

Abusing undocumented features to spoof PE section headers

 secret.club/2023/06/05/spoof-pe-sections.html

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Introduction

Some time ago, I accidentally came across some interesting behaviour in PE files while debugging an unrelated project. I noticed that setting the `SectionAlignment` value in the NT header to a value lower than the page size (4096) resulted in significant differences in the way that the image is mapped into memory. Rather than following the usual procedure of parsing the section table to construct the image in memory, the loader appeared to map the entire file, including the headers, into memory with read-write-execute (RWX) permissions - the individual section headers were completely ignored.

As a result of this behaviour, it is possible to create a PE executable without any sections, yet still capable of executing its own code. The code can even be self-modifying if necessary due to the write permissions that are present by default.

One way in which this mode could potentially be abused would be to create a fake section table - on first inspection, this would appear to be a normal PE module containing read-write/read-only data sections, but when launched, the seemingly NX data becomes executable.

While I am sure that this technique will have already been discovered (and potentially abused) in the past, I have been unable to find any documentation online describing it. MSDN does briefly mention that the `SectionAlignment` value can be less than the page size, but it doesn't elaborate any further on the implications of this.

Inside the Windows kernel

A quick look in the kernel reveals what is happening. Within `MiCreateImageFileMap`, we can see the parsing of PE headers - notably, if the `SectionAlignment` value is less than 0x1000, an undocumented flag (0x200000) is set prior to mapping the image into memory:

```

if(v29->SectionAlignment < 0x1000)
{
    if((SectionFlags & 0x80000) != 0)
    {
        v17 = 0xC000007B;
        MiLogCreateImageFileMapFailure(v36, v39, *(unsigned int *)
(v29 + 64), DWORD1(v99));
        ImageFailureReason = 55;
        goto LABEL_81;
    }
    if(!MiLegacyImageArchitecture((unsigned __int16)v99))
    {
        v17 = 0xC000007B;
        ImageFailureReason = 56;
        goto LABEL_81;
    }
    SectionFlags |= 0x200000;
}
v40 = MiBuildImageControlArea(a3, v38, v29, (unsigned int)&v99, SectionFlags,
(__int64)&FileSize, (__int64)&v93);

```

If the aforementioned flag is set, `MiBuildImageControlArea` treats the entire file as one single section:

```

if((SectionFlags & 0x200000) != 0)
{
    SectionCount = 1;
}
else
{
    SectionCount = a4->NumberOfSections + 1;
}
v12 = MiAllocatePool(64, 8 * (7 * SectionCount + (((unsigned __int64)
(unsigned int)MiFlags >> 13) & 1)) + 184, (SectionFlags & 0x200000) != 0 ? 0x61436D4D
: 0x69436D4D);

```

As a result, the raw image is mapped into memory with all PTEs assigned `MM_EXECUTE_READWRITE` protection. As mentioned previously, the `IMAGE_SECTION_HEADER` list is ignored, meaning a PE module using this mode can have a `NumberOfSections` value of 0. There are no obvious size restrictions on PE modules using this mode either - the loader will allocate memory based on the `SizeOfImage` field and copy the file contents accordingly. Any excess memory beyond the size of the file will remain blank.

Demonstration #1 - Executable PE with no sections

The simplest demonstration of this technique would be to create a generic “loader” for position-independent code. I have created the following sample headers by hand for testing:


```
0x10, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00,
```

```
// (code goes here)
```

```
};
```

These headers contain a `SectionAlignment` value of 0x200 (rather than the usual 0x1000), a `SizeOfImage` value of 0x100000 (1MB), a blank section table, and an entry-point positioned immediately after the headers. Aside from these values, there is nothing special about the remaining fields:

```

(DOS Header)
  e_magic           : 0x5A4D
  ...
  e_lfanew         : 0x40
(NT Header)
  Signature        : 0x4550
  Machine          : 0x8664
  NumberOfSections : 0x0
  TimeDateStamp    : 0x0
  PointerToSymbolTable : 0x0
  NumberOfSymbols  : 0x0
  SizeOfOptionalHeader : 0xF0
  Characteristics : 0x22
  Magic            : 0x20B
  MajorLinkerVersion : 0xE
  MinorLinkerVersion : 0x1D
  SizeOfCode       : 0x0
  SizeOfInitializedData : 0x0
  SizeOfUninitializedData : 0x0
  AddressOfEntryPoint : 0x148
  BaseOfCode       : 0x0
  ImageBase        : 0x140000000
  SectionAlignment : 0x200
  FileAlignment    : 0x200
  MajorOperatingSystemVersion : 0x6
  MinorOperatingSystemVersion : 0x0
  MajorImageVersion : 0x0
  MinorImageVersion : 0x0
  MajorSubsystemVersion : 0x6
  MinorSubsystemVersion : 0x0
  Win32VersionValue : 0x0
  SizeOfImage      : 0x100000
  SizeOfHeaders    : 0x148
  CheckSum         : 0x0
  Subsystem        : 0x2
  DllCharacteristics : 0x8160
  SizeOfStackReserve : 0x100000
  SizeOfStackCommit : 0x1000
  SizeOfHeapReserve : 0x100000
  SizeOfHeapCommit : 0x1000
  LoaderFlags      : 0x0
  NumberOfRvaAndSizes : 0x10
  DataDirectory[0] : 0x0, 0x0
  ...
  DataDirectory[15] : 0x0, 0x0
(Start of code)

```

For demonstration purposes, we will be using some position-independent code that calls `MessageBoxA`. As the base headers lack an import table, this code must locate and load all dependencies manually - `user32.dll` in this case. This same payload can be used in both 32-bit and 64-bit environments:

BYTE bMessageBox[939] =

```
{
    0x8B, 0xC4, 0x6A, 0x00, 0x2B, 0xC4, 0x59, 0x83, 0xF8, 0x08, 0x0F, 0x84,
    0xA0, 0x01, 0x00, 0x00, 0x55, 0x8B, 0xEC, 0x83, 0xEC, 0x3C, 0x64, 0xA1,
    0x30, 0x00, 0x00, 0x00, 0x33, 0xD2, 0x53, 0x56, 0x57, 0x8B, 0x40, 0x0C,
    0x33, 0xDB, 0x21, 0x5D, 0xF0, 0x21, 0x5D, 0xEC, 0x8B, 0x40, 0x1C, 0x8B,
    0x00, 0x8B, 0x78, 0x08, 0x8B, 0x47, 0x3C, 0x8B, 0x44, 0x38, 0x78, 0x03,
    0xC7, 0x8B, 0x48, 0x24, 0x03, 0xCF, 0x89, 0x4D, 0xE8, 0x8B, 0x48, 0x20,
    0x03, 0xCF, 0x89, 0x4D, 0xE4, 0x8B, 0x48, 0x1C, 0x03, 0xCF, 0x89, 0x4D,
    0xF4, 0x8B, 0x48, 0x14, 0x89, 0x4D, 0xFC, 0x85, 0xC9, 0x74, 0x5F, 0x8B,
    0x70, 0x18, 0x8B, 0xC1, 0x89, 0x75, 0xF8, 0x33, 0xC9, 0x85, 0xF6, 0x74,
    0x4C, 0x8B, 0x45, 0xE8, 0x0F, 0xB7, 0x04, 0x48, 0x3B, 0xC2, 0x74, 0x07,
    0x41, 0x3B, 0xCE, 0x72, 0xF0, 0xEB, 0x37, 0x8B, 0x45, 0xE4, 0x8B, 0x0C,
    0x88, 0x03, 0xCF, 0x74, 0x2D, 0x8A, 0x01, 0xBE, 0x05, 0x15, 0x00, 0x00,
    0x84, 0xC0, 0x74, 0x1F, 0x6B, 0xF6, 0x21, 0x0F, 0xBE, 0xC0, 0x03, 0xF0,
    0x41, 0x8A, 0x01, 0x84, 0xC0, 0x75, 0xF1, 0x81, 0xFE, 0xFB, 0xF0, 0xBF,
    0x5F, 0x75, 0x74, 0x8B, 0x45, 0xF4, 0x8B, 0x1C, 0x90, 0x03, 0xDF, 0x8B,
    0x75, 0xF8, 0x8B, 0x45, 0xFC, 0x42, 0x3B, 0xD0, 0x72, 0xA9, 0x8D, 0x45,
    0xC4, 0xC7, 0x45, 0xC4, 0x75, 0x73, 0x65, 0x72, 0x50, 0x66, 0xC7, 0x45,
    0xC8, 0x33, 0x32, 0xC6, 0x45, 0xCA, 0x00, 0xFF, 0xD3, 0x8B, 0xF8, 0x33,
    0xD2, 0x8B, 0x4F, 0x3C, 0x8B, 0x4C, 0x39, 0x78, 0x03, 0xCF, 0x8B, 0x41,
    0x20, 0x8B, 0x71, 0x24, 0x03, 0xC7, 0x8B, 0x59, 0x14, 0x03, 0xF7, 0x89,
    0x45, 0xE4, 0x8B, 0x41, 0x1C, 0x03, 0xC7, 0x89, 0x75, 0xF8, 0x89, 0x45,
    0xE8, 0x89, 0x5D, 0xFC, 0x85, 0xDB, 0x74, 0x7D, 0x8B, 0x59, 0x18, 0x8B,
    0x45, 0xFC, 0x33, 0xC9, 0x85, 0xDB, 0x74, 0x6C, 0x0F, 0xB7, 0x04, 0x4E,
    0x3B, 0xC2, 0x74, 0x22, 0x41, 0x3B, 0xCB, 0x72, 0xF3, 0xEB, 0x5A, 0x81,
    0xFE, 0x6D, 0x07, 0xAF, 0x60, 0x8B, 0x75, 0xF8, 0x75, 0x8C, 0x8B, 0x45,
    0xF4, 0x8B, 0x04, 0x90, 0x03, 0xC7, 0x89, 0x45, 0xEC, 0xE9, 0x7C, 0xFF,
    0xFF, 0xFF, 0x8B, 0x45, 0xE4, 0x8B, 0x0C, 0x88, 0x03, 0xCF, 0x74, 0x35,
    0x8A, 0x01, 0xBE, 0x05, 0x15, 0x00, 0x00, 0x84, 0xC0, 0x74, 0x27, 0x6B,
    0xF6, 0x21, 0x0F, 0xBE, 0xC0, 0x03, 0xF0, 0x41, 0x8A, 0x01, 0x84, 0xC0,
    0x75, 0xF1, 0x81, 0xFE, 0xB4, 0x14, 0x4F, 0x38, 0x8B, 0x75, 0xF8, 0x75,
    0x10, 0x8B, 0x45, 0xE8, 0x8B, 0x04, 0x90, 0x03, 0xC7, 0x89, 0x45, 0xF0,
    0xEB, 0x03, 0x8B, 0x75, 0xF8, 0x8B, 0x45, 0xFC, 0x42, 0x3B, 0xD0, 0x72,
    0x89, 0x33, 0xC9, 0xC7, 0x45, 0xC4, 0x54, 0x65, 0x73, 0x74, 0x51, 0x8D,
    0x45, 0xC4, 0x88, 0x4D, 0xC8, 0x50, 0x50, 0x51, 0xFF, 0x55, 0xF0, 0x6A,
    0x7B, 0x6A, 0xFF, 0xFF, 0x55, 0xEC, 0x5F, 0x5E, 0x5B, 0xC9, 0xC3, 0x90,
    0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90, 0x90,
    0x48, 0x89, 0x5C, 0x24, 0x08, 0x48, 0x89, 0x6C, 0x24, 0x10, 0x48, 0x89,
    0x74, 0x24, 0x18, 0x48, 0x89, 0x7C, 0x24, 0x20, 0x41, 0x54, 0x41, 0x56,
    0x41, 0x57, 0x48, 0x83, 0xEC, 0x40, 0x65, 0x48, 0x8B, 0x04, 0x25, 0x60,
    0x00, 0x00, 0x00, 0x33, 0xFF, 0x45, 0x33, 0xFF, 0x45, 0x33, 0xE4, 0x45,
    0x33, 0xC9, 0x48, 0x8B, 0x48, 0x18, 0x48, 0x8B, 0x41, 0x30, 0x48, 0x8B,
    0x08, 0x48, 0x8B, 0x59, 0x10, 0x48, 0x63, 0x43, 0x3C, 0x8B, 0x8C, 0x18,
    0x88, 0x00, 0x00, 0x00, 0x48, 0x03, 0xCB, 0x8B, 0x69, 0x24, 0x44, 0x8B,
    0x71, 0x20, 0x48, 0x03, 0xEB, 0x44, 0x8B, 0x59, 0x1C, 0x4C, 0x03, 0xF3,
    0x8B, 0x71, 0x14, 0x4C, 0x03, 0xDB, 0x85, 0xF6, 0x0F, 0x84, 0x80, 0x00,
    0x00, 0x00, 0x44, 0x8B, 0x51, 0x18, 0x33, 0xC9, 0x45, 0x85, 0xD2, 0x74,
    0x69, 0x48, 0x8B, 0xD5, 0x0F, 0x1F, 0x40, 0x00, 0x0F, 0xB7, 0x02, 0x41,
    0x3B, 0xC1, 0x74, 0x0D, 0xFF, 0xC1, 0x48, 0x83, 0xC2, 0x02, 0x41, 0x3B,
    0xCA, 0x72, 0xED, 0xEB, 0x4D, 0x45, 0x8B, 0x04, 0x8E, 0x4C, 0x03, 0xC3,
    0x74, 0x44, 0x41, 0x0F, 0xB6, 0x00, 0x33, 0xD2, 0xB9, 0x05, 0x15, 0x00,
}
```

```

0x00, 0x84, 0xC0, 0x74, 0x35, 0x0F, 0x1F, 0x00, 0x6B, 0xC9, 0x21, 0x8D,
0x52, 0x01, 0x0F, 0xBE, 0xC0, 0x03, 0xC8, 0x42, 0x0F, 0xB6, 0x04, 0x02,
0x84, 0xC0, 0x75, 0xEC, 0x81, 0xF9, 0xFB, 0xF0, 0xBF, 0x5F, 0x75, 0x08,
0x41, 0x8B, 0x3B, 0x48, 0x03, 0xFB, 0xEB, 0x0E, 0x81, 0xF9, 0x6D, 0x07,
0xAF, 0x60, 0x75, 0x06, 0x45, 0x8B, 0x23, 0x4C, 0x03, 0xE3, 0x41, 0xFF,
0xC1, 0x49, 0x83, 0xC3, 0x04, 0x44, 0x3B, 0xCE, 0x72, 0x84, 0x48, 0x8D,
0x4C, 0x24, 0x20, 0xC7, 0x44, 0x24, 0x20, 0x75, 0x73, 0x65, 0x72, 0x66,
0xC7, 0x44, 0x24, 0x24, 0x33, 0x32, 0x44, 0x88, 0x7C, 0x24, 0x26, 0xFF,
0xD7, 0x45, 0x33, 0xC9, 0x48, 0x8B, 0xD8, 0x48, 0x63, 0x48, 0x3C, 0x8B,
0x94, 0x01, 0x88, 0x00, 0x00, 0x00, 0x48, 0x03, 0xD0, 0x8B, 0x7A, 0x24,
0x8B, 0x6A, 0x20, 0x48, 0x03, 0xF8, 0x44, 0x8B, 0x5A, 0x1C, 0x48, 0x03,
0xE8, 0x8B, 0x72, 0x14, 0x4C, 0x03, 0xD8, 0x85, 0xF6, 0x74, 0x77, 0x44,
0x8B, 0x52, 0x18, 0x0F, 0x1F, 0x44, 0x00, 0x00, 0x33, 0xC0, 0x45, 0x85,
0xD2, 0x74, 0x5B, 0x48, 0x8B, 0xD7, 0x66, 0x0F, 0x1F, 0x44, 0x00, 0x00,
0x0F, 0xB7, 0x0A, 0x41, 0x3B, 0xC9, 0x74, 0x0D, 0xFF, 0xC0, 0x48, 0x83,
0xC2, 0x02, 0x41, 0x3B, 0xC2, 0x72, 0xED, 0xEB, 0x3D, 0x44, 0x8B, 0x44,
0x85, 0x00, 0x4C, 0x03, 0xC3, 0x74, 0x33, 0x41, 0x0F, 0xB6, 0x00, 0x33,
0xD2, 0xB9, 0x05, 0x15, 0x00, 0x00, 0x84, 0xC0, 0x74, 0x24, 0x66, 0x90,
0x6B, 0xC9, 0x21, 0x8D, 0x52, 0x01, 0x0F, 0xBE, 0xC0, 0x03, 0xC8, 0x42,
0x0F, 0xB6, 0x04, 0x02, 0x84, 0xC0, 0x75, 0xEC, 0x81, 0xF9, 0xB4, 0x14,
0x4F, 0x38, 0x75, 0x06, 0x45, 0x8B, 0x3B, 0x4C, 0x03, 0xFB, 0x41, 0xFF,
0xC1, 0x49, 0x83, 0xC3, 0x04, 0x44, 0x3B, 0xCE, 0x72, 0x92, 0x45, 0x33,
0xC9, 0xC7, 0x44, 0x24, 0x20, 0x54, 0x65, 0x73, 0x74, 0x4C, 0x8D, 0x44,
0x24, 0x20, 0xC6, 0x44, 0x24, 0x24, 0x00, 0x48, 0x8D, 0x54, 0x24, 0x20,
0x33, 0xC9, 0x41, 0xFF, 0xD7, 0xBA, 0x7B, 0x00, 0x00, 0x00, 0x48, 0xC7,
0xC1, 0xFF, 0xFF, 0xFF, 0xFF, 0x41, 0xFF, 0xD4, 0x48, 0x8B, 0x5C, 0x24,
0x60, 0x48, 0x8B, 0x6C, 0x24, 0x68, 0x48, 0x8B, 0x74, 0x24, 0x70, 0x48,
0x8B, 0x7C, 0x24, 0x78, 0x48, 0x83, 0xC4, 0x40, 0x41, 0x5F, 0x41, 0x5E,
0x41, 0x5C, 0xC3

```

```
};
```

As a side note, several readers have asked how I created this sample code (previously used in another project) which works correctly in both 32-bit and 64-bit modes. The answer is very simple: it begins by storing the original stack pointer value, pushes a value onto the stack, and compares the new stack pointer to the original value. If the difference is 8, the 64-bit code is executed - otherwise, the 32-bit code is executed. While there are certainly more efficient approaches to achieve this outcome, this method is sufficient for demonstration purposes:

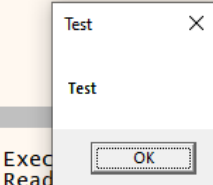
```

mov eax, esp    ; store stack ptr
push 0         ; push a value onto the stack
sub eax, esp   ; calculate difference
pop ecx       ; restore stack
cmp eax, 8    ; check if the difference is 8
je 64bit_code
32bit_code:
xxxx
64bit_code:
xxxx

```

By appending this payload to the original headers above, we can generate a valid and functional EXE file. The provided PE headers contain a hardcoded `SizeOfImage` value of `0x100000` which allows for a maximum payload size of almost 1MB, but this can be increased if necessary. Running this program will display our message box, despite the fact that the PE headers lack any executable sections, or any sections at all in this case:

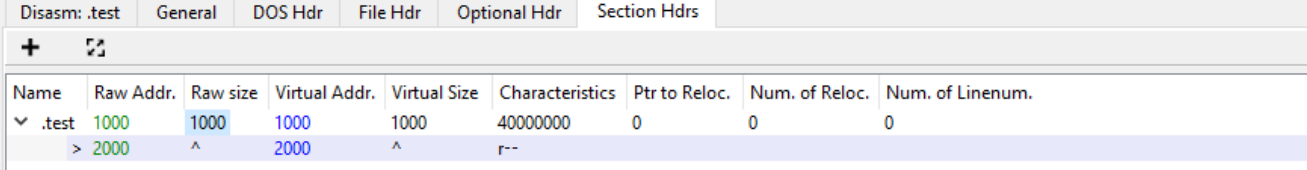
07FF423EA5000	000000000000FB000	Reserved (00007FF423EA0000)	MAP	--	--
07FF423FA0000	0000000100020000	Reserved	PRV	--	--
07FF523FC0000	0000000002000000	Reserved	PRV	--	--
07FF525FC0000	0000000000001000		PRV	-RW--	-RW--
07FF525FD0000	0000000000001000		MAP	-R---	-R---
07FF525FE0000	0000000000023000		MAP	-R---	-R---
07FF7902F0000	0000000001000000	no_section_64.exe	IMG	ERWC-	ERWC-
07FFF13520000	0000000000001000	textshaping.dll	IMG	-R---	ERWC-
07FFF13521000	000000000004B000	".text"	IMG	ER---	ERWC-
07FFF1356C000	000000000005B000	".rdata"	IMG	-R---	ERWC-
07FFF135C7000	0000000000001000	".data"	IMG	-RW--	ERWC-



Demonstration #2 - Executable PE with spoofed sections

Perhaps more interestingly, it is also possible to create a fake section table using this mode as mentioned earlier. I have created another EXE which follows a similar format to the previous samples, but also includes a single read-only section:

0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
1000	8B	C4	6A	00	2B	C4	59	83	F8	08	0F	84	A0	01	00	00	. Ä j . + Ä Y . ø														
1010	55	8B	EC	83	EC	3C	64	A1	30	00	00	00	33	D2	53	56	U . i . i < d ; 0 3 Ò S V														
1020	57	8B	40	0C	33	DB	21	5D	F0	21	5D	EC	8B	40	1C	8B	W . @ . 3 Û ! j ð ! j i . @ . .														
1030	00	8B	78	08	8B	47	3C	8B	44	38	78	03	C7	8B	48	24	. . x . . G < . D 8 x . Ç . H ð														
1040	03	CF	89	4D	E8	8B	48	20	03	CF	89	4D	E4	8B	48	1C	. Ì . M è . H . . Ì . M ä . H .														
1050	03	CF	89	4D	F4	8B	48	14	89	4D	FC	85	C9	74	5F	8B	. Ì . M ô . H . . M ü . È t _														
1060	70	18	8B	C1	89	75	F8	33	C9	85	F6	74	4C	8B	45	E8	p . . Ä . u ø 3 È . ö t L . È è														
1070	0F	E7	04	48	3B	C2	74	07	41	3B	CE	72	F0	EB	37	8B	. . . H ; Ä t . A ; Ì r ø è 7 .														
1080	45	E4	8B	0C	88	03	CF	74	2D	8A	01	BE	05	15	00	00	E ä Ì t - . . ¼														
1090	84	C0	74	1F	6B	F6	21	0F	BE	C0	03	F0	41	8A	01	84	. Ä t . k ö ! . ¼ Ä . ø A . . .														
10A0	C0	75	F1	81	FE	FB	F0	BF	5F	75	74	8B	45	F4	8B	1C	Ä u ñ . þ ú ø ç _ u t . E ö . .														



The main payload has been stored within this read-only section and the entry-point has been updated to `0x1000`. Under normal circumstances, you would expect the program to crash immediately with an access-violation exception due to attempting to execute read-only memory. However, this doesn't occur here - the target memory region contains RWX permissions and the payload is executed successfully:

620000	0000000000001000		MAP	-R---	-R---
630000	0000000000023000		MAP	-R---	-R---
760000	0000000000001000	fake_readonly_section.exe	IMG	ERWC-	ERWC-
761000	0000000000001000	".test"	IMG	ERW--	ERWC-
920000	0000000000001000	kernelbase.dll	IMG	-R---	ERWC-
921000	0000000000133000	".text"	IMG	ER---	ERWC-

Notes

The sample EXE files can be downloaded [here](#).

The proof-of-concepts described above involve appending the payload to the end of the NT headers, but it is also possible to embed executable code within the headers themselves using this technique. The module will fail to load if the `AddressOfEntryPoint` value is less than the `SizeOfHeaders` value, but this can easily be bypassed since the `SizeOfHeaders` value is not strictly enforced. It can even be set to 0, allowing the entry-point to be positioned anywhere within the file.

It is possible that this feature was initially designed to allow for very small images, enabling the headers, code, and data to fit within a single memory page. As memory protection is applied per-page, it makes sense to apply RWX to all PTEs when the virtual section size is lower than the page size - it would otherwise be impossible to manage protections correctly if multiple sections resided within a single page.

I have tested these EXE files on various different versions of Windows from Vista to 10 with success in all cases. Unfortunately it has very little practical use in the real world as it won't deceive any modern disassemblers - nonetheless, it remains an interesting concept.