernelBaseLdr = GetModule64LdrTable( L"KERNELBASE.dll" );
lpsKernelBaseLdr->LoadCount = 0xffff;
lpsKernelBaseLdr->Flags += LDRP_ENTRY_PROCESSED | LDRP_PROCESS_ATTACH_CALLED;
```
This concludes the solution to the problem with loading and initializing kernel32.dll within the 64bit mode of a wow64 application. From this point on (given that any other initializations are made) you can execute any kind of code you see fit.

Loading an External 64bit Payload DLL (HeavenInjector)

Finally, once kernel32.dll gets loaded, the environment is ready to accommodate any external libraries which can be loaded using the LoadLibrary function. For the purposes of this POC code a payload library has been coded which makes use of the *CreateRemoteThread* API function to inject a library in a 64bit application. The payload library name is *payload.dll* and is included, along with it's source code in the file attached to this post.

The payload code loads this library using *LoadLibrary* and calls the exported function *InjectLibrary.*

Finally, the POC as a whole is consisted of the following files:

- heaveninject.exe 32bit Executable which receives as arguments a 64bit process id and the library pathname to inject into that process. It switches to 64bit using the *w64wow64* library, loads payload.dll using the exported function *LoadLibrary64A* from within *w64wow64* and executes the *InjectLibrary* function who's pointer is retrieved using *GetProcAddress64* again from within the *w64wow64* library.
- payload.dll A 64bit library responsible for injecting a library into the 64bit process.
- a.dll A POC Hello World library which is injected into the process.

Final Notes

The provided proof of concept code is experimental and might require some additional coding to support some external system libraries that might not be initialized properly.

Conclusions

In conclusion to this rather enormous post, we note that the techniques described above can be used or abused as an anti-reverse engineering technique on 32bit applications rendering the code executed in 64bit inaccessible by a 32bit debugger such as Ollydbg. Additionally, this technique can also have devastating results on usermode sandbox or hooking technologies that might install hooks on 32bit system or application libraries. Thank you for reading :).

Downloads

W64oWoW64 Library:<https://github.com/georgenicolaou/W64oWoW64>

HeavenInjector: <https://github.com/georgenicolaou/HeavenInjector>

POC Video for HeavenInjector: http://www.youtube.com/watch?v=Z1c_OrW7VaQ