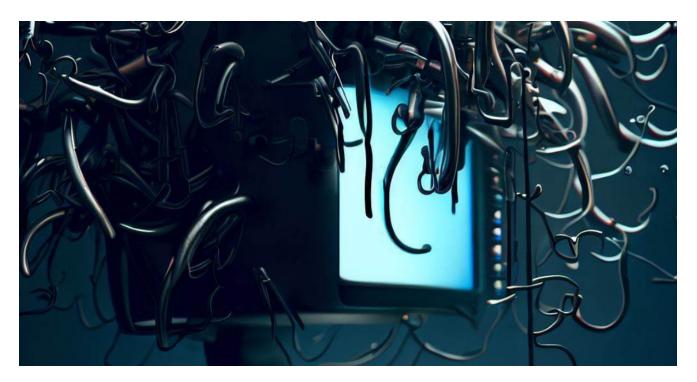
# Understanding Syscalls: Direct, Indirect, and Cobalt Strike Implementation

d01a.github.io/syscalls

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## Syscalls? Why?

- To Bypass user-mood hooks. why?
  - For Hiding a code inside a legitimate process (Process Injection)
  - Avoiding EDR alerts!

### **User-mood Hooks**

Hooking user-mode functions by placing a jump to another code section. EDRs use hooks to check the function parameters. For example, if you are trying to change the memory protections of some data to add executable protections. This is a very suspicious activity so EDRs will be alert to that. Most Hooks are on the lowest level of the user-mode interface in **ntdll.dll** which are the system calls.

# **Direct syscalls**

Windows has a defined schema of how syscalls are used. Most of the documented windows APIs are just a wrapper of a lower-level Functions in ntdll.dll which are compiled to a syscall with the right SSN (System Service Number). To look at how Nt\* version of the higher-level API is implemented.

0:018> uf NtOpenProcess ntdll!NtOpenProcess: 00007ffa`4874d4c0 4c8bd1 00007ffa`4874d4c3 b826000000 00007ffa`4874d4c8 f604250803fe7f0 (00000000`7ffe0308)],1	mov mov 1 test	r10,rcx eax,26h byte ptr [SharedUserData+0x308
00007ffa`4874d4d0 7503	jne	ntdll!NtOpenProcess+0x15
(00007ffa`4874d4d5) Branch		
ntdll!NtOpenProcess+0x12:		
00007ffa`4874d4d2 0f05	syscall	
00007ffa`4874d4d4 c3	ret	
ntdll!NtOpenProcess+0x15: 00007ffa`4874d4d5 cd2e	int	2Eh
00007ffa`4874d4d7 c3	ret	

At address 00007ffa~4874d4d2 there syscall instruction. This instruction transfers the execution to the system-handler at the kernel. The handler is specified using pre-defined SSN number loaded into EAX Register (In this case EAX =  $0 \times 26$  at address

00007ffa~4874d4c3). So, to make a syscall The SSN associated. The code stub of the syscalls is simple.

mov r10, rcx
mov eax,
<syscall\_number>
syscall
ret

Now, the missing thing is the syscall\_number. These numbers are changing based on the Build version of windows. There are some techniques to get these numbers.

### 1. SysWhispers

<u>SysWhispers</u> That generate the table of these numbers in the form of a header file and assembly file that can be embedded in the code. The generated code contains syscall number for multiple versions, The right windows build version is detected at runtime using PEB structure.

```
+0x118 OSMajorVersion :
Uint4B
+0x11c OSMinorVersion :
Uint4B
+0x120 OSBuildNumber :
Uint2B
....
```

The assembly code generated (Full document at <u>example-output</u>)

```
NtOpenProcess PROC
        mov rax, gs:[60h]
                                                ; Load PEB into RAX.
NtOpenProcess_Check_X_X_XXXX:
                                            ; Check major version.
        cmp dword ptr [rax+118h], 5
        je NtOpenProcess_SystemCall_5_X_XXXX
        cmp dword ptr [rax+118h], 6
        je NtOpenProcess_Check_6_X_XXXX
        cmp dword ptr [rax+118h], 10
        je NtOpenProcess_Check_10_0_XXXX
        jmp NtOpenProcess_SystemCall_Unknown
NtOpenProcess_Check_6_X_XXXX:
                                            ; Check minor version for Windows
Vista/7/8.
        cmp dword ptr [rax+11ch], 0
```

```
je NtOpenProcess_Check_6_0_XXXX
        cmp dword ptr [rax+11ch], 1
        je NtOpenProcess_Check_6_1_XXXX
        cmp dword ptr [rax+11ch], 2
        je NtOpenProcess_SystemCall_6_2_XXXX
        cmp dword ptr [rax+11ch], 2
        je NtOpenProcess_SystemCall_6_3_XXXX
        jmp NtOpenProcess_SystemCall_Unknown
NtOpenProcess_Check_6_0_XXXX:
                                           ; Check build number for Windows
Vista.
        cmp dword ptr [rax+120h], 6000
        je NtOpenProcess_SystemCall_6_0_6000
        cmp dword ptr [rax+120h], 6001
        je NtOpenProcess_SystemCall_6_0_6001
        cmp dword ptr [rax+120h], 6002
        je NtOpenProcess_SystemCall_6_0_6002
        jmp NtOpenProcess_SystemCall_Unknown
                                            ; Check build number for Windows 7.
NtOpenProcess_Check_6_1_XXXX:
        cmp dword ptr [rax+120h], 7600
        je NtOpenProcess_SystemCall_6_1_7600
        cmp dword ptr [rax+120h], 7601
        je NtOpenProcess_SystemCall_6_1_7601
        jmp NtOpenProcess_SystemCall_Unknown
                                            ; Check build number for Windows 10.
NtOpenProcess_Check_10_0_XXXX:
        cmp dword ptr [rax+120h], 10240
        je NtOpenProcess_SystemCall_10_0_10240
        cmp dword ptr [rax+120h], 10586
        je NtOpenProcess_SystemCall_10_0_10586
        . . .
```

- 1. **SSN code stub** This technique doesn't Look for SSN number, instead it gets the code stub of the required API. This can be done by opening the PE file and parsing the Export table of ntdll
- 2. **Extract SSN** It Extract the SSN from ntdll by parsing the Export table. The difference between it and the previous one is that it only extracts the syscall number. Both methods load ntdll.dll from the disk first using win32 API OpenFile which might be hooked. <u>hell's gate</u> for more.

- 3. Syscalls' number sequence This method take advantage of the SSNs are in a sequence for example if a syscall number is 0x26 the following will be 0x27 and so on. This relies also on the fact that not all the system calls are hooked! So, to get the SSN of a function, you need to find the nearest unhooked syscall. this was presented by <u>halos gate</u>. But This is not valid in newer versions of Windows as the SSNs sequence is no longer valid.
- 4. Parallel loading This is an interesting technique explained in this <u>blog</u>. It uses windows feature introduced in windows 10 to load DLLs through multiple threads instead of one in older versions of windows. It was found that the syscall stub of native Functions NtOpenFile(), NtCreateSection(), ZwQueryAttributeFile(), ZwOpenSection() and Z wMapViewOfFile() -There is a lot of things happens between the two actions, detailed explanation in the previously mentioned <u>blog</u> -are copied into LdrpThunkSignature array. This is done to check the integrity of the functions' code. These APIs' syscall numbers can be used to load a new version of ntdll.dll from the disk and avoid any user-mood hooks.
- 5. Sorting by system call address This technique uses the relation between the address of the system call stub and the SSN. It is known as <u>FreshyCalls</u>. In simple words, it walks the Export Address Table of ntdll and saves the Name -or a hash of the name- and Address of each entry in a table. Then, it sorts the entries by the addresses in ascending order. It was found that the first function NtAccessCheck (by address) has an SSN = 0

0:007> uf NtAccessCheck ntdll!NtAccessCheck: 00007ffa`4874d000 4c8bd1 mov r10,rcx 00007ffa`4874d003 b80000000 mov eax,0 00007ffa`4874d008 f604250803fe7f01 test byte ptr [SharedUserData+0x308 (00000000`7ffe0308)],1 00007ffa`4874d010 7503 jne ntdll!NtAccessCheck+0x15 (00007ffa`4874d015) Branch ntdll!NtAccessCheck+0x12: 00007ffa`4874d012 0f05 syscall 00007ffa`4874d014 c3 ret ntdll!NtAccessCheck+0x15: 00007ffa`4874d015 cd2e int 2Eh 00007ffa`4874d017 c3 ret

and if we unassembled the next function by adding one to the last address (as ret opcode is one byte) we will get that the next function's SSN is 1!

0:007> uf 00007ffa`4874d017+1 ntdll!NtAccessCheck+0x18: 00007ffa`4874d018 0f1f84000000000 nop dword ptr [rax+rax] 00007ffa`4874d020 4c8bd1 mov r10,rcx 00007ffa`4874d023 b801000000 eax,1 mov 00007ffa`4874d028 f604250803fe7f01 test byte ptr [SharedUserData+0x308 (00000000`7ffe0308)],1 00007ffa`4874d030 7503 ntdll!NtWorkerFactoryWorkerReady+0x15 jne (00007ffa`4874d035) Branch ntdll!NtWorkerFactoryWorkerReady+0x12: 00007ffa`4874d032 0f05 syscall 00007ffa`4874d034 c3 ret ntdll!NtWorkerFactoryWorkerReady+0x15: 00007ffa`4874d035 cd2e 2Eh int 00007ffa`4874d037 c3 ret

So, by sorting the functions by the addresses, we have the SSN. for the code, look at <u>MDSec</u> (8. Sorting by System Call Address) blog or see <u>FreshlyCalls</u> implementation. The execution of the system call is not direct by calling syscall instruction. Instead. It uses the method explained below. Briefly, it uses the syscall instructions from ntdll.

## Indirect syscalls

All the methods described are workarounds to get the system call number without getting caught. syscall instruction reveals that some suspicious activity is going on. This is done using KPROCESS!InstrumentationCallback in windows.

```
0:030> dt _kprocess
ntdll!_KPROCESS
+0x000 Header :
_DISPATCHER_HEADER
...
+0x3d8 InstrumentationCallback : Ptr64
Void
...
+0x3f8 EndPadding : [8] Uint8B
```

Any time the windows is done with a syscall and returns to user-mode, it checks this member it is not NULL, the execution will be transferred to that pointer. To check if the syscall is legit, the return address after finishing the syscall is checked to see if it is not from a valid place. If the address is in the address space of the process running, it's not a legitimate place to make a syscall. This check was done by ScyllaHide to detect manual syscalls, the source code can be found <u>here</u>.

```
if (InterlockedOr(TlsGetInstrumentationCallbackDisabled(), 0x1) == 0x1)
        return ReturnVal; // Do not recurse
    const PVOID ImageBase = NtCurrentPeb()->ImageBaseAddress;
    const PIMAGE_NT_HEADERS NtHeaders = RtlImageNtHeader(ImageBase);
    if (NtHeaders != nullptr && ReturnAddress >= (ULONG_PTR)ImageBase &&
        ReturnAddress < (ULONG_PTR)ImageBase + NtHeaders-
>OptionalHeader.SizeOfImage)
    {
        // Syscall return address within the exe file
        ReturnVal = (ULONG_PTR)(ULONG)STATUS_PORT_NOT_SET;
        // Uninstall ourselves after we have completed the sequence { NtOIP,
NtQIP }. More NtSITs will follow but we can't do anything about them
        NumManualSyscalls++;
        if (NumManualSyscalls >= 2)
        {
            InstallInstrumentationCallbackHook(NtCurrentProcess, TRUE);
        }
    }
    InterlockedAnd(TlsGetInstrumentationCallbackDisabled(), 0);
    return ReturnVal;
}
```

It checks the return address of the successful system call. If it resides on the address space of the binary we are running, it is an indication of manual system call.

**The Solution** The solution to this hooking method is done by <u>Bouncy Gate</u> and <u>Recycled</u> <u>Gate</u> method. The idea is quite simple, it is an adjusted version of <u>Hell's Gate</u>. Instead of directly executing <u>syscall</u> instruction and getting caught by static signatures and system call callbacks described above, the author replaces the <u>syscall</u> instruction with a trampoline jump (JMP) to a <u>syscall</u> instruction address from <u>ntdll.dll</u>. now there is no direct <u>syscall</u> instruction and the system call originated from a legitimate place <u>ntdll</u>. This is also implemented in <u>SysWhispers3</u>. To get the address of the syscall instruction in <u>ntdll</u> we can parse the export table and search for syscall, ret opcodes OF <u>05</u> OC or the constant pattern of syscalls in ntdll can be used to get the syscall address. If the function is not hooked, the syscall instruction is on offset 0x12 from the function's address, we can verify that by comparing the opcodes.

# Indirect syscalls in Cobalt Strike

The sample from <u>Dodo's blog</u> Where he already analyzed how indirect syscalls implemented in Cobalt Strike. for easy access, here is <u>UnpacMe Results</u>

<u>020b20098f808301cad6025fe7e2f93fa9f3d0cc5d3d0190f27cf0cd374bcf04</u>. The sample is packed. The unpacking process is easy. Just put a breakpoint on VirtualProtect and get the base address (First Argument). Function sub\_18001B6B0 contains the important part, system call SSN retrieving and execution methods. You can get to this function by following the call instruction to rax which contains a qword memory area or a call to the qword directly. These locations are populated with addresses of the required APIs in this function. We can see multiple calls to sub\_18001A73C with arguments: qword\_\*, a hash (such as 0B12B7A69h), variable passed to the function sub\_18001A7F4 and another allocated memory which is also passed to sub\_18001A7F4.

.text:00000018001B6E5	and	[rbp+1E70h+arg_0], 0
.text:00000018001B6ED	mov	edi, 1F4h
.text:000000018001B6F2	lea	r8, [rbp+1E70h+arg_0]
.text:00000018001B6F9	lea	<pre>rcx, [rsp+1F70h+var_1F40] ; void *</pre>
.text:000000018001B6FE	mov	edx, edi
.text:000000018001B700	call	sub_18001A7F4
.text:00000018001B705	mov	rbx, [rbp+1E70h+arg_0]
.text:00000018001B70C	lea	rax, qword_18004FFA0
.text:000000018001B713	lea	rcx, [rsp+1F70h+var_1F40]
.text:00000018001B718	mov	r8, rbx
.text:00000018001B71B	mov	r9d, 0B12B7A69h
.text:00000018001B721	mov	edx, edi
.text:00000018001B723	mov	[rsp+1F70h+var_1F50], rax
.text:00000018001B728	call	sub_18001A73C
.text:00000018001B72D	lea	rax, qword_18004FFB8
.text:00000018001B734	lea	rcx, [rsp+1F70h+var_1F40]
.text:00000018001B739	mov	r9d, 0C508CF8Bh
.text:00000018001B73F	mov	r8, rbx
.text:00000018001B742	mov	edx, edi
.text:00000018001B744	mov	[rsp+1F70h+var_1F50], rax
.text:00000018001B749	call	sub_18001A73C
.text:00000018001B74E	lea	rax, qword_18004FFD0
.text:00000018001B755	lea	rcx, [rsp+1F70h+var_1F40]
.text:00000018001B75A	mov	r9d, 35AF2123h
.text:00000018001B760	mov	r8, rbx
.text:00000018001B763	mov	edx, edi
.text:00000018001B765	mov	[rsp+1F70h+var_1F50], rax
.text:00000018001B76A	call	sub_18001A73C
.text:00000018001B76F	lea	rax, qword_18004FFE8
.text:00000018001B776	lea	rcx, [rsp+1F70h+var_1F40]
.text:00000018001B77B	mov	r9d, 085988908h
.text:00000018001B781	mov	r8, rbx
.text:00000018001B784	mov	edx, edi
.text:00000018001B786	mov	[rsp+1F70h+var_1F50], rax
.text:00000018001B78B	call	sub_18001A73C
.text:00000018001B790	lea	rax, qword_180050000
.text:00000018001B797	lea	rcx, [rsp+1F70h+var_1F40]
.text:00000018001B79C	mov	r9d, 765EF075h
.text:00000018001B7A2	mov	r8, rbx
.text:00000018001B7A5	mov	edx, edi
.text:00000018001B7A7	mov	[rsp+1F70h+var_1F50], rax
.text:00000018001B7AC	call	sub_18001A73C

Function sub\_18001A73C is to resolve the function address (syscall stub address) by the hash. And function sub\_18001A7F4 used to populate the list with the system call SSN and system call stub. So, sub\_18001A7F4 is our target. In the following picture is the beginning of the function.

sub	rsp, 30h
mov	rax, gs:30h
mov	r13, rcx
mov	r8d, 20202020h
mov	<pre>r9, [rax+_TEB.ProcessEnvironmentBlock]</pre>
mov	r10, [r9+_PEB.Ldr]
add	r10, _PEB_LDR_DATA.InLoadOrderModuleList
mov	r9, [r10]
	; CODE XREF: sub_18001A7F4+5B↓j
	; sub_18001A7F4+65↓j
cmp	r9, r10
jz	loc_18001A9E4
	mov mov mov add mov

The function starts with getting a pointer to the first entry in InLoadOrderModuleList structure by going through reading the Process Environment Block (PEB). here in the picture, r10 is holding the current entry of the structure and r9 is like a variable to get each

entry, this is the breaking condition of the loop as the <u>LIST\_ENTRY</u> structure wrap around itself (doubly linked list).

The next step is to get the Export directory of ntdll.dll but first, get ntdll address in memory.

.text:00000018001A83B	mov	rdi, [r9+LDR DATA TABLE ENTRY.DllBase]
.text:00000018001A83F	mov	r9, [r9]
.text:00000018001A842	movsxd	rax, [rdi+IMAGE DOS HEADER.e lfanew]
.text:00000018001A846	mov	ecx, [rax+rdi+IMAGE NT HEADERS.OptionalHeader.DataDirectory.VirtualAddress]
.text:000000018001A84D	test	ecx, ecx
.text:000000018001A84F	iz	short loc 18001A832
.text:00000018001A851	lea	rsi, [rdi+rcx] ; data dir RVA + BaseAddress = address of IMAGE DATA DIRECTORY
.text:00000018001A851		; First member is IMAGE EXPORT DIRECTORY
.text:00000018001A855	cmp	[rsi+IMAGE_EXPORT_DIRECTORY.NumberOfNames], 0
.text:00000018001A859	jz	short loc_18001A832
.text:00000018001A85B	mov	ecx, [rsi+IMAGE_EXPORT_DIRECTORY.Name]
.text:00000018001A85E	mov	eax, [rcx+rdi] ; DLL name
.text:00000018001A861	or	eax, r8d
.text:00000018001A864	стр	eax, 'ldtn'
.text:000000018001A869	jnz	short loc_18001A832
.text:00000018001A86B	mov	eax, [rcx+rdi+4]
.text:00000018001A86F	or	eax, r8d
.text:00000018001A872	cmp	eax, 'ld.l'
.text:00000018001A877	jnz	short loc_18001A832
.text:00000018001A879	movzx	eax, word ptr [rcx+rdi+8]
.text:00000018001A87E	or	ax, 20h ; ' '
.text:00000018001A882	cmp	ax, 'l'
.text:00000018001A886	jnz	short loc_18001A832
.text:00000018001A888	mov	ebx, [rsi+IMAGE_EXPORT_DIRECTORY.AddressOfFunctions]
.text:00000018001A88B	mov	<pre>ebp, [rsi+IMAGE_EXPORT_DIRECTORY.AddressOfNames]</pre>
.text:00000018001A88E	mov	<pre>eax, [rsi+IMAGE_EXPORT_DIRECTORY.AddressOfNameOrdinals]</pre>
.text:00000018001A891	mov	r8d, edx
.text:00000018001A894	add	rax, rdi
.text:00000018001A897	add	rbx, rdi
.text:00000018001A89A	add	rbp, rdi
.text:00000018001A89D	mov	rcx, r13 ; void *
.text:00000018001A8A0	shl	r8, 4 ; Size
.text:00000018001A8A4	xor	edx, edx ; Val
.text:00000018001A8A6	mov	[rsp+68h+AddressOfFunctions], rbx
.text:00000018001A8AB	mov	[rsp+68h+AddressOfNameOrdinals], rax
.text:00000018001A8B3	mov	[rsp+68h+AddressOfNames], rbp
.text:00000018001A8B8	xor	r15d, r15d
.text:00000018001A8BB	call	memset

It is looking for the right module in the InLoadOrderModuleList by going through each entry, the flink is a pointer to LDR\_DATA\_TABLE\_ENTRY where we can get a pointer to the module. By parsing the module (going through PE file headers) to get the name of the DLL which resides in the Export directory (First member) which is the first member of IMAGE\_DATA\_DIRECTORY structure. It is then tested to see if it is the target module (ntdll). If the module is ntdll, it saves a pointer to AddressOfFunctions, AddressOfNames and AddressOfNameOrdinals. A memory region of size 0x1f40 is then zeroed as it will hold the structures of the system call information needed. The next part is checking the function prefix Ki and Zw. It looks for only one function prefixed by Ki with the hash 8DCD4499h, but I couldn't find function with this hash (using debugger). Then, a call to a hashing function is made. The hashing function is simple.

.text:00000018001A7C8	sub_18001A7C8	proc nea	ar	; CODE	XREF: sub_18001A7F4+	+F5↓p
.text:00000018001A7C8				; sub_1	.8001A7F4+138↓p	
<pre>4.text:000000018001A7C8</pre>		xor	r9d, r9d			
.text:000000018001A7CB		mov	r8, rcx			
.text:00000018001A7CE		mov	eax, 52964EE9h			
.text:000000018001A7D3		cmp	[rcx], r9b			
.text:000000018001A7D6		jz	short locret_180	001A7F2		
.text:000000018001A7D8						
.text:000000018001A7D8	loc_18001A7D8:			; CODE	XREF: sub_18001A7C8+	+28 <b>↓</b> j
.text:000000018001A7D8		movzx	ecx, word ptr [	rcx]		
.text:000000018001A7DB		mov	edx, eax			
.text:000000018001A7DD		inc	r9d			
.text:000000018001A7E0		ror	edx, 8			
.text:000000018001A7E3		add	edx, ecx			
.text:000000018001A7E5		mov	ecx, r9d			
.text:000000018001A7E8		add	rcx, r8			
.text:000000018001A7EB		xor	eax, edx			
.text:000000018001A7ED		стр	byte ptr [rcx],			
.text:000000018001A7F0		jnz	short loc_18001/	A7D8		
.text:000000018001A7F2						
.text:000000018001A7F2	locret_18001A7F	2:		; CODE	XREF: sub_18001A7C8+	⊧E↑j
.text:000000018001A7F2		rep ret	n			
.text:000000018001A7F2	sub_18001A7C8	endp				
.text:000000018001A7F2						

It uses 0x52964EE9 as an initial key value to start the process then:

- Get 2-bytes of the Function name (little endian).
- Rotate the key by 8 (2 characters).
- Add the key and the 2-bytes of the name.
- Increment the counter by 1 (Resulting that all the chars in between the start and end taken two times in the calculation for example ZwOpenProcess will take wz in the first iteration and Ow in the second and so on).
- The result of the addition is XORed with the key to produce the new key. The hash value returned is the last result of the XOR operation.

The resulting value is stored in the following form, in the pre-allocated space.

Address	He	•															ASCII
00000000013D6C0	62	7F	CD	64	40	19	05	00	40	19	62	77	00	00	00	00	b.id@@.bw
																	jPP.bw
00000000013D6E0																	
00000000013D6F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000000013D700	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000000013D710	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
00000000013D720	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

- The first **DWORD** is the hash.
- The second DWORD is the Relative Virtual Address (RVA) of the system call0.
- The third QWORD is the Virtual Address (VA) of the system call stub (RVA + ntdll Base Address).

So, it can be written as:

```
struct syscall_info {
DWORD API_hash;
DWORD
syscall_stub_RVA;
QWORD
syscall_stub_address;
};
```

After populating the structure with the addresses. The structure elements are being sorted by the RVA of the system call stub (second entry in the structure).

+• • •	.text:00000018001A97F .text:000000018001A983 .text:00000018001A986	lea xor test	r11d, [r15-1] ; counter r10d, r10d r11d, r11d
		jz	short loc_18001A9E4
	.text:000000018001A988 .text:000000018001A988 loc 18001A988:		; CODE XREF: sub 18001A7F4+1EE↓j
e		mov	
•		sub	eax, r10d
		dec	
i secci i		jz	short loc_18001A9DC
		lea	r8, [r13+syscall_info.syscall_stub_RVA]
		lea	<pre>r9, [r13+(syscall_info.syscall_stub_RVA+10h)] ; RVA of the Next element in the structure ebx, eax</pre>
1 ( )		mov	ebx, eax
	.text:000000018001A99F loc 18001A99F:		; CODE XREF: sub 18001A7F4+1E6↓j
1 1 m = ▶●	.text:000000018001A99F	mov	esi, [r8]
•		mov	edi, [r9]
•		cmp	esi, edi
		jbe	short loc_18001A9CF
		mov	eax, [r9-4]
		mov	edx, [r8-4]
		mov	rcx, [r8+4]
		mov	[r8-4], eax
•		mov	rax, [r9+4]
•		mov mov	[r8], edi [r9-4], edx
•		mov	[r9], esi
•		mov	[r9+4], rcx
•		mov	[r8+4], rax
	.text:000000018001A9CF loc_18001A9CF:		; CODE XREF: sub_18001A7F4+1B3^j
		add	
	.text:000000018001A9D3	add	r8, 10h
		dec	
		jnz	short loc_18001A99F
1.1	.text:00000018001A9DC		
· · · · · • •	.text:000000018001A9DC loc_18001A9DC: .text:000000018001A9DC	inc	; CODE XREF: sub_18001A7F4+19F↑j r10d
•		cmp	r10d, r11d
•		jb	short loc 18001A98B
	1 1 000000100011054		

After the sorting algorithm is done, the memory structure look like the following:

Address	Hex	ASCII
00000000012D300		w.ª;@@.Nw
00000000012D310	FD D6 91 2E 50 13 05 00 50 13 4E 77 00 00 00 00	ýÖPP.Nw
00000000012D320	B3 D8 6C F1 60 13 05 00 60 13 4E 77 00 00 00 00	*Ø1ñ``.Nw
00000000012D330	FB FD 5E C5 70 13 05 00 70 13 4E 77 00 00 00 00	ûý^Âpp.Nw
00000000012D340	8A 36 BD A7 80 13 05 00 80 13 4E 77 00 00 00 00	. 6½§Nw
00000000012D350	CB EF 59 D1 90 13 05 00 90 13 4E 77 00 00 00 00	ËïYÑNw
00000000012D360	96 44 C7 9C A0 13 05 00 A0 13 4E 77 00 00 00 00	.DÇNw
00000000012D370	D1 EC 6F 3D B0 13 05 00 B0 13 4E 77 00 00 00 00	Ňìo=°°.Nw
00000000012D380	DC C9 B3 2A C0 13 05 00 C0 13 4E 77 00 00 00 00	ÜÉ**ÀÀ.Nw
00000000012D390	AA 8F 35 A4 D0 13 05 00 D0 13 4E 77 00 00 00 00	a.5¤DD.Nw
00000000012D3A0		ï¤Inàà.Nw
00000000012D3B0	7A E9 E8 C0 F0 13 05 00 F0 13 4E 77 00 00 00 00	zéèÀðð.Nw
00000000012D3C0	E1 1F 52 00 00 14 05 00 00 14 4E 77 00 00 00 00	á.RNw
00000000012D3D0	33 FC BF 84 10 14 05 00 10 14 4E 77 00 00 00 00	3ü¿Nw
00000000012D3E0		ËÆBINw
00000000012D3F0	91 62 0D 79 30 14 05 00 30 14 4E 77 00 00 00 00	.b.y00.Nw
00000000012D400	A9 D7 1B EA 40 14 05 00 40 14 4E 77 00 00 00 00	0x. è@@. Nw
000000000120410	49 DE 70 ED 50 14 05 00 50 14 4E 77 00 00 00 00	ODDVP P NW

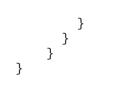
The first address is the address to the Lowest address ZwMapUserPhysicalPagesScatter (Could be different at newer versions of windows) at address 00000000774E1340 If we see the system call SSN of it:

00000000774E1340	4C:8BD1	mov r10,rcx	ZwMapUserPhysicalPagesScatter
00000000774E1343	B8 00000000	mov eax,0	
00000000774E1348	0F05	syscall	
00000000774E134A	C3	ret	

system call number is zero. This is how it gets the SSN for any function, by iterating the structure to get the right hash, the counter will be used to get the SSN (SSN = counter). So far, this is remarkably like <u>MDSec</u> (8. Sorting by System Call Address) implementation of the technique known as <u>FreshlyCalls</u>. We could rewrite the technique using MDSec implementation as follows:

```
#define RVA2VA(Type, DllBase, Rva) (Type)((ULONG_PTR) DllBase + Rva)
static
void
GetSyscallList(PSYSCALL_LIST List) {
    PPEB_LDR_DATA
                            Ldr;
    PLDR_DATA_TABLE_ENTRY LdrEntry;
    PIMAGE_DOS_HEADER DosHeader;
PIMAGE_NT_HEADERS NtHeaders;
                           i, j, NumberOfNames, VirtualAddress, Entries=0;
    DWORD
    PIMAGE_DATA_DIRECTORY DataDirectory;
    PIMAGE_EXPORT_DIRECTORY ExportDirectory;
                            Functions;
    PDWORD
    PDWORD
                            Names;
    PWORD
                            Ordinals;
                            DllName, FunctionName;
    PCHAR
    PVOID
                            DllBase;
    PSYSCALL_ENTRY
                            Table;
    SYSCALL_ENTRY
                            Entry;
    11
    // Get the DllBase address of NTDLL.dll
    // NTDLL is not guaranteed to be the second in the list.
    // so it's safer to loop through the full list and find it.
    Ldr = (PPEB_LDR_DATA)NtCurrentTeb()->ProcessEnvironmentBlock->Ldr;
    // For each DLL loaded
    for (LdrEntry=(PLDR_DATA_TABLE_ENTRY)Ldr->Reserved2[1];
         LdrEntry->DllBase != NULL;
         LdrEntry=(PLDR_DATA_TABLE_ENTRY)LdrEntry->Reserved1[0])
    {
      DllBase = LdrEntry->DllBase;
      DosHeader = (PIMAGE_DOS_HEADER)DllBase;
      NtHeaders = RVA2VA(PIMAGE_NT_HEADERS, DllBase, DosHeader->e_lfanew);
      DataDirectory = (PIMAGE_DATA_DIRECTORY)NtHeaders-
>OptionalHeader.DataDirectory;
      VirtualAddress =
DataDirectory[IMAGE_DIRECTORY_ENTRY_EXPORT].VirtualAddress;
      if(VirtualAddress == 0) continue;
      ExportDirectory = (PIMAGE_EXPORT_DIRECTORY) RVA2VA(ULONG_PTR, DllBase,
VirtualAddress);
      DllName = RVA2VA(PCHAR,DllBase, ExportDirectory->Name);
```

```
if((*(ULONG*)DllName | 0x20202020) != 'ldtn') continue;
      if((*(ULONG*)(DllName + 4) | 0x20202020) == 'ld.l') break;
    }
    NumberOfNames = ExportDirectory->NumberOfNames;
    Functions = RVA2VA(PDWORD,DllBase, ExportDirectory->AddressOfFunctions);
              = RVA2VA(PDWORD,DllBase, ExportDirectory->AddressOfNames);
    Names
    Ordinals = RVA2VA(PWORD, DllBase, ExportDirectory->AddressOfNameOrdinals);
    Table
              = List->Table;
    do {
      FunctionName = RVA2VA(PCHAR, DllBase, Names[NumberOfNames-1]);
      if(*(USHORT*)FunctionName == 'iK' && HashSyscall(FunctionName) ==
0x8DCD4499) {
        Table[Entries].API_Hash = HashSyscall(&FunctionName);
        Table[Entries].syscall_stub_RVA = Functions[Ordinals[NumberOfNames-1]];
        Table[Entries].syscall_stub_address = RVA2VA(void,
DllBase,Functions[Ordinals[NumberOfNames-1]]);
        Entries++;
        if(Entries == MAX_SYSCALLS) break;
      if(*(USHORT*)FunctionName == 'wZ') {
        Table[Entries].API_Hash = HashSyscall(&FunctionName);
        Table[Entries].syscall_stub_RVA = Functions[Ordinals[NumberOfNames-1]];
        Table[Entries].syscall_stub_address = RVA2VA(void,
DllBase,Functions[Ordinals[NumberOfNames-1]]);
        Entries++;
        if(Entries == MAX_SYSCALLS) break;
      }
    } while (--NumberOfNames);
    11
    // Save total number of system calls found.
    11
    List->Entries = Entries;
    11
    // Sort the list by address in ascending order.
    11
    for(i=0; i<Entries - 1; i++) {</pre>
      for(j=0; j<Entries - i - 1; j++) {</pre>
        if(Table[j].syscall_stub_RVA > Table[j+1].syscall_stub_RVA) {
          11
          // Swap entries.
          11
          Entry.Hash = Table[j].Hash;
          Entry.Address = Table[j].Address;
          Table[j].API_Hash = Table[j+1].API_Hash;
          Table[j].syscall_stub_RVA = Table[j+1].syscall_stub_RVA;
          Table[j].syscall_stub_address = Table[j+1].syscall_stub_address;
          Table[j+1].API_Hash = Entry.API_Hash;
          Table[j+1].syscall_stub_RVA = Entry.syscall_stub_RVA;
          Table[j+1].syscall stub_address = Entry.syscall_stub_address;
```



The next thing is to use the structure to get the SSN. and syscall instruction to call. This is done by function sub\_18001A73C.

.text:00000018001A73C .text:00000018001A73E .text:000000018001A740 .text:000000018001A745 .text:000000018001A746 .text:000000018001A74A .text:000000018001A74A	test jz mov push sub xor mov	edx, edx short locret_18001A788 [rsp+counter], rbx rdi rsp, 20h edi, edi rax, rcx
.text:00000018001A74F .text:00000018001A74F <b>loc_18001A74F:</b> .text:00000018001A74F .text:00000018001A752 .text:00000018001A754 .text:000000018001A756 .text:00000018001A756 .text:00000018001A755 .text:00000018001A756 .text:00000018001A756 .text:00000018001A756	cmp jz inc add cmp jb jmp	; CODE XREF: get_api+20↓j [rax], r9d short loc_18001A760 edi rax, 10h edi, edx short loc_18001A74F short loc_18001A781
.text:000000018001A760 ;	mov add mov mov call	; CODE XREF: get_api+16†j eax, edi rax, rax rbx, [rcx+rax*8+8] ; API address rcx, rbx get_syscall_ret_address

The function takes the following parameters:

- The array of structures that has the system call info (called syscall\_info above)
- constant value 0x1F4 the maximum length of the structure members (structure size = 0x1F4 \* 0x10).
- Pre-Allocated memory
- The function hash.
- Global variable to get the system call SSN and stub. The function is simple, it searches the populated structure to find the given hash. If it's found, the counter value is taken and to get the Address of the system call stub. To get the address, the base address of the structure is added to the offset multiplied by 0x10 (struct size) and add 8 to get the last QWORD.

```
API_Address = *(STRUCT_BASE_ADDR + COUNTER * 0x10
+ 8)
```

The address the passed to get\_syscall\_ret\_address to get the syscall ret addresses to use it to execute the system call to bypass the callback mentioned before (call stack tracing is be used to detect this trick).

.text:000000018001A78C		
.text:000000018001A78C	mov	[rsp+arg_8], 50Fh ; syscall
.text:000000018001A793	MOVZX	r8d, [rsp+arg_8]
.text:000000018001A799	mov	r9b, 0C3h ; ret
.text:000000018001A79C	xor	eax, eax
.text:000000018001A79E		
.text:000000018001A79E lc	bc_18001A79E:	; CODE XREF: get_syscall_ret_address+28↓j
.text:000000018001A79E	movsxd	rdx, eax
.text:000000018001A7A1	стр	r8w, [rdx+rcx]
.text:000000018001A7A6	jnz	short loc_18001A7AF
.text:000000018001A7A8	стр	r9b, [rdx+rcx+2]
.text:000000018001A7AD	jz	short loc_18001A7B9
.text:000000018001A7AF		
.text:000000018001A7AF lc	bc_18001A7AF:	; CODE XREF: get_syscall_ret_address+1A^j
.text:000000018001A7AF	inc	eax
.text:000000018001A7B1	стр	eax, 20h ; ' '
.text:000000018001A7B4	j1	short loc_18001A79E
.text:000000018001A7B6	xor	eax, eax
.text:000000018001A7B8	retn	
.text:000000018001A7B9 ;		
.text:000000018001A7B9		
.text:000000018001A7B9 lc	bc_18001A7B9:	; CODE XREF: get_syscall_ret_address+21^j
.text:000000018001A7B9	cdqe	
.text:000000018001A7BB	add	rax, rcx
.text:000000018001A7BE	retn	
.text:00000018001A7BE ge	et_syscall_ret_address	; endp

The global variable is used to store:

- QWORD to store System call address (function address at ntdll)
- QWORD to store syscall, ret instruction sequence address.
- DWORD to store system call number SSN. We can rewrite it as follows:

```
struct
syscall_required_addresses {
QWORD syscall_stub_address;
QWORD
syscall_intruction_address;
DWORD syscall_number;
};
```

(Creative names I know :) )



There are some choices to call the required function. This is done based on the value at a global variable (0x18004BC6C):

- 1 : Direct call using the first member of the structure (Address of the function in ntdll)
- 2 : Indirect system call using trampoline jump using the system call number and the syscall address stored before.

```
.text:00000018001D0BF ; __int64 sub_18001D0BF()
.text:00000018001D0BF sub_18001D0BF proc near ; CODE XREF: sub_18001B11C+5C↑p
.text:00000018001D0BF mov r11, cs:syscall_inst_addr
.text:00000018001D0C6 mov eax, cs:SSN
.text:00000018001D0CC mov r10, rcx
.text:00000018001D0CF jmp r11
.text:00000018001D0CF sub_18001D0BF endp
.text:00000018001D0CF
```

anything else: Direct call to Win32 API.

# **Detecting syscalls**

System calls can be used to bypass user mood hooks but there are other methods to detect Direct and Indirect syscalls. To detect Direct system calls, Windows provides a large set of callback functions, one of them is KPROCESS!InstrumentationCallback. This callback is triggered whenever the system returns from the kernel mode to user mode. This could be used to check the return address of the syscall which reveals the location of syscall instruction execution. This location should be ntdll but in case of the direct system calls, it will be from the .text section of the PE file. This was used by ScyllaHide. Indirect system calls solved this problem by getting the address of syscall instruction in ntdll and jump to it. To detect indirect syscalls the call stack tracing method can be used to check from where the system call originated -before jumping to ntdll-. This also can be bypassed by creating a new thread to get a new call stack using callback functions like TpAllocWork and RtlQueueWorkItem. If you want to know more about this, you can read Hiding In PlainSight 1&2

Note: This was personal notes I wrote when I was learning about syscalls, if there's anything not accurate, please let me know

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